

CMPT 473
Software Testing, Reliability and Security

Symbolic Execution

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Symbolic Execution

- As we have seen, building constraints that model code can be useful

CBMC was able to prove
certain errors couldn't exist!

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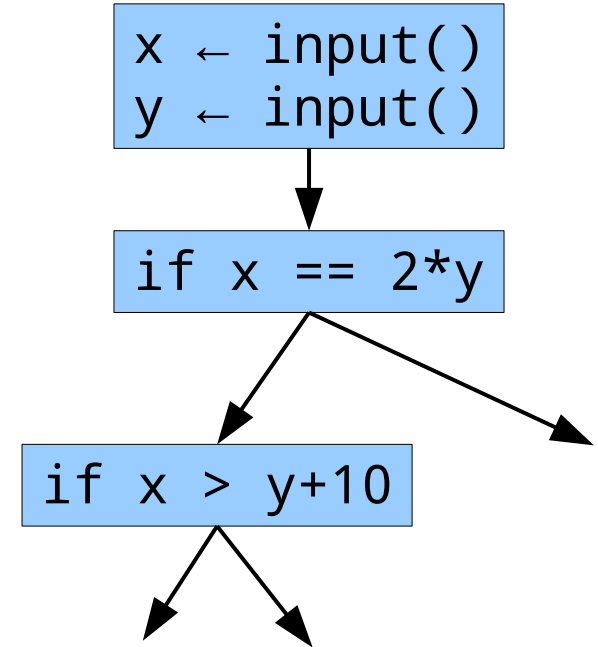
Symbolic Execution

- As we have seen, building constraints that model code can be useful
- With care, we can even use constraints to generate all inputs that are “interesting”
- Techniques for supporting this are known as ***symbolic execution***
 - (SymEx)

Symbolic Execution

- An approach for generating test inputs.

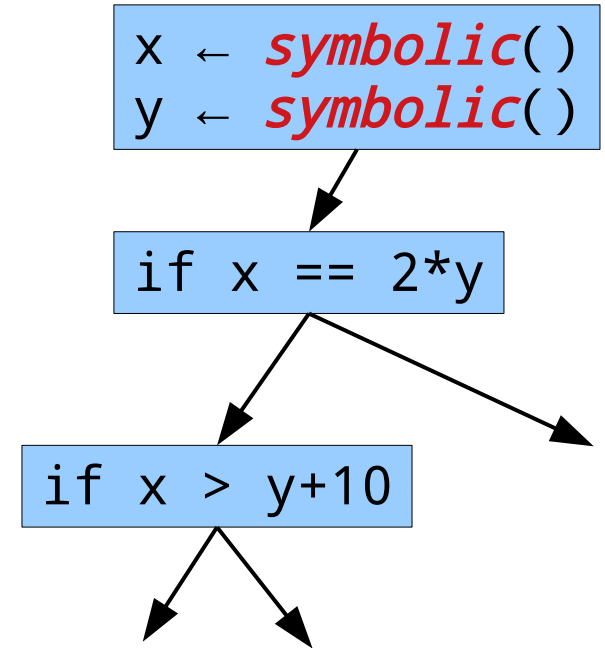
[Cadar & Sen, 2013]



Symbolic Execution

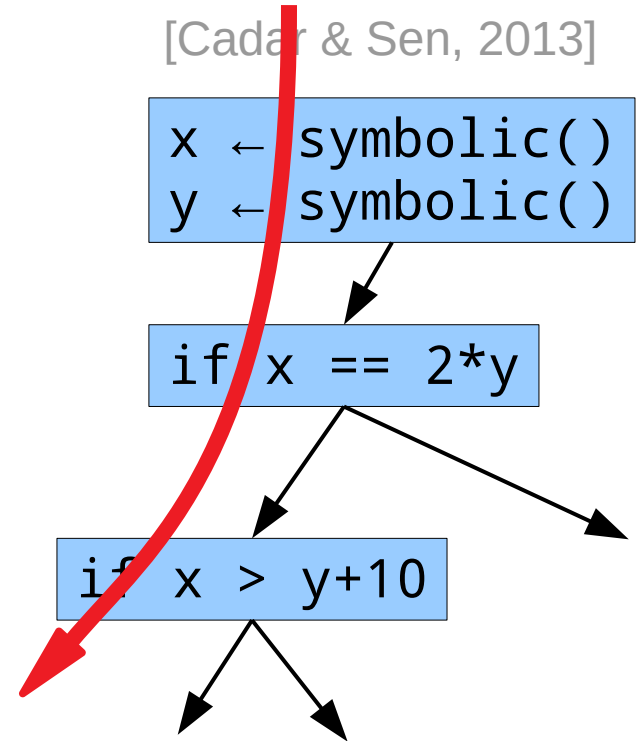
- An approach for generating test inputs.
- Replace the concrete inputs of a program with symbolic values

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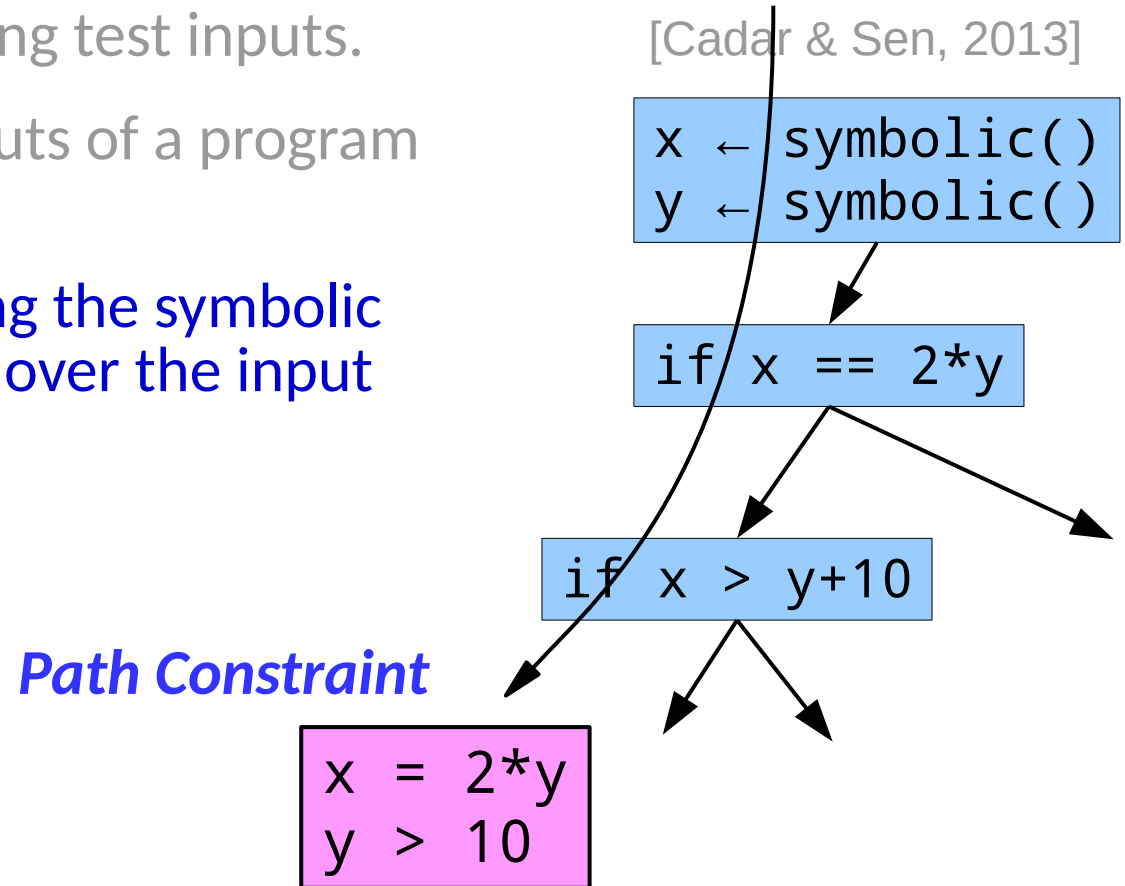
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A path constraint represents all executions along that path

