Declarative Memory Services

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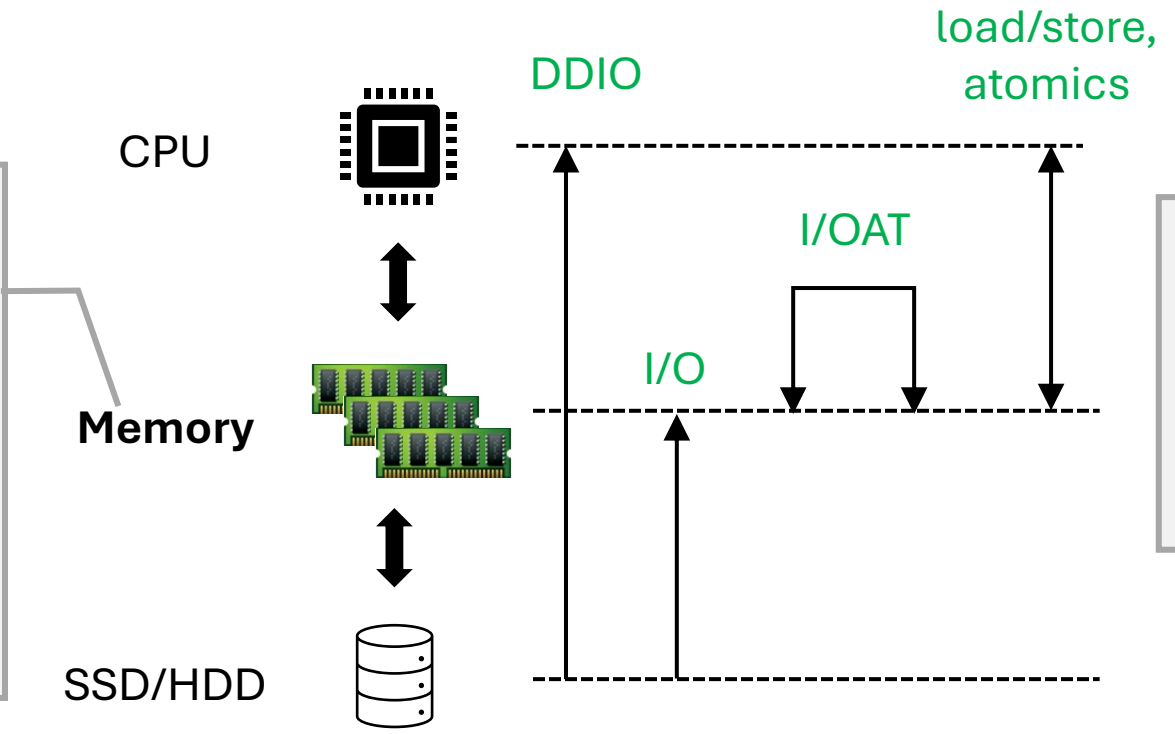
“Memory” traditionally…

Properties:
- Single-node
- Byte-addressable
- Volatile DRAM
- Coherent
- ~100ns low latency
- High bandwidth
- Passive

