

Doubly Connect Edge List (DCEL)

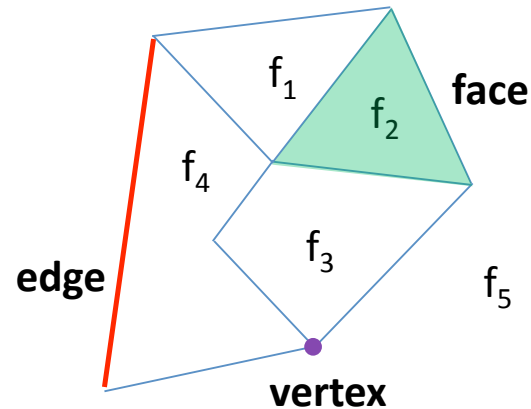
Notes from the book by de Berg, Van
Krevald, Overmars, and Schwarzkopf.

pp. 29-39

Doubly Connected Edge List (DCEL)

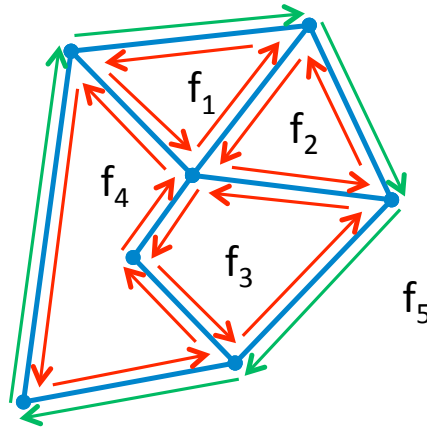
- DCEL is one of the most commonly used representations for planar subdivisions such as Voronoi diagrams.
- It is an edge-based structure which links together the three sets of records:
 - **Vertex**
 - **Edge**
 - **Face**
- It facilitates traversing the faces of planar subdivision, visiting all the edges around a given vertex

Doubly Connected Edge List (DCEL)



- Record for each face, edge, and vertex
 - Geometric information
 - Topological information
 - Attribute information
- Half-edge structure

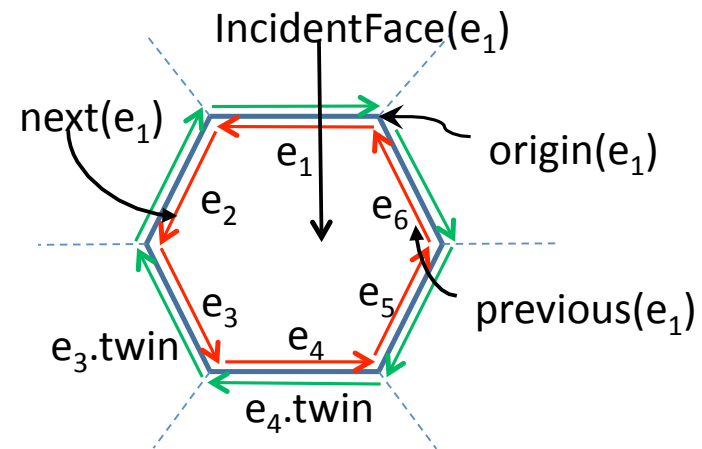
Doubly Connected Edge List (DCEL)



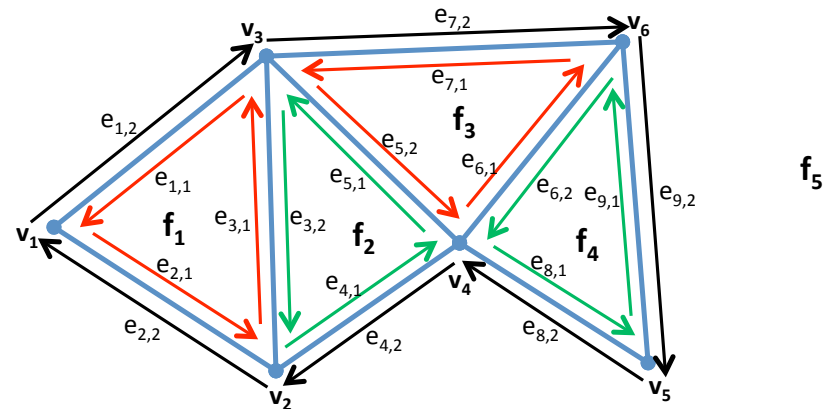
- **Main ideas:**
 - Edges are oriented counterclockwise inside each face
 - Since an edge borders two faces, each edge is replaced by two half-edges, one for each face