Path Constraint

```
x = 2*y
y > 10
```

[Cadar & Sen, 2013]

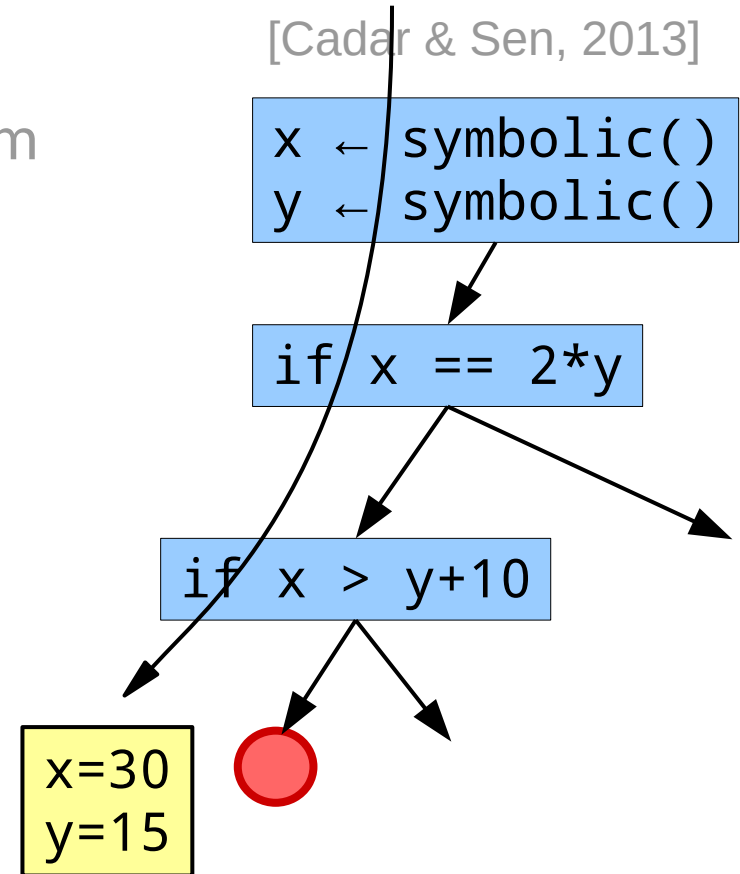
```
x ← symbolic()
y ← symbolic()
```

```
if x == 2*y
```

```
if x > y+10
```

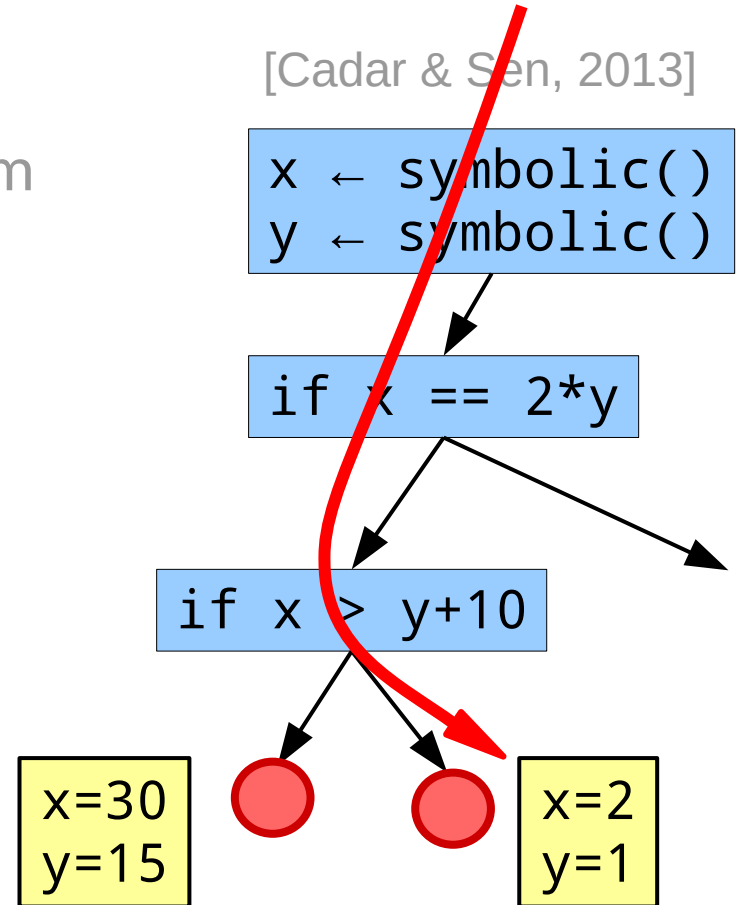
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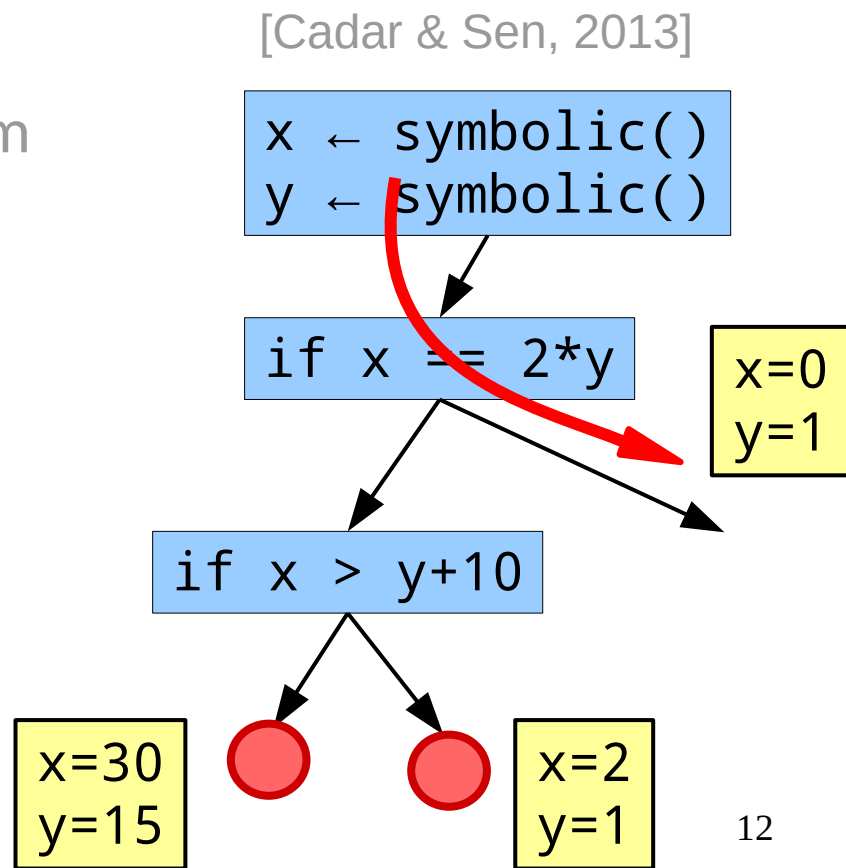
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What is $wp(\varphi_1) \wedge \neg wp(\varphi_2)$?