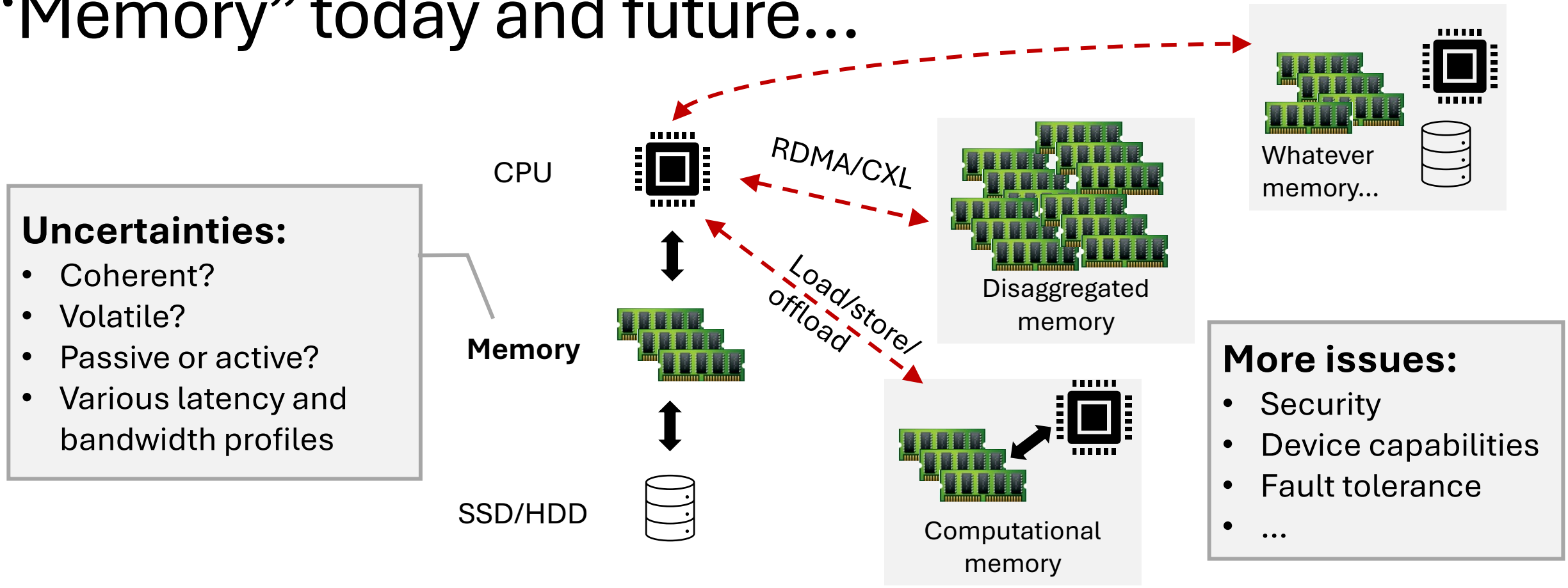


Issues:
- NUMA-awareness
- Allocator performance
- Cache-conscious
- …

Relative tractable primitives and tools + imperative programming

Life was ok.

“Memory” today and future…



Uncertainties:

- Coherent?
- Volatile?
- Passive or active?
- Various latency and bandwidth profiles

More issues:

- Security
- Device capabilities
- Fault tolerance
- …

Intractable primitives → highly complex, imperative programming

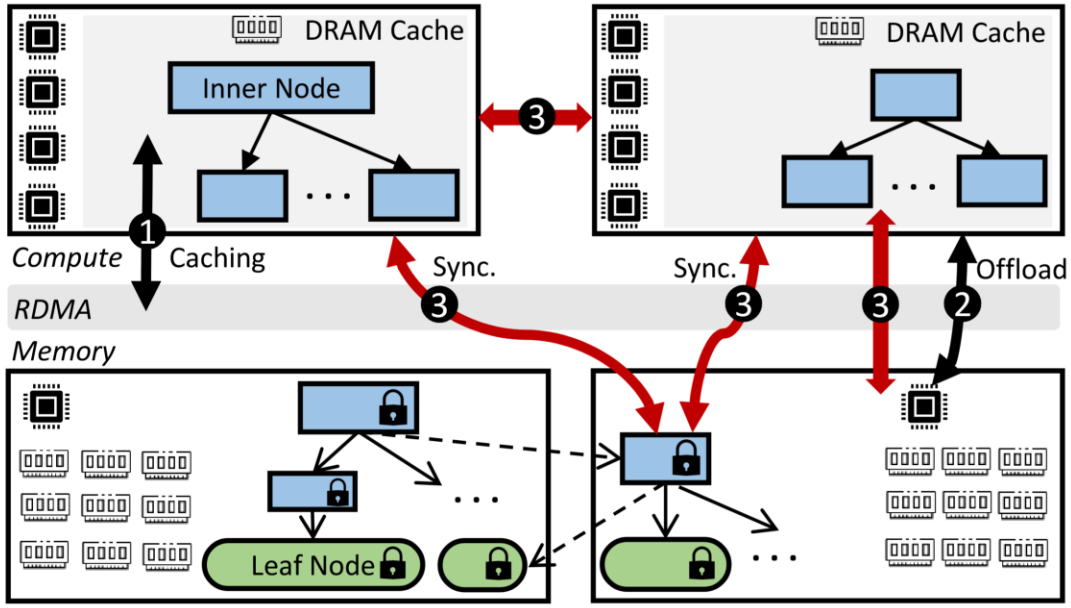
Life is hard.

Case Study:

Adapting a B+-Tree for disaggregated memory

(1) Longer latency, should cache:

- Which B+-tree nodes to cache?
- Is there coherence between compute servers?



(2) Memory has CPU, should offload:

- How much CPU do I have?
- What operations to offload?

(3) Data placement + replication:

- Who can access which data?
- How to partition?

