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ISP-Friendly Peer Matching without ISP Collaboration

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Introduction

- P2P Systems (file sharing, live streaming, VoD, ...) are very popular nowadays
 - Generating a major fraction of Internet traffic
- In most of these systems, peers offer limited capacity and reliability ->
- Client peer typically served by multiple senders
- Choosing "good" senders is important for
 - Clients: better performance
 - P2P System: more capacity, better scalability
 - Network (ISPs): reduced load on inter-/intra-ISP links



Our Work: Peer Matching Problem

- Given set of potential senders for a receiver, find subset of them that minimize the load on ISPs
 - Load = traffic on inter-ISP links (expensive) and on intra-ISP links
- General formulation/solution
 can be used in different Systems
 - BitTorrent-like: use in tracker
 - Gnutella-like: use in peers
 - Peer-assisted CDN: use in managing servers



Our Approach

- Find nearby peers in terms of
 - AS hops → reduce inter-ISP load
 - PoP (Point of Presence) distance → reduce intra-ISP load
- Efficiently estimate AS paths &PoP distances
 - Small memory/CPU requirements
 - Use public data (BGP tables, GeoIP, ...), no probing
- No infrastructure need to be deployed
- No need for cooperation from ISPs
- No need to modify client software



Related Work

- Provider Portal for P2P (P4P) [Xie 08]
 - ISPs deploy servers (iTrackers) to guide peer matching
 - Require cooperation from ISPs and infrastructure
- Using DNS redirections [Choffnes 08]
 - Peers observe similar delay to CDN replicas → close to each other
 - Require modifying clients
- Match peers within the same AS [Bindal 06]
 - Specific to BitTorrent
 - Peers outside AS are chosen randomly



Overview



- Large ISP: PoPs at many locations
- S3 is better for R than S2, though same AS distance



Overview

- We need to estimate AS hops between peers
 - Need AS path inference algorithm (valley-free policy)
 - Modified Dijkstra Algorithm [Mao 05] → very expensive, took ~2 days of running time, huge memory footprint
 - Our optimization: runs 1-3 hours, < 1% memory
- We also need to infer PoP topology
 - Propose simple method to cluster IP prefixes
 - Use GeoIP databases to map clusters to PoPs



Peer Matching



 Distance oracle: concise data structure, enables online matching in O(1) steps



AS Path Inference

- Original algorithm [Mao 05]
 - O(V³), V number of ASes
 - Currently, there are **28,000**+ ASes [CAIDA]
 - Very expensive
- Observation
 - Many ASes have small degree (client/stub ASes)
 - 83% of ASes have 1 or very few C2P links [CAIDA]
 - → their AS paths can be inferred using their providers' paths
 - → stub ASes can be removed from the AS graph



AS Path Inference





AS Path Inference

- Estimate AS paths for "core" ASes first
 - Using any algorithm
- Then paths starting/end at a transit AS
- Then, the rest
- Time Complexity: O(T³ + T S + S²)
 - T small fraction (<18%) of ASes → much faster
 - 3 hours (instead of 2 days) for <u>whole matrix</u>





Distance Oracle

- Maintaining the whole distance matrix is costly
 - Assume each entry takes 1 byte
 - #of ASes = 28,594 → 817 MB for distance matrix
- Idea: Maintain only the "core" matrix, other entries can be computed online
- Tradeoff: small core vs. more online computations



Distance Oracle: Tradeoff

Level	Core Size (#ASes)	Memory (MB)	Online Computations (Average)
0	28,594	817	0.0
1	5,144	26	2.6
2	2,949	8	12.4
3	2,471	6	97.0

- We can use a distance oracle of size 8 MB, and perform < 13 operations online on average</p>
 - #operations depends on AS degree, not size of AS graph



PoP Distance: Overview

- Infer approximatePoP topology of ASes
- Algorithm overview:
 - Collect all IP prefixes from BGP tables and map them to their ASes
 - For each AS, cluster IP prefixes into one or more PoPs
 - Analyze BGP updates to infer connectivity among PoP in same AS
 - Combine PoP topology with AS graph



Evaluation

- Trace-based simulations
- Use real IPs from
 - BitTorrent tracker: ~150,000
 - Online content provider (CBC -- Canadian Broadcasting Corporation); ~ 160,000
- BGP tables from RouteViews& RIPE
- AS relationship data from CAIDA
- Use GeoIP [maxmind.com] for IP/AS/location mapping



Evaluation

- Simulate multi-sender sessions for objects with different popularities: 1% -- 10%
- Change number of chosen senders: 10 100
- Implement
 - **ISPF:** ISP-Friendly Matching
 - **ISPF-Lite:** ISP-Friendly Matching, without PoP info
 - AS: Match based on AS distance only
 - **Prefix:** Match based on longest common IP prefix
 - Uses IPs of peers → almost free
 - **Random:** currently used in many systems





- Some gain from (the free) Prefix
- Significant gain from ISPF/ISPF-Lite





Even more gain for rare objects





ISPF achieves shorter distances



Dynamics of AS Graph



- CAIDA AS relationship data for a full year
- run AS inference algorithm once a month or so

Conclusions and Future Work

- Peer matching is important in multi-sender systems
 - P2P systems, distributed streaming, CDN, pCDN, ...
- Peer matching algorithms based on
 - AS distance
 - PoP distance
- Do not need infrastructure, modifying peers, ISP cooperation
- Efficient computation and storage of AS distance
 - ~ 8 MB (<1% of whole matrix) for 28,000+ ASes



Conclusions and Future Work

- Currently implementing in pCDN (for CBC) and in BitTorrent tracker
- Exploring link cost model
 - Assign weights to inter-AS links
- Exploring more efficient methods for discovering & storing PoP topology



Thank You!

Questions??

More info at:

http://nsl.cs.sfu.ca/