How Can We Solve Constraints?

- SMT Solvers
 - Satisfiability Modulo Theories
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(assert (= x (* 2 y)))
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(check-sat)
(get-model)
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$x=22$
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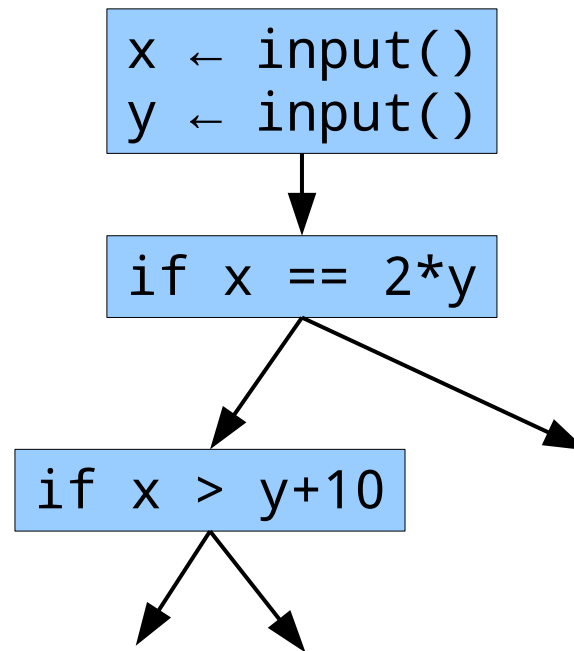
Try it online:

<http://www.rise4fun.com/Z3/tutorial/>

Exploring the Execution Tree

- The possible paths of a program form an *execution tree*.