Doubly Connected Edge List (DCEL)

- The vertex record of a vertex v stores the coordinates of v . It also stores a pointer $\text{IncidentEdge}(v)$ to an arbitrary half-edge that has v as its origin
- The face record of a face f stores a pointer to some half-edge on its boundary which can be used as a starting point to traverse f in counterclockwise order
- The half-edge record of a half-edge e stores pointer to:
 - **Origin** (e)
 - Twin of e , $e.\text{twin}$ or $\text{twin}(e)$
 - The face to its left ($\text{IncidentFace}(e)$)
 - **Next**(e) : next half-edge on the boundary of $\text{IncidentFace}(e)$
 - **Previous**(e) : previous half-edge

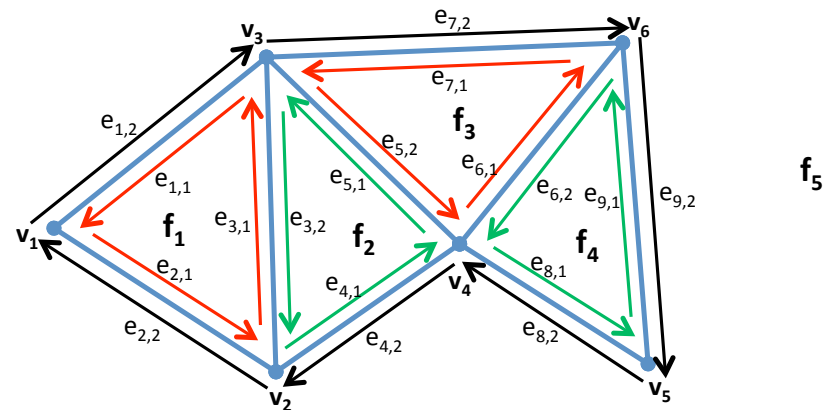


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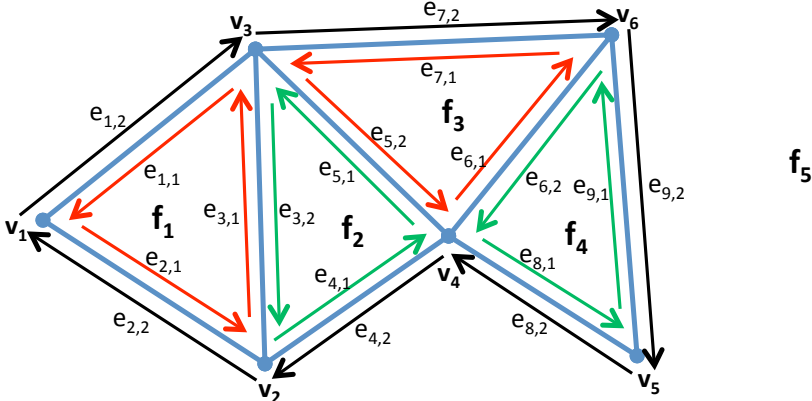
Vertex	Coordinates	IncidentEdge
v_1	(x_1, y_1)	$e_{2,1}$
v_2	(x_2, y_2)	$e_{4,1}$
v_3	(x_3, y_3)	$e_{3,2}$
v_4	(x_4, y_4)	$e_{6,1}$
v_5	(x_5, y_5)	$e_{9,1}$
v_6	(x_6, y_6)	$e_{7,1}$

