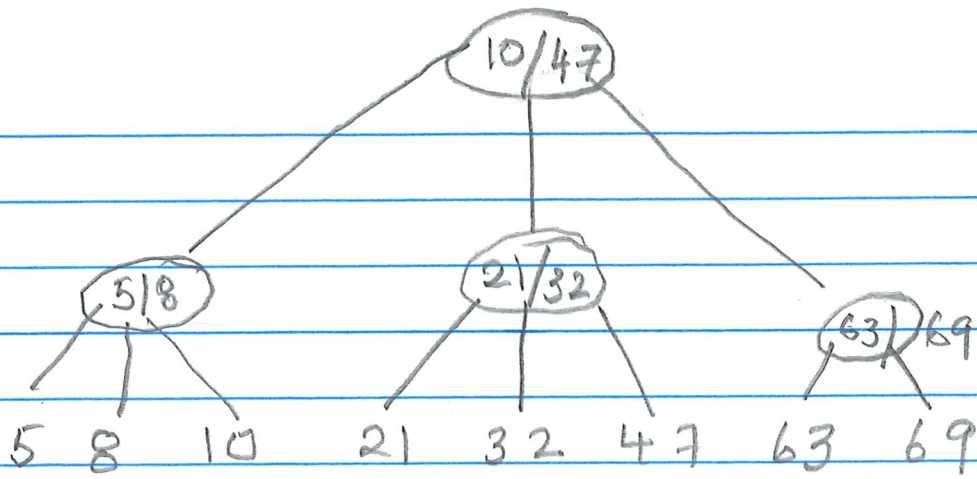


2-3 Trees

- leaf nodes contain data item & all leaves are at the same level.
- each internal nonleaf node
 - has 2 or 3 children
 - two search values
 - largest item in left subtree
 - largest item in middle subtree
- Smallest 2-3 tree
 - only one leaf node.
- balanced & ordered.



height of tree

- n data items

- all nodes have two children

- height $\leq \lfloor \log_2(n+1) \rfloor - 1$

- all nodes have three children

- height $= \lfloor \log_3(2n+1) \rfloor - 1$

Time - For a general tree, $O(\text{height})$

$$\log_3(2n+1) - 1 \leq \text{height} \leq \lfloor \log_2(n+1) \rfloor - 1$$

Time to search 1 key: $O(\text{height})$

Time to list n elements in sorted order
: $O(n)$.

Insertions in 2-3 trees

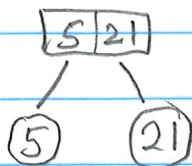
- insert new leaf in the appropriate place.
- Repeat until all nonleaf nodes have 2 or 3 children
 - if there is a node with 4 children, split the parent into two parent nodes with 2 children each
 - if the root node is split, add a new root
- Adjust the search values along the insertion paths.

Example:

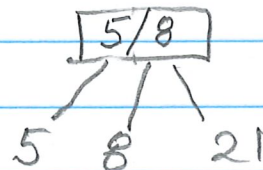
Insert 5



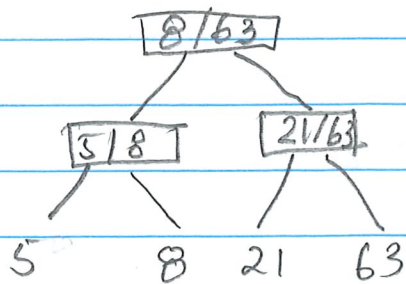
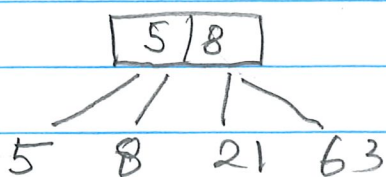
Insert 21