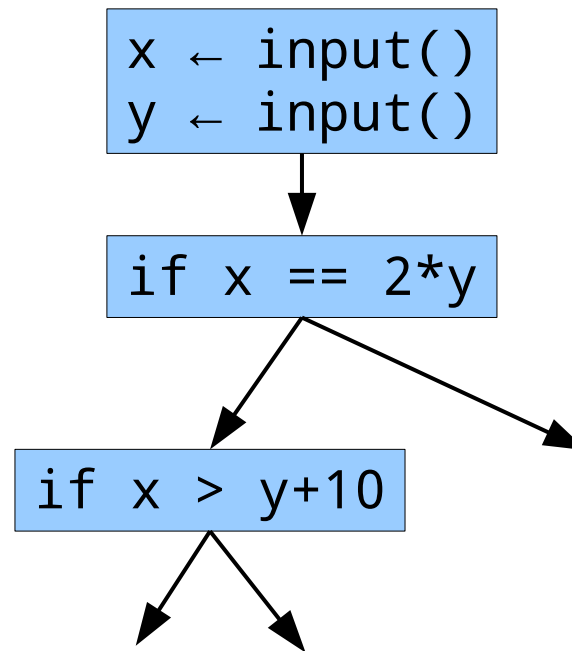
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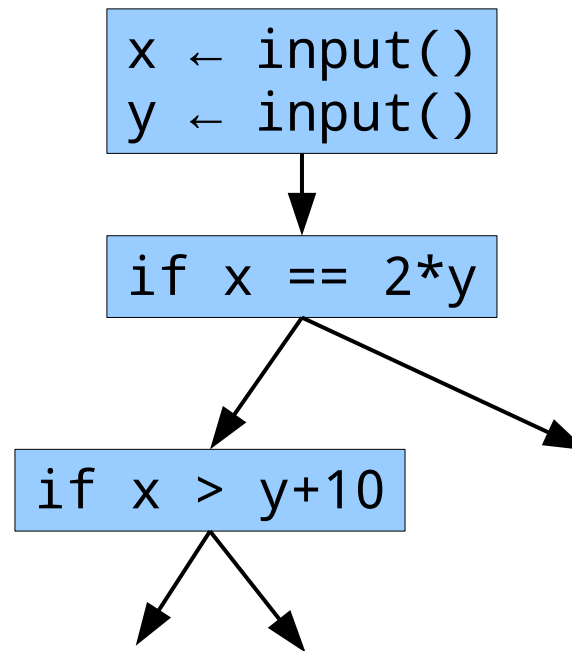
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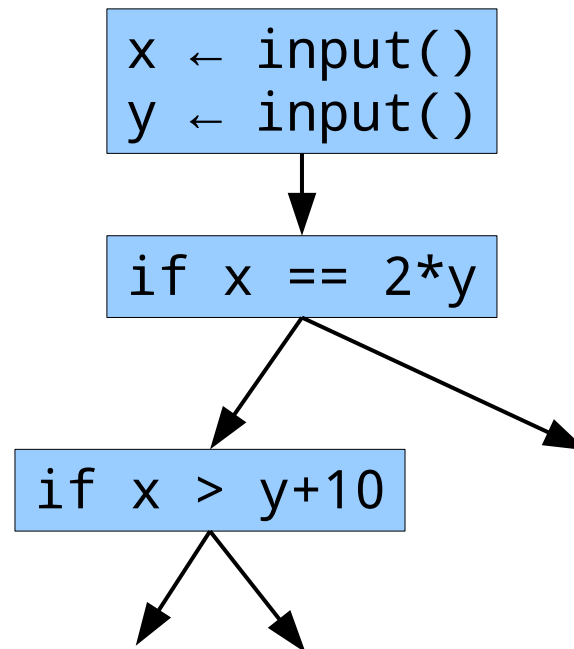
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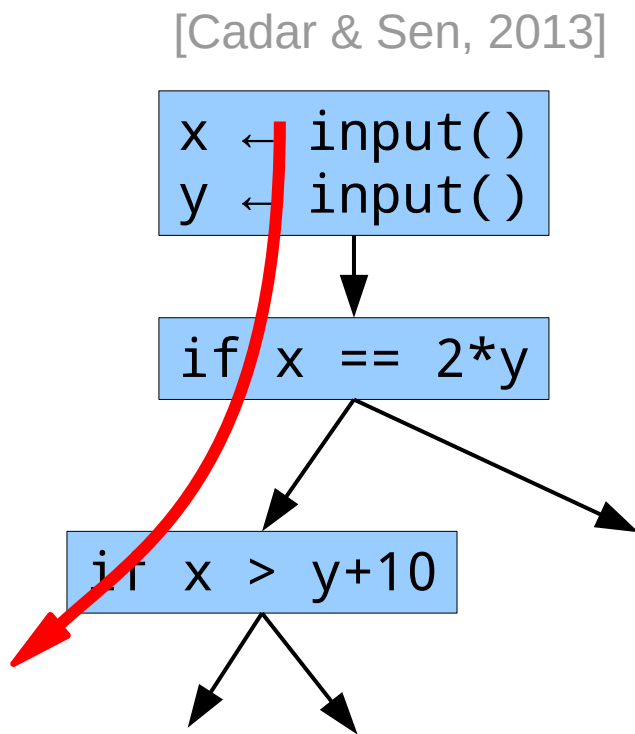
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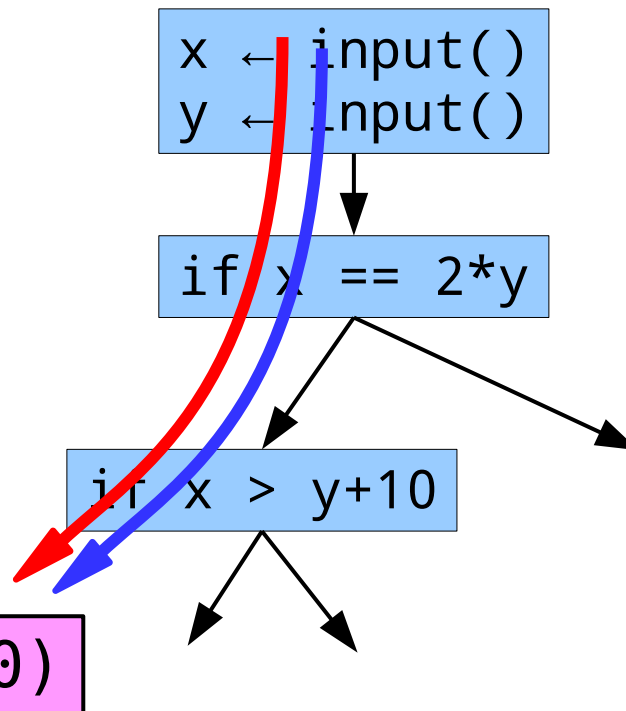
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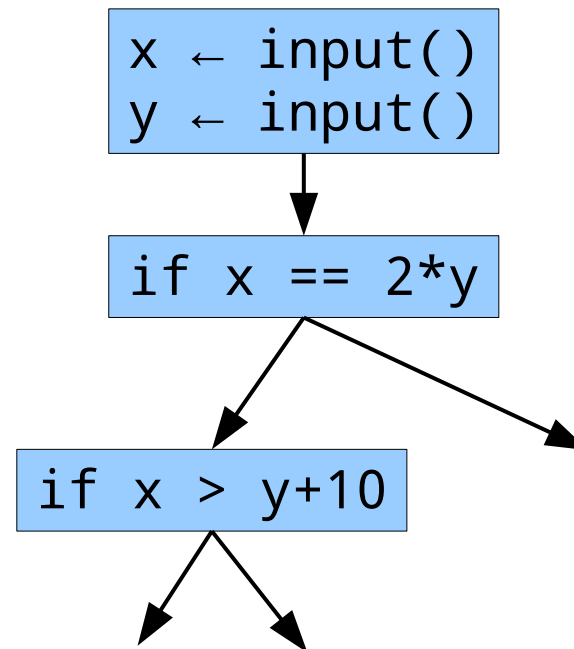
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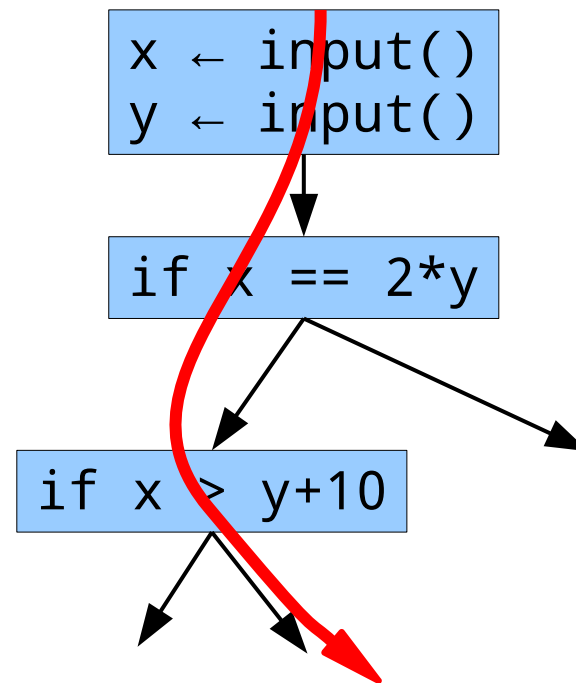


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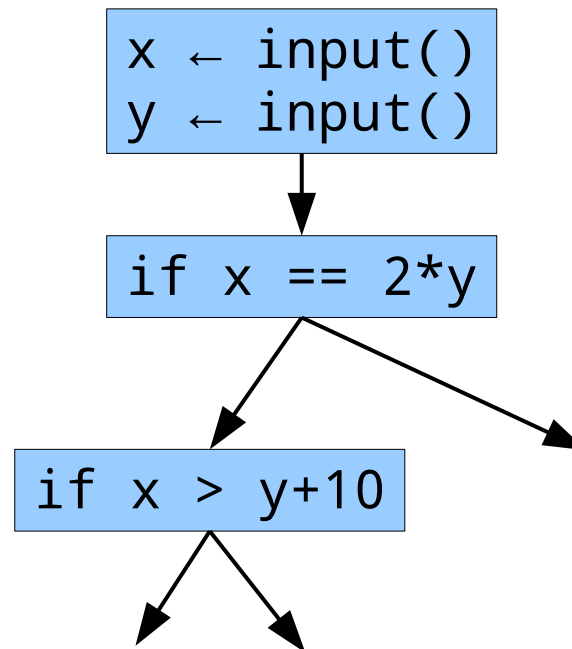