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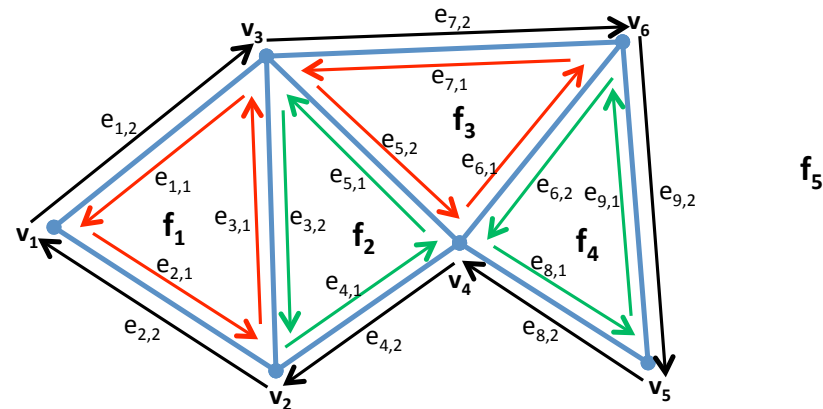
Face	Edge
f_1	$e_{1,1}$
f_2	$e_{5,1}$
f_3	$e_{5,2}$
f_4	$e_{8,1}$
f_5	$e_{9,2}$

Doubly Connected Edge List (DCEL)



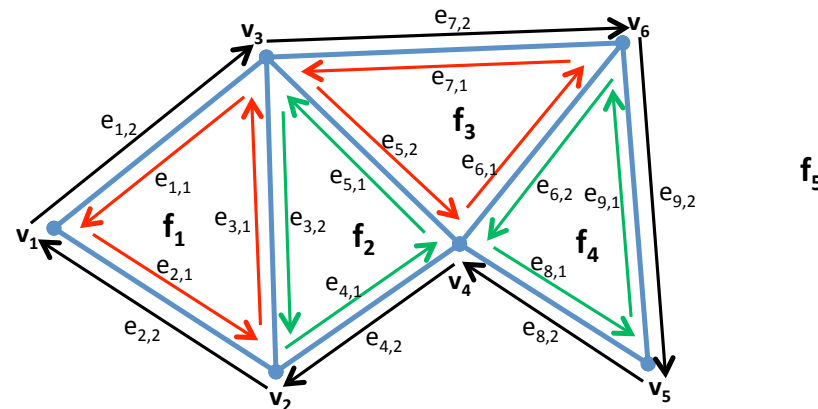
Half-edge	Origin	Twin	IncidentFace	Next	Previous
$e_{3,1}$	v_2	$e_{3,2}$	f_1	$e_{1,1}$	$e_{2,1}$
$e_{3,2}$	v_3	$e_{3,1}$	f_2	$e_{4,1}$	$e_{5,1}$
$e_{4,1}$	v_2	$e_{4,2}$	f_2	$e_{5,1}$	$e_{3,2}$
$e_{4,2}$	v_4	$e_{4,1}$	f_5	$e_{2,2}$	$e_{8,2}$
...

Doubly Connected Edge List (DCEL)



- **Storage space requirement:**
 - Linear in the number of vertices, edges, and faces

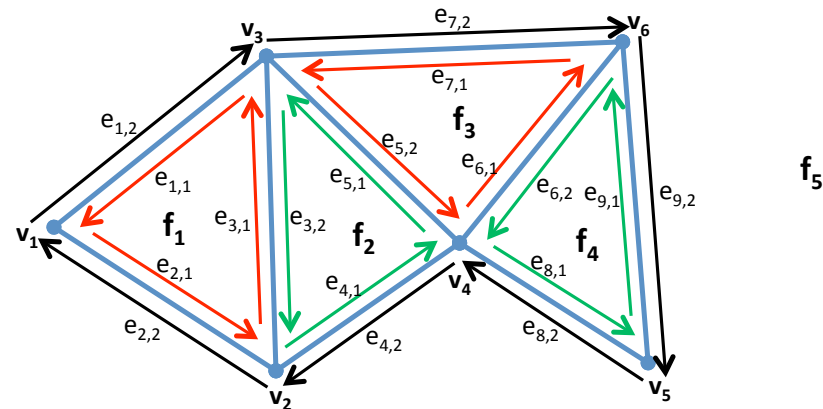
Doubly Connected Edge List (DCEL)