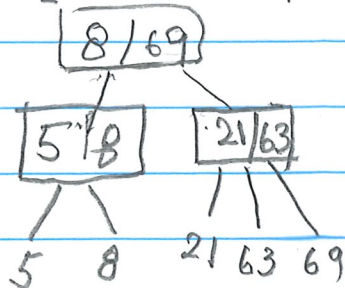
Insert 8



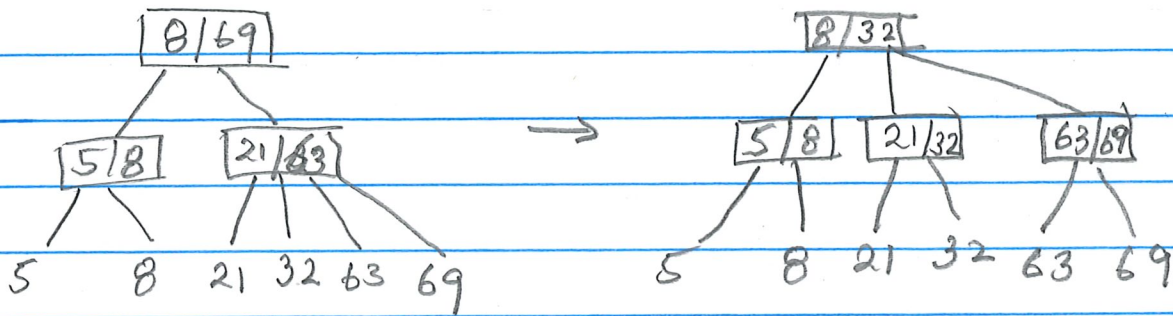
Insert 63



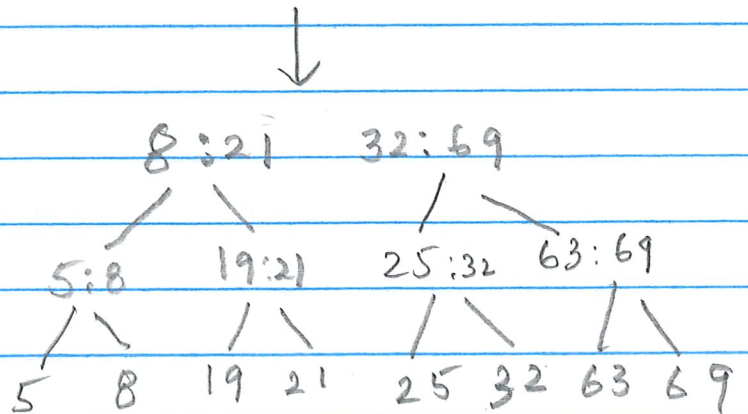
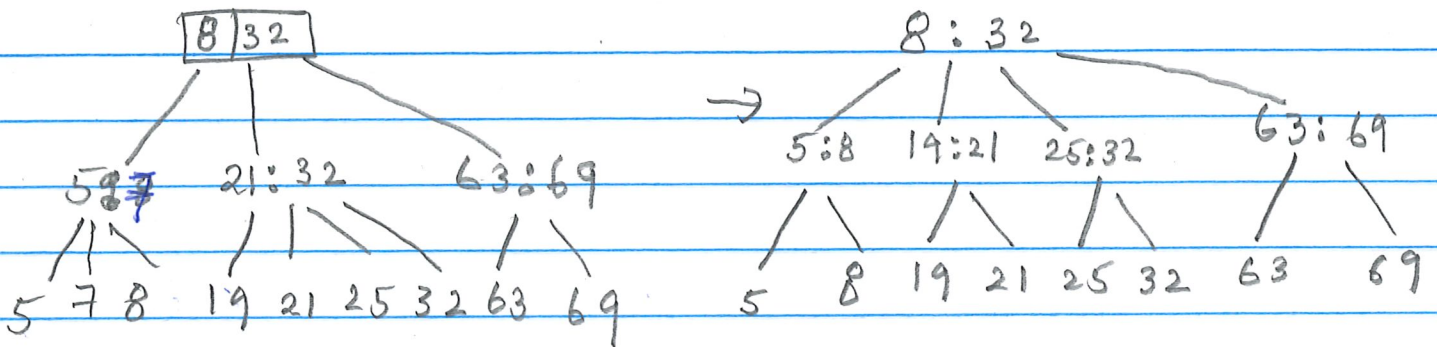
Insert 69