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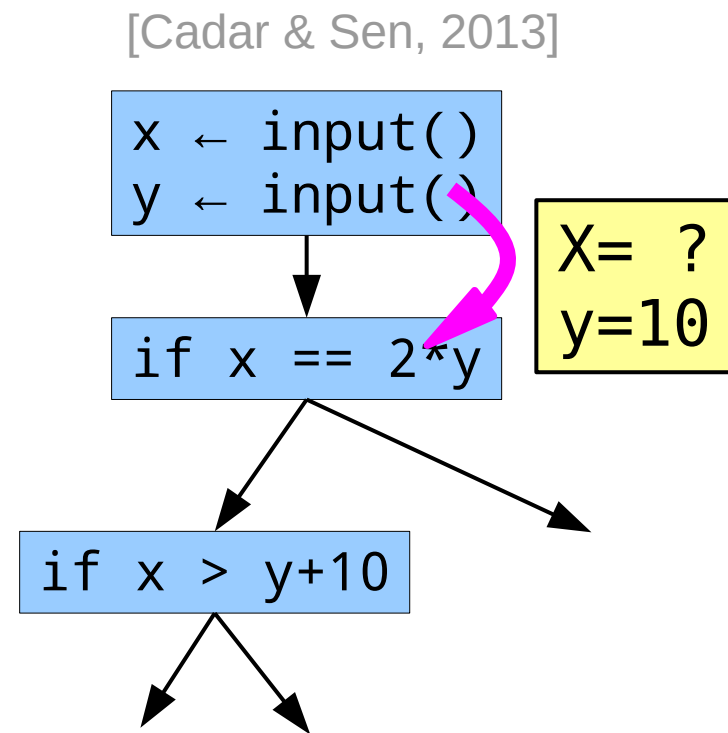
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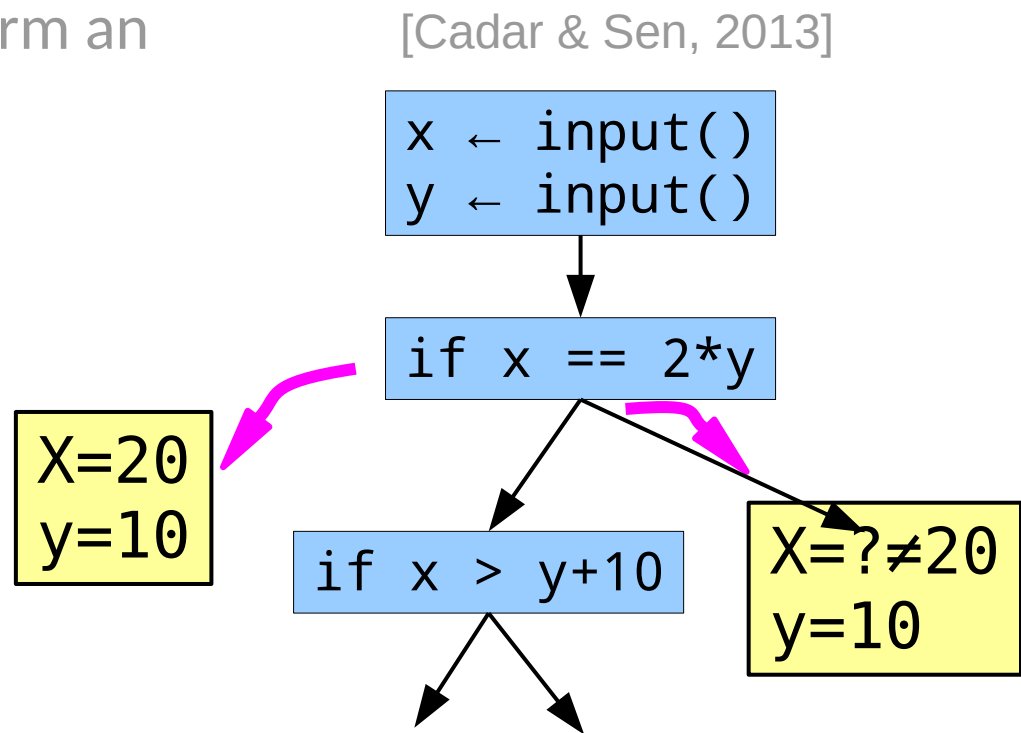
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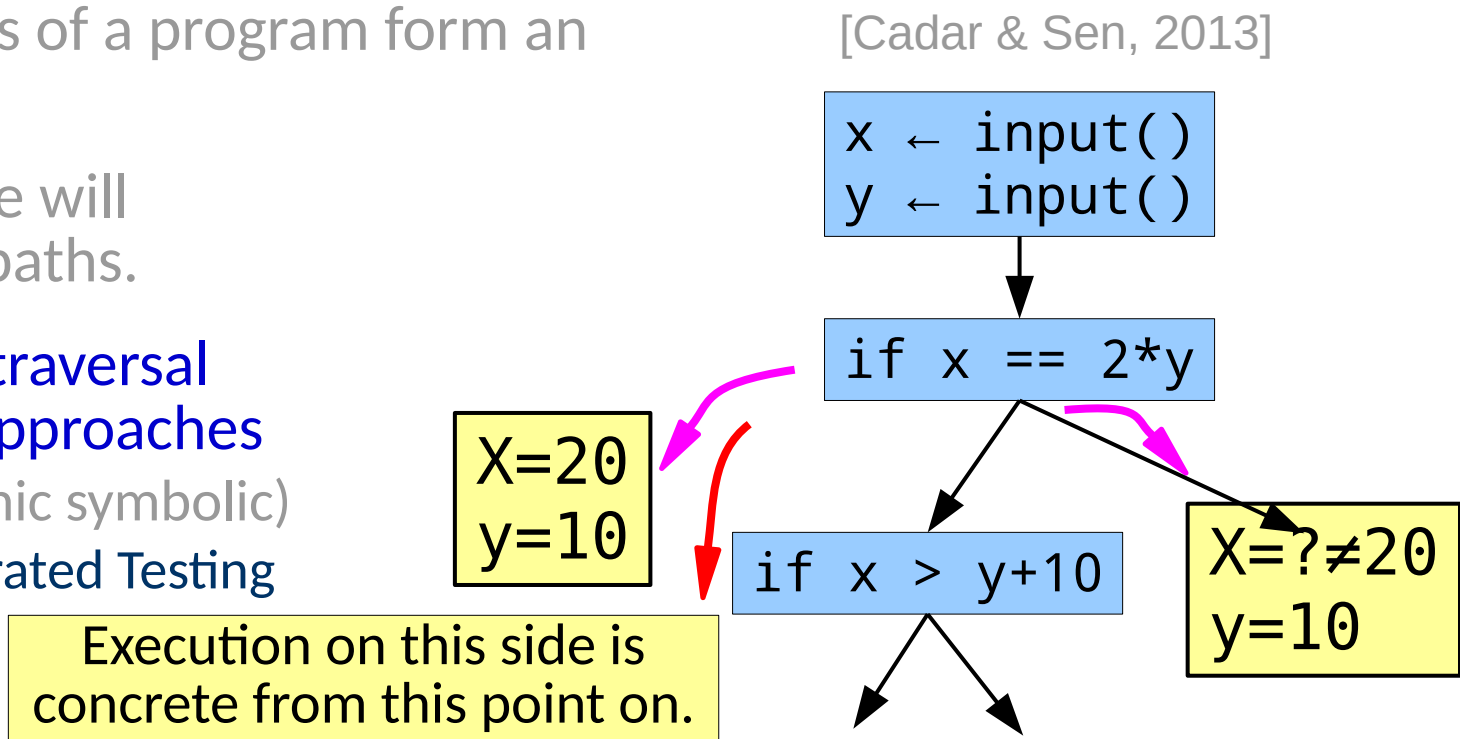
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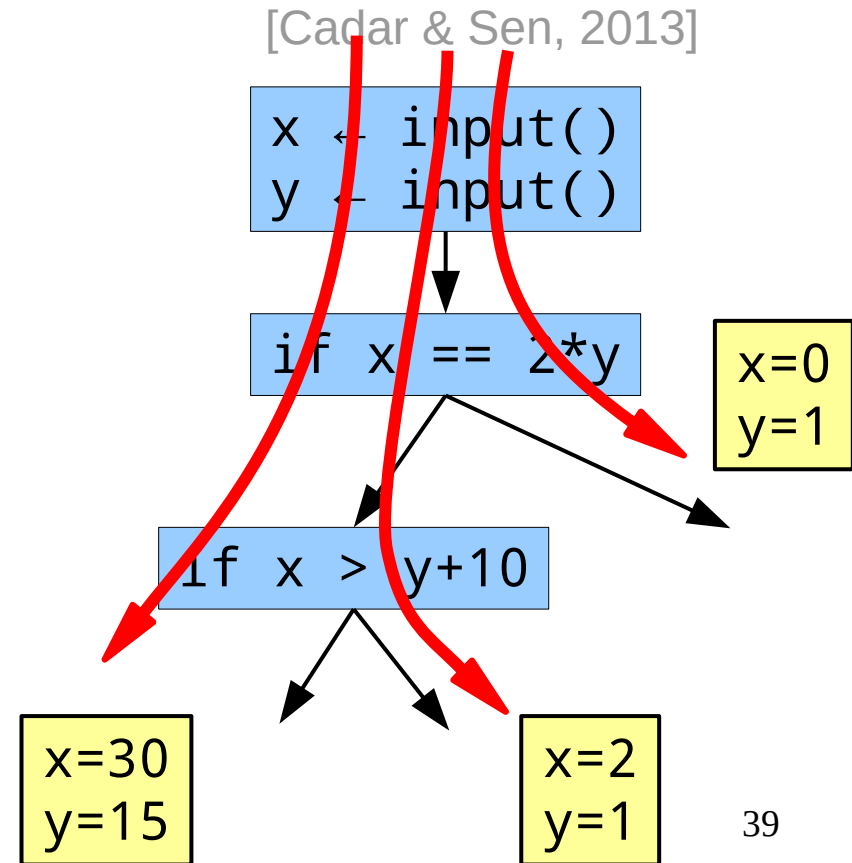