- **Operations:**

- Walk around the boundary of a given face in CCW order
- Access a face from an adjacent one
- Visit all the edges around a given vertex

Doubly Connected Edge List (DCEL)



- Interesting Queries:

- Given a DCEL description, a line L and a half-edge that this line cuts, efficiently find all the faces cut by L .

Doubly Connected Edge List (DCEL)

- **Traversing face f :**

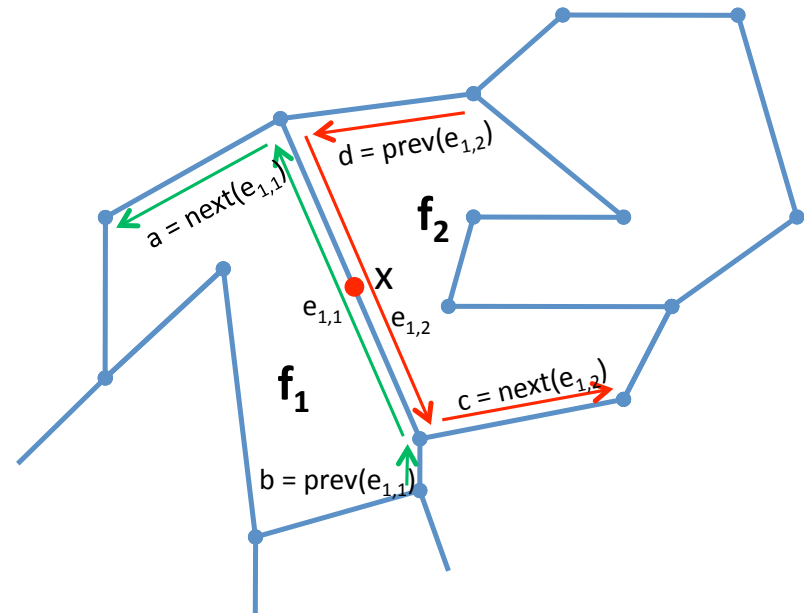
- Given: an edge of f

1. Determine the half-edge e incident on f
2. $\text{start_edge} \leftarrow e$
3. While $\text{next}(e) \neq \text{start_edge}$ then
 $e \leftarrow \text{next}(e)$

Doubly Connected Edge List (DCEL)

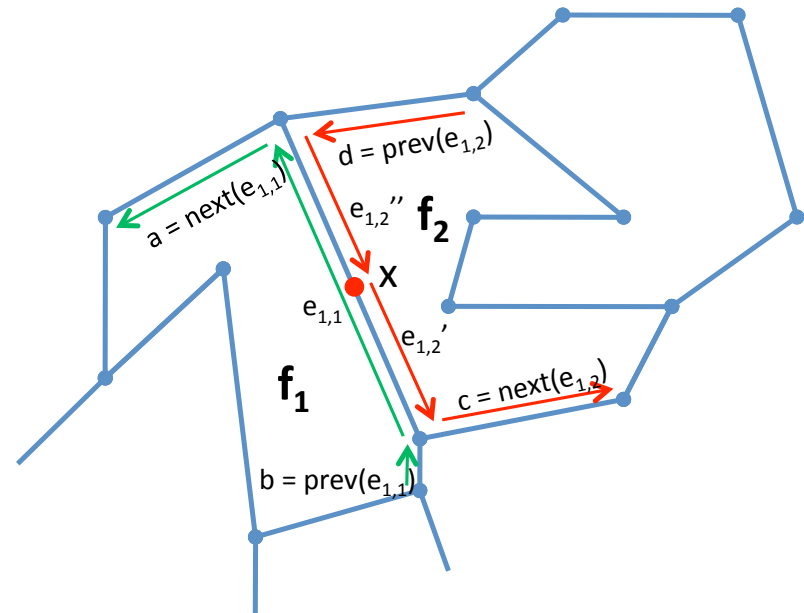
- **Traversing all edges incident on a vertex v**
 - Note: we only output the half-edges whose origin is v
 - Given: a half-edge e with the origin at v
 1. $\text{start_edge} \leftarrow e$
 2. While $\text{next}(\text{twin}(e)) \neq \text{start_edge}$ then
 $e \leftarrow \text{next}(\text{twin}(e))$