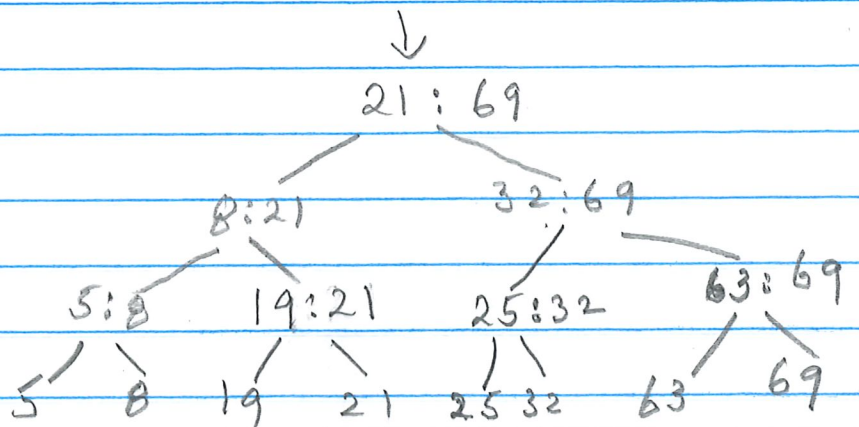
Insert 32



Insert 7, 19, 25



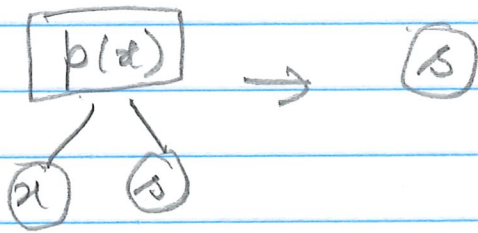
Time to insert
is $O(\text{height})$



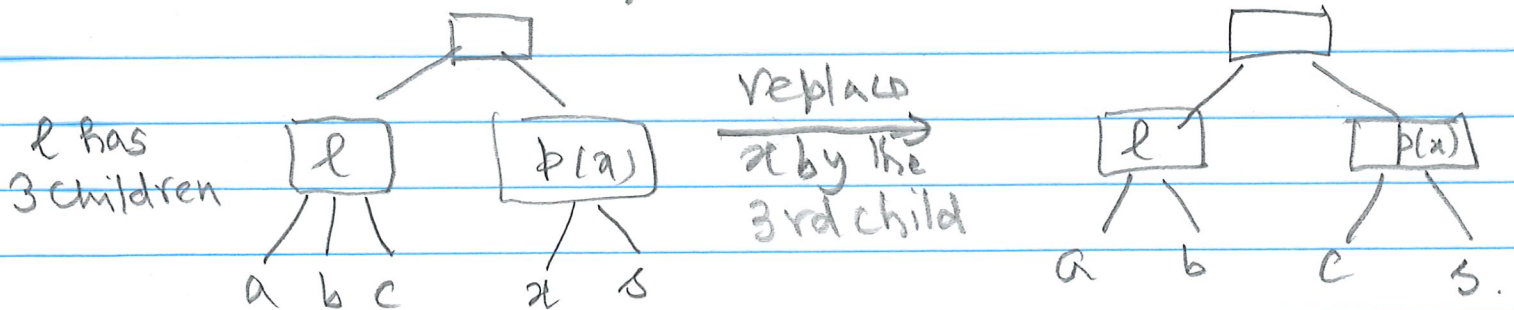
Deletions in 2-3 trees

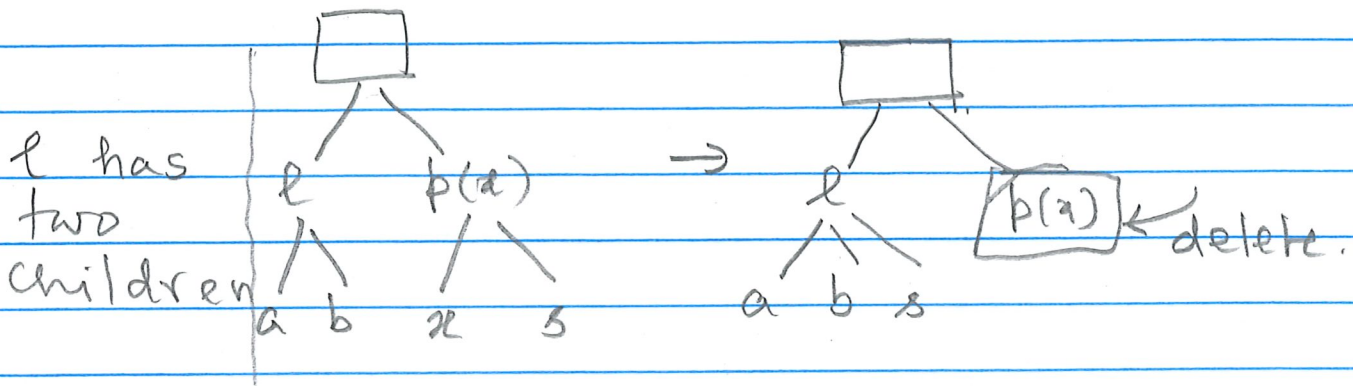
Delete x from the tree; $p(x)$: parent of x .

- if x is the root: delete x
- if $p(x)$ has 3 children, delete x .
- if $p(x)$ has 2 children (one is x + the other child is s)
 - if $p(x)$ is the root



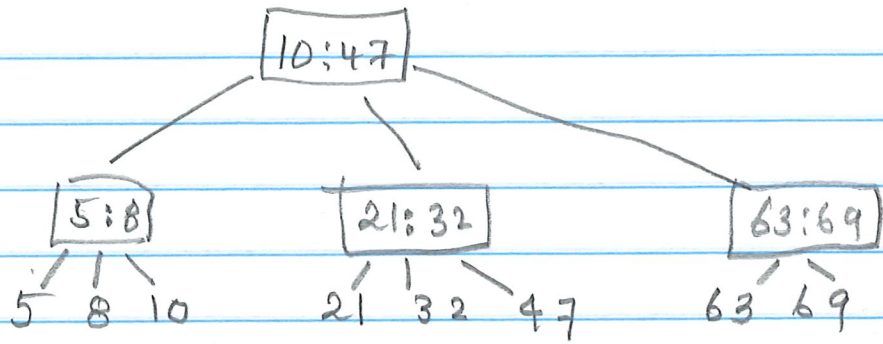
- $p(x)$ is not the root node. Let l be the left sibling of $p(x)$ & r be the right sibling of $p(x)$; (note l or r may not exist)