(Some) Applications

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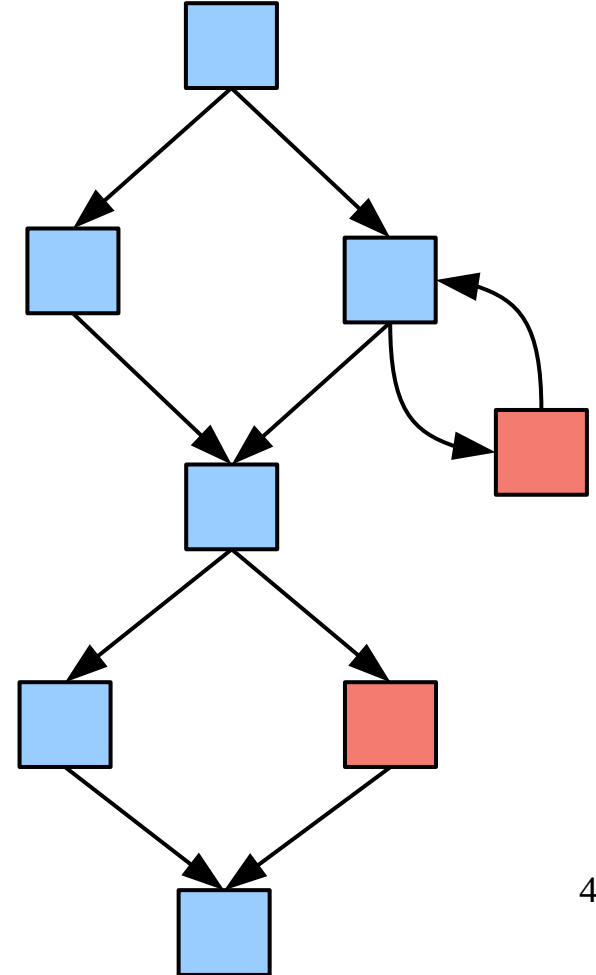