Adding a Vertex



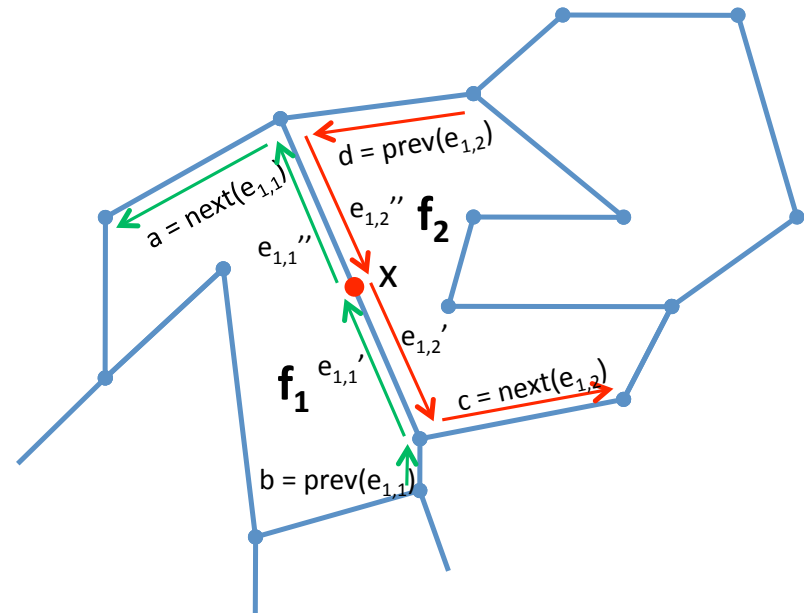
Adding a Vertex

- New vertex x
- New edges: $e_{1,2}'$ and $e_{1,2}''$
- $\text{IncidentEdge}(x) = e_{1,2}'$
- $\text{Origin}(e_{1,2}') = x$
- $\text{Next}(e_{1,2}') = \text{next}(e_{1,2})$
- $\text{Prev}(e_{1,2}') = e_{1,2}''$
- $\text{IncidentFace}(e_{1,2}') = f_2$
- $\text{Origin}(e_{1,2}'') = \text{origin}(e_{1,2})$
- $\text{Next}(e_{1,2}'') = e_{1,2}'$
- $\text{Prev}(e_{1,2}'') = \text{prev}(e_{1,2})$
- $\text{IncidentFace}(e_{1,2}'') = f_2$
- $\text{Next}(\text{Prev}(e_{1,2})) = e_{1,2}''$
- $\text{Prev}(\text{Next}(e_{1,2})) = e_{1,2}'$
- Delete edge $e_{1,2}$