* DEX: Scalable Range Indexing on Disaggregated Memory, VLDB 2024

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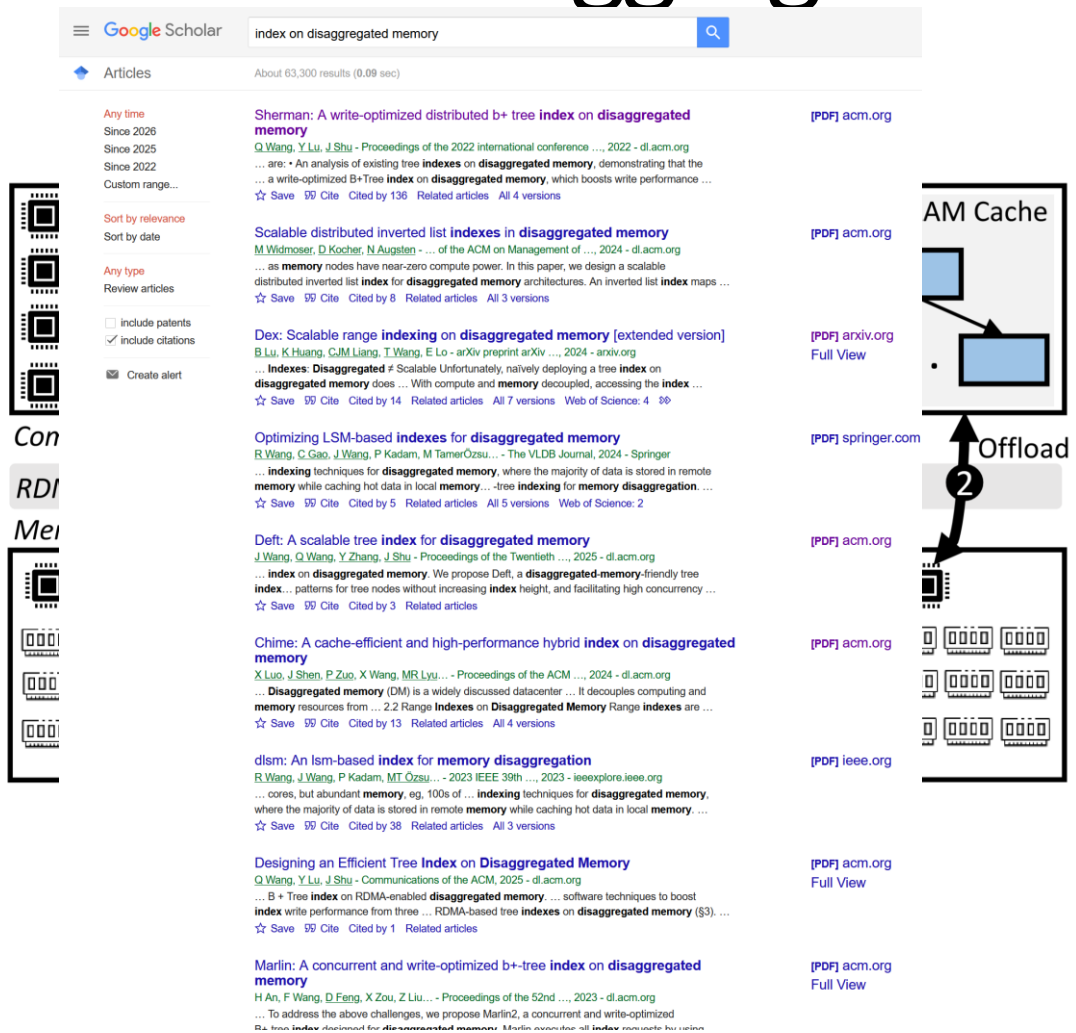

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Hand-coded decisions
Unsustainable (*more cases in paper*).

Would be nice to be more *declarative*

- Decouple device-specific logic from high-level design
  - “I want this function to be offloaded, if possible”
  - “Latency to access this memory block should not exceed 5ms”
- Simplify programming for today and future, unknown architectures
  - Same DBMS design, any hardware
- Better cross-device optimizations

How to get there?