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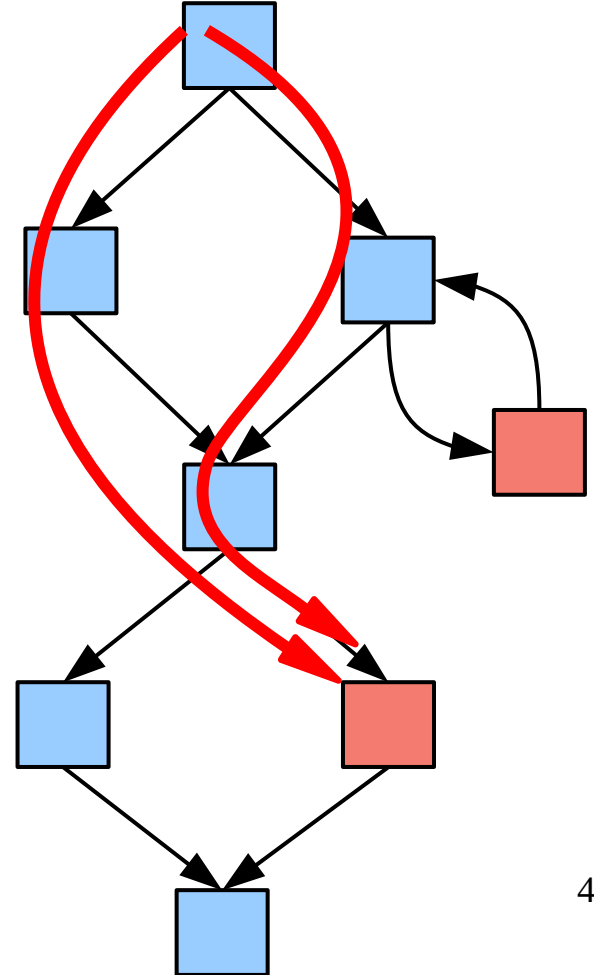
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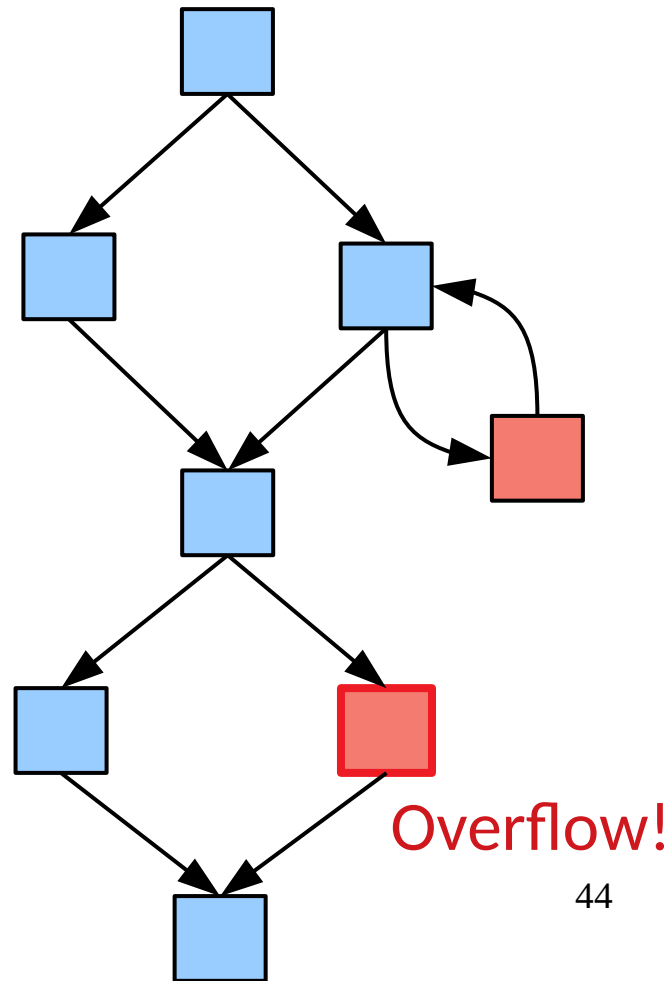


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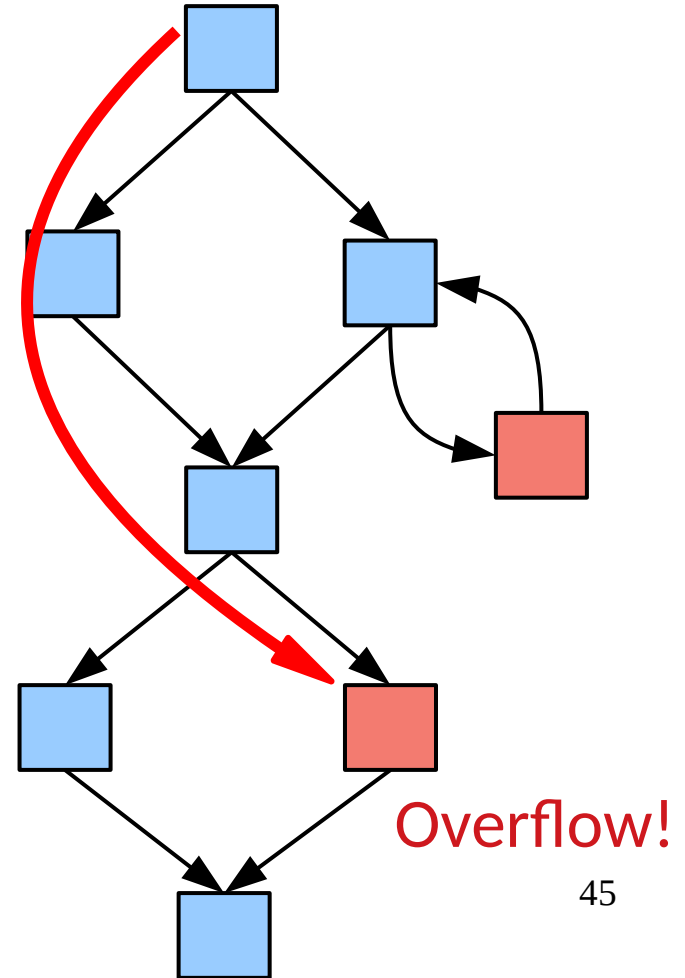
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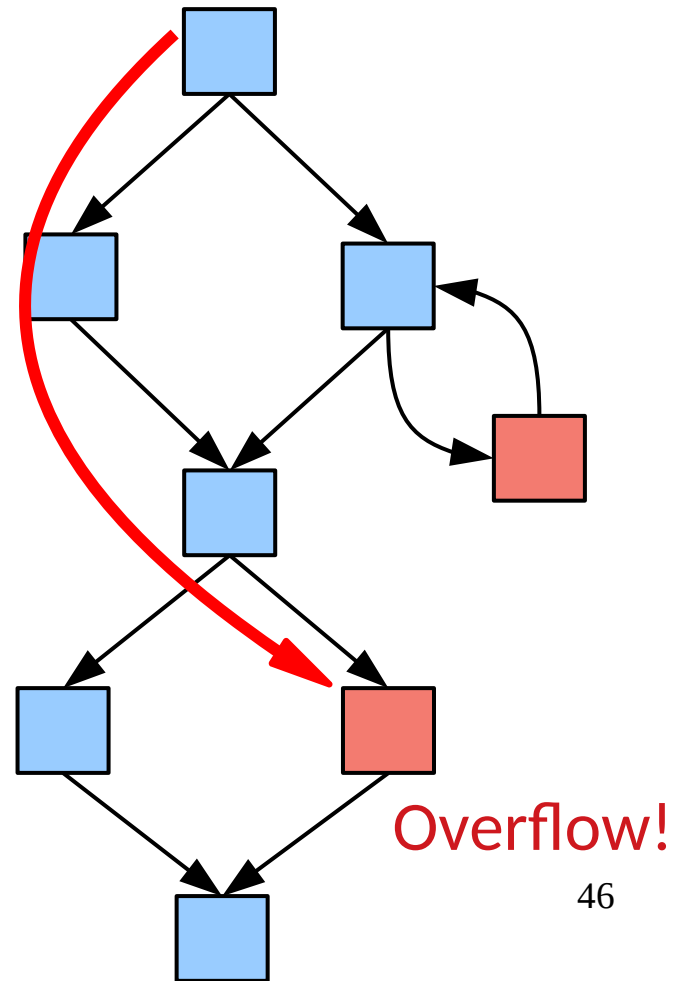


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Input \vdash Overflow \wedge StartsShellcode

This is the core process for
Darpa Cybersecurity Grand Challenge entries!