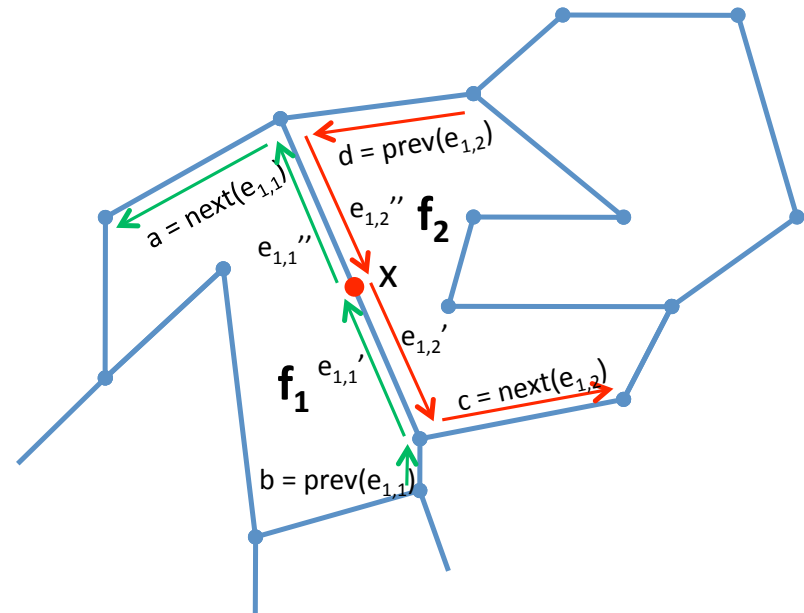
Adding a Vertex

- New edges: $e_{1,1}'$ and $e_{1,1}''$
- $\text{Origin}(e_{1,1}') = \text{origin}(e_{1,1})$
- $\text{Next}(e_{1,1}') = e_{1,1}''$
- $\text{Prev}(e_{1,1}') = \text{prev}(e_{1,1})$
- $\text{IncidentFace}(e_{1,1}') = f_1$
- $\text{Origin}(e_{1,1}'') = e_{1,1}'$
- $\text{Next}(e_{1,1}'') = \text{next}(e_{1,1})$
- $\text{Prev}(e_{1,1}'') = e_{1,1}'$
- $\text{IncidentFace}(e_{1,1}'') = f_1$
- $\text{Next}(\text{prev}(e_{1,1})) = e_{1,1}'$
- $\text{Prev}(\text{next}(e_{1,1})) = e_{1,1}''$
- $\text{Twin}(e_{1,2}') = e_{1,1}'$
- $\text{Twin}(e_{1,1}') = e_{1,2}'$
- $\text{Twin}(e_{1,2}'') = e_{1,1}''$
- $\text{Twin}(e_{1,1}'') = e_{1,2}''$
- Delete edge $e_{1,1}$

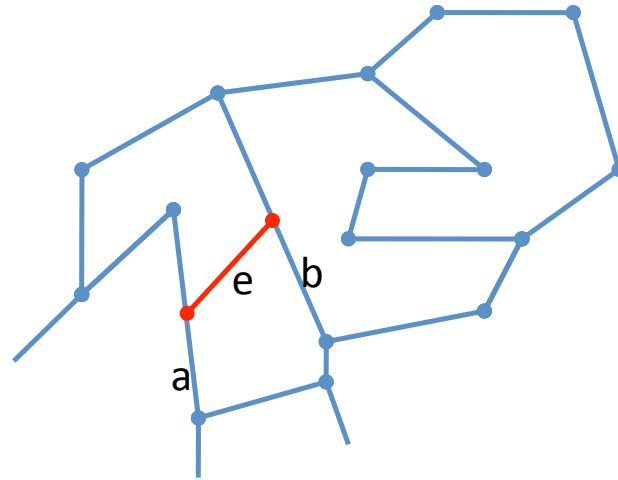


Adding a Vertex

- If $e_{1,1}$ was starting edge of f_1 , need to change it to either one of the new edges
- If $e_{1,2}$ was starting edge of f_2 , need to change it to either one of the new edges



Other Operations on DCEL



- **Add an Edge**

- Planar subdivision
- e is added
- DCEL can be updated in constant time once the edges a and b are known