

CMPT 473
Software Quality Assurance

Program Analysis Tools

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Fixing bugs is costly

Why?

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- Once developers have moved on, finding the root cause of a bug is difficult
- Bugs that escape into the wild have real world impact
 - Unintended car acceleration
 - Spacecraft crashes
 - Security leaks
 - ...

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Why do we still have bugs?

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 - Even prove that certain bugs are not present
 - Identify bad styles that may lead to bugs

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 - Set of tools/techniques that allow computers to automatically reason about the behavior of programs
- Push the burden of understanding programs onto computers
 - People have trouble with repetitive, subtle behavior
 - Computers excel at it

For example

```
if ((err = update(&ctx, &server)) != 0)
    goto fail;
if ((err = update(&ctx, &params)) != 0)
    goto fail;
    goto fail;
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Why should a computer be able to find it?

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 - *Rules* can determine what is buggy or not

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BUG: Both branches of the if statement have the same target

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- *Static analysis* tools
 - Examine the source code or binary and reason about **all possible executions**
 - Best at identifying bugs that haven't struck yet but might in the future

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This one is tougher....

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The halting problem strikes *again*....

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 - Static approaches are limited by the halting problem
- The results are imperfect
 - False positives – Warnings about bugs that don't actually exist
 - False negatives – Missing warnings for bugs that do exist
- Learning how to use these tools effectively can take practice

But what can they actually do?

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Read for Wednesday

But what can they actually do?

- You've already seen the PVS-Studio examples
- Many tools are freely available:
 - *Lint
 - FindBugs
 - Clang Static Analyzer
 - ESC/Java
 - Valgrind
 - Clang Sanitizers
 - ... (and more on the course web page)

Taking a look at Valgrind

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Does not work for Java or Python by default. Why?!

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 - Used extensively at google (chrome, ...)

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- What about the static analysis tools?

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 - And a plug-in system for recognizing new ones.

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 - Uses *abstract interpretation* to simulate many different paths through the program at once
 - Generates summaries showing exactly how errors *may* occur
 - Many automatically recognized bugs
 - And a plug-in system for recognizing new ones.
 - Poorly organized & asserted code yields many errors

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I would argue that a Java project not using FindBugs is a broken project!

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You must eventually figure out
that the ghost isn't real

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 - Want to determine whether warnings are real

This takes a lot of work & happens every time.
Can we do better?

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Blacklisting & suppression allows us to “remember” false positives & prevent them in the future....

[DEMO]

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Have you seen / heard of such tools before?

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Any ideas?

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But they are getting better!
Used extensively in safety critical systems.