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- Test driven model checking (Yogi)
- ...

The latest testing & verification services from MS are built around these techniques.

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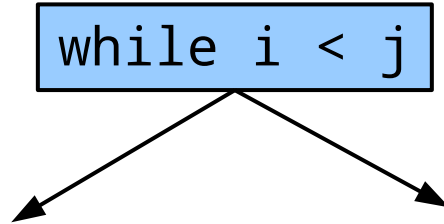
Let's revisit a familiar example...

Challenges

- Path Explosion
- Challenging constraints
- Constraint representations & domain knowledge

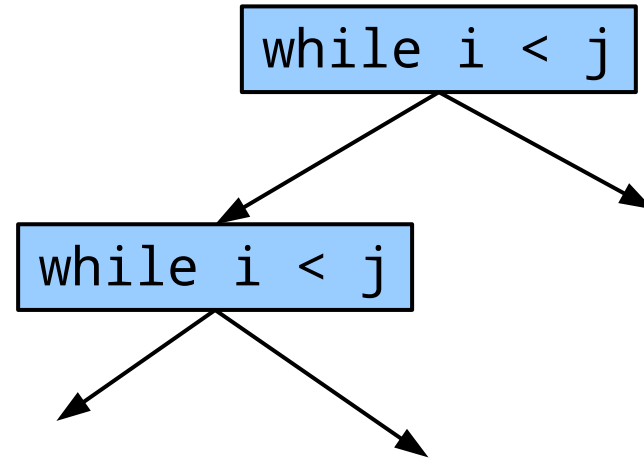
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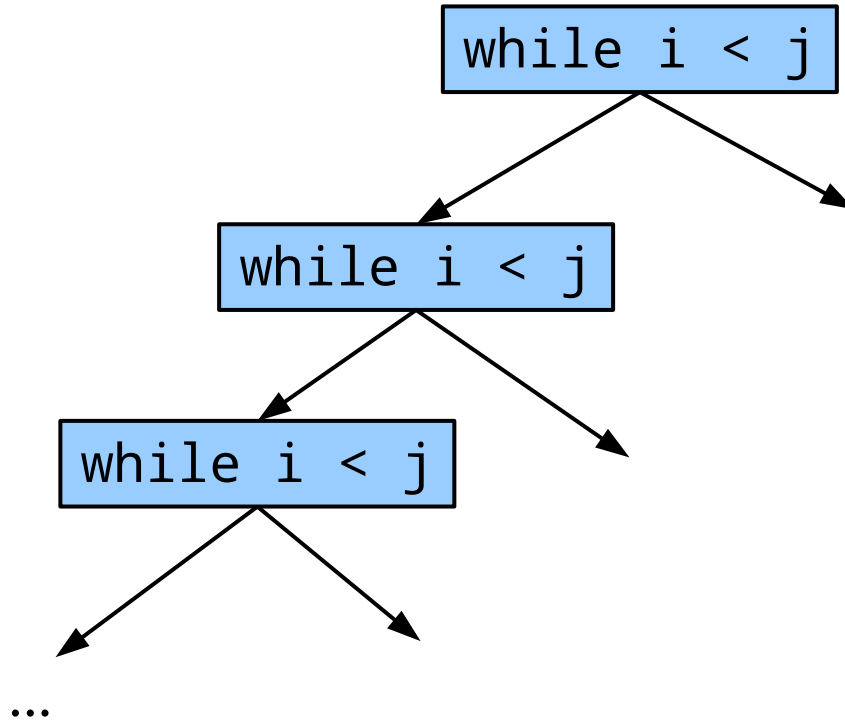
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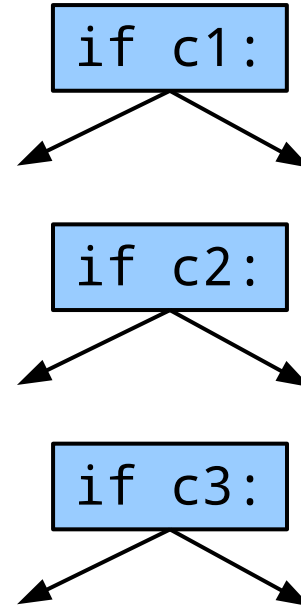
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Path Explosion

- Loops
- **Combinatorial Explosion**



Path Explosion

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State of the art techniques carefully use summarization & representations to minimize these

Path Explosion

- Loops
- Combinatorial Explosion
- Strategies
 - Search heuristics
 - Memoization

Challenging Constraints

- Intuitively, we cannot solve all constraints

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if hash(password) == y:  
    print("how odd")
```

What would it imply if we could?

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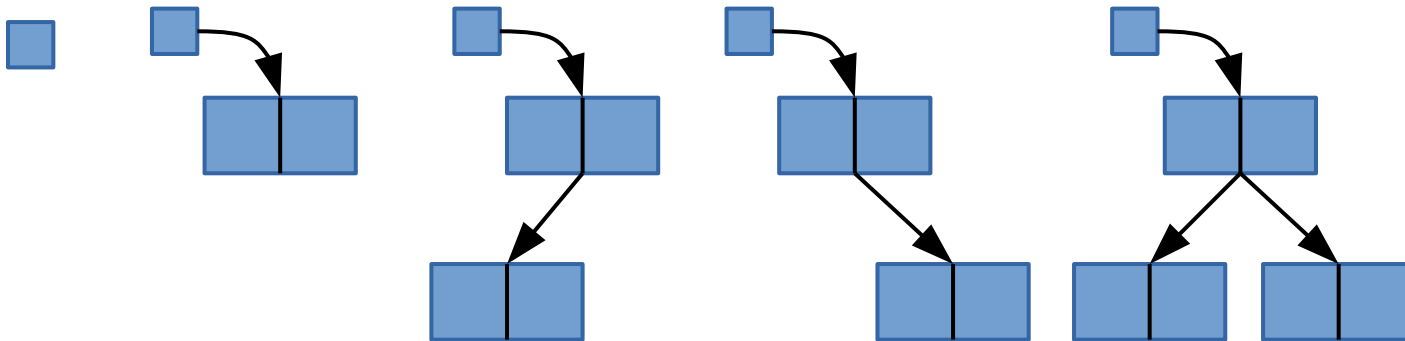
How do these affect our ability to explore the execution tree?

Domain Knowledge

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 - A linear arrangement of memory?
 - Combinatorial aliasing relation pairs?

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- Can we carefully explore interesting structures?
 - Korat style enumeration



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Try it out:

- 1) <https://github.com/klee/klee>
- 2) Symbolic PathFinder
- 3) <http://research.microsoft.com/Pex/>
- 4) <http://angr.io/>