CMPT 745
Software Engineering

# What are programs? 

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- Team mates
- Compilers
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- Different programs have different requirements
- Performance over all
- Security!


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- Different programs have different requirements
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- Security!
- We cannot reason about programs in only one way


## Program Representation

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0101011
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Why might binaries be good for security tasks?

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「-..
Bar.c Baz.c
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- Difficult models:
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- Relationships can be hard to extract
- Often used when relating to comments or specs

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## Program Representation

- Before we can reason about programs, we must have a vocabulary and a model to analyze
- Difficult models:
- Compiled binaries
- Source code
- A good representation should make explicit the relationships you want to analyze


## Program Representation

Core graph representations for analysis:

1) Abstract Syntax Trees
2) Control Flow Graphs
3) Program Dependence Graphs
4) Call Graphs
5) Points-to Graphs
6) Emerging Representations for ML

## 1) Abstract Syntax Trees

- Lifts the source into a canonical tree form


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\begin{gathered}
\text { for i in range }(5,10): \\
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\end{gathered}
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## 1) Abstract Syntax Trees

- Lifts the source into a canonical tree form
- Internal nodes are operators, statements, etc.

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& \text { for i in range }(5,10): \\
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## 1) Abstract Syntax Trees

- Lifts the source into a canonical tree form
- Internal nodes are operators, statements, etc.
- Leaves are values, variables, operands


## for $i$ in range $(5,10)$ : $\mathrm{a}[\mathrm{i}]=\mathrm{i}$ * 5



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- Used for syntax analysis \& transformation:

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for i in range(5,10):
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- Training prediction/completion models

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## But:

1) The same program may still be spelled many ways
2) Some information is implicit rather than explicit


## 2) Control Flow Graphs

- Express the possible decisions and possible paths through a program

```
cond = input()
if cond:
    a = foo()
else:
    a = bar()
print(a)
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- Express the possible decisions and possible paths through a program
- Basic Blocks (Nodes) are straight line code
- Edges show how decisions can lead to different basic blocks
- Paths through the graph are potential paths through the program

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## 2) Control Flow Graphs (CFGs)

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i = 1
while i < N:
    i = i + 1
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The 'while' is gone

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Why is the ' $i f$ ' in a separate block?


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> What would the CFG of the equivalent 'for' look like?


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- Instruction X depends on Y if Y can influence X
- Nodes are instructions
- An edge $Y \rightarrow X$ shows that $Y$ influences $X$


## 3)Program Dependence Graph (PDG)

- A Program Dependence Graph captures how instructions can influence each other
- Instruction $X$ depends on $Y$ if $Y$ can influence $X$
- 2 main types of influence:
- Data dependence
- Control dependence


## Data Dependence

X data depends on Y if

- There exists a path from $Y$ to $X$ in the CFG



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Preliminary: X dominates Y if

- every path from the entry node to Y passes X
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1) sum $=0$
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3) while i < N:
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5) sum $=$ sum $+i$
6) print (sum)
$\operatorname{IDOM}(6)=$ ?
7) sum = 0
8) $i=1$


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4)i = i + 1
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## Control Dependence

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```
1) sum = 0
```

2) $i=1$
3) while $i$ What does this mean intuitively?
4) i =
5) sum $=$ sum $+i$
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- There exists a path from X to Y s.t. Y post dominates every node between X and Y .
- $Y$ does not strictly post dominate $X$


## Control Dependence (Finally)

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- Definition 1:
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What is CD(5)? CD(3)

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1) sum $=0$
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4) i $=i+1 \%$
5) if $0==i \% 2:$
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What is $\operatorname{CD}(7)$ ?


## Control Dependence

- There exists a path from X to Y s.t. Y post dominates every node between $X$ and $Y$.
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1) if $X$ or $Y$ :
2) $\operatorname{print}(X)$
3) print(Y)

What is $\mathrm{CD}(2)$ ?

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1) if $X$ or $Y$ :
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- Testing: How can I reach a statement?


## 3)Program Dependence Graph(PDG)

## The PDG is the combination of

- The control dependence graph
- The data dependence graph

Recall: Edges identify potential influence

- Debuggin Can you see challenges that may arise
- Security: when using the PDG in practice?
- Testing: How can I reach a statement?


## 4) Call Graph (Multigraph)

- Captures the composition of a program
- Nodes are functions
- Edges show possible calls



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What does this capture?

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How should we handle function pointers?


## 5) Points-to Graphs

Pointers / indirection create two difficult problems:

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- Multiple variables may denote the same memory location



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x.lock()
y.unlock()
```


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- Aliasing
- Multiple variables may denote the same memory location
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- One variable may potentially denote several different targets in memory.
$x . \operatorname{lock}()$
y. $\cdot$ unlock( $)$

```
x = password
...
broadcast(y)
```


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1) }r=C(
2) p.f = r
3) t = C()
4) if ...:
5) rr = p
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- Machine learning is seen as a value driver for many tasks, but using it effectively to reason about software is still challenging


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- Machine learning is seen as a value driver for many tasks, but using it effectively to reason about software is still challenging
- Trying simple models should always be considered first e.g. simple feed forward networks can work better [Yedida 2021]


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- Machine learning is seen as a value driver for many tasks, but using it effectively to reason about software is still challenging
- Trying simple models should always be considered first
- Bug fix \& close time estimation [Yedida 2021]
- Project planning \& analytics [Krishna 2020]
- Recognizing actionable compiler warnings [Yang 2020]


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- Observe:

Many engineering tasks require discrete \& symbolic reasoning.

- ML is classically better on non symbolic problems.


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- Observe:

Many engineering tasks require discrete \& symbolic reasoning.

- ML is classically better on non symbolic problems.
- Bridging the gap is an area of open research (neurosymbolic, ...)
- Solutions that do not require a priori implementation are desirable
- But different models \& pipelines arise to aid in reasoning about software


## 6) Emerging Representations for ML



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neurosymbolic models for synthesis
[Nye 2020]

## 6) Emerging Representations for ML


code2vec
[Alon 2018]

42.77\%
$33.74 \%$
8.86\%


## 6) Emerging Representations for ML

Seq2DRNN
Encoder-Decoders
[Alvarez-Melis 2017, Gu 2019]


## Representing Program Executions

## Execution Representations

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- All possible program behaviors at once
- Usually projected onto the CFG
- Execution representations are dynamic
- Only the behavior of a single real execution
- Multiple instances of an instruction occur multiple times


## Control Flow Trace



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## Dynamic Dependence Graph



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Slicing (static or dynamic) computes a transitive closure of dependences


## Dynamic Dependence Graph



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- Debugging: What caused a bug?
- Security: How did sensitive information leak?
- Testing: What tests need to be run based on a change?

Prioritizing, pruning, \& bundling information is often critical when applying slicing

## Summary

- Different tasks may benefit from representing programs in different ways
- Thinking of the right representation for the task you have is important