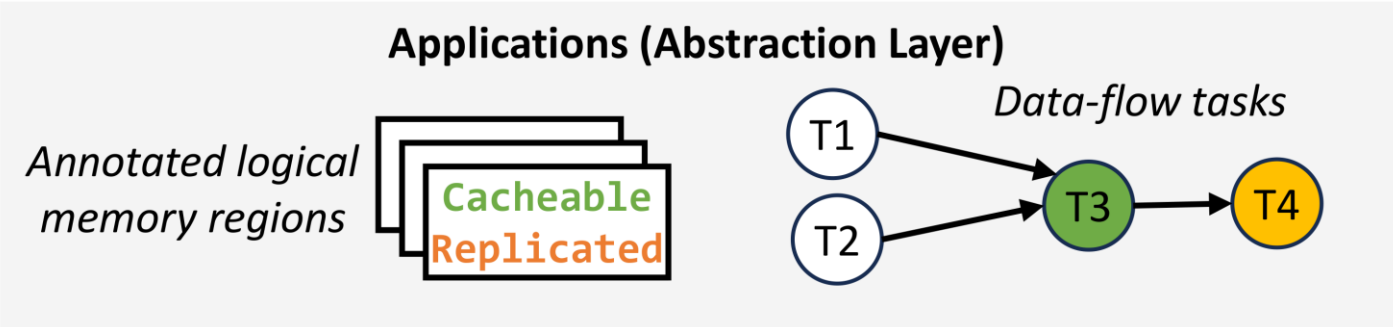
Vision: Declarative Memory Services

Three-layer design:

- Abstraction Layer
  - Developers work with “logical memory regions” and data flows
  - Annotate with desired properties
- Calibration Layer
  - Discover and index device capabilities
  - Expose device primitives and APIs
- Memory Services Layer
  - A set of generic “memory services” that well use memory devices
  - Jointly optimize for the application based on annotations

Caveat: yet to implement, this is pure vision!

Declarative Abstraction layer



B+-tree node definition:

```
struct InternalNode {
  KV kv_pairs[MAX_KV];
  int key_count;
  …
};
```

Declare desired properties

Previously:

```
InternalNode *n = allocate(…)

// hand-made decision to cache it
cache.insert(n);
```

Now with DMS:

```
[cacheable, coherent, latency < 10μs]
InternalNode *n = allocate(…)

// placed in coherent, compute-side memory, by DMS
~~cache.insert(n);~~
```

Data flows work similarly:

- Properties attached to tasks, enforced by DMS runtime

Physical design and logical functionality decoupled

Calibration Layer

- Discovers and track device capabilities, provide APIs
- Key component: device catalogue
  - A table that evolves with hardware changs

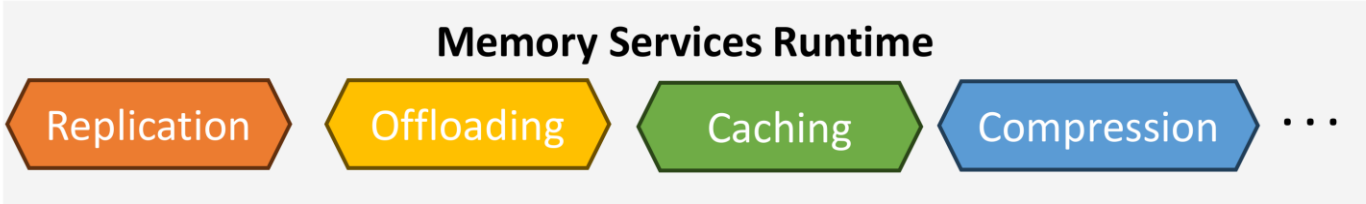
Device	Capabilities	APIs	Characteristics
Local DRAM	Coherence Byte-addressable	dram-load, dram-store, dram-dsa, atomics…	… x GBps within socket, under y load…
CXL DRAM	Partial coherence Byte-addressable	cxl-load, cxl-store…	… 300ns best - 1us worst latency…
Membrane (computational memory)	Compute Byte-addressable	pim-load, pim-store, pim-offload…	… x ns latency with host…

Implemented and maintained by DMS developers

Challenging

Memory Services Layer

- Use device catalogue APIs to build services



- DEX example:
  - Services needed: data placement and caching
  - Upon allocation: place data based on annotated desired properties
  - Runtime: lightweight metadata tracking for caching
- Customized policies possible
  - “Please don’t evict parent node before child node”
  - “Please use this encoding scheme for such and such data”

Research Challenges and Agenda

- Device Characterization
  - Beyond simple stats: e.g., latency behaviour under varying load levels
  - Self-evolving the device catalogue with new hardware
- Properties → Services: When to pick which implementation?
- SLA Guarantees
  - Memory services monitor metrics, and migrate between services to meet SLO
  - How to deal with conflicting SLAs?
    - E.g., tenants prioritizing throughput vs. latency
- DMS Deployment
  - DMS requires non-trivial information (global and local server) to work
- Correctness and Debugging
  - DMS-based programs are declarative
  - How to verify their correctness and debug them? Tools for exploring why an SLO was missed?

Summary

- Memory is heterogeneous: complexity arises with more features
  - Current approach to leveraging memory devices is unsustainable
  - Hand-crafted with low-level primitives
  - Getting worse as hardware evolves

- **Declarative Memory Services**
  - Developers specify logical functionality
  - Calibration layer discovers and characterises devices
  - Memory services provide physical implementations and optimizations

Thank you!