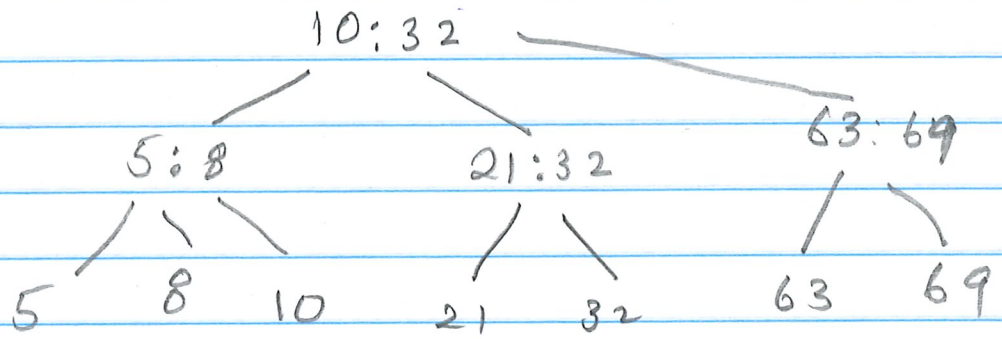


- Steps :
- remove x
 - combine $p(x)$ with l
 - make s a child of l
 - rename $p(x)$ to x
 - recursively remove x

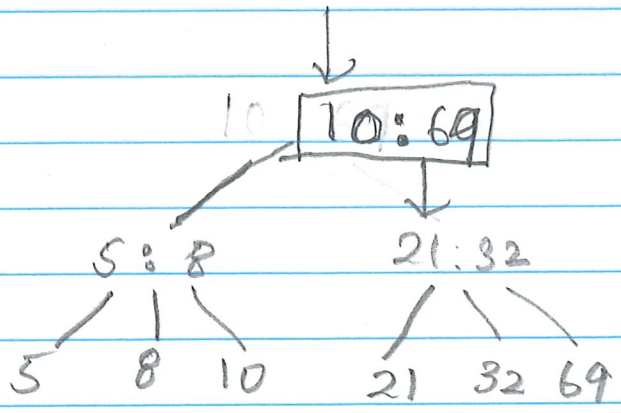
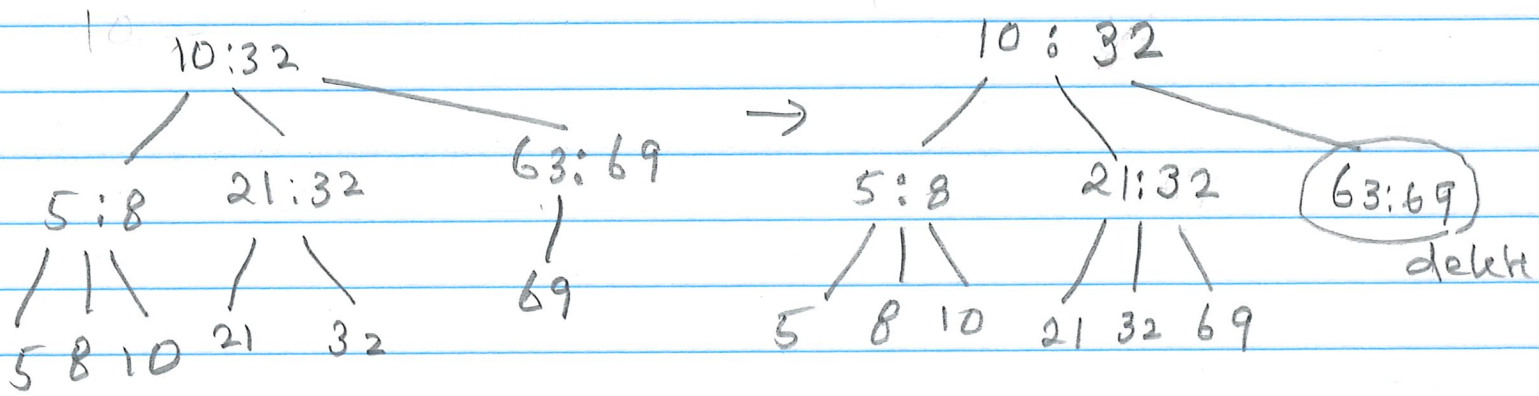
Example



delete 47



delete 63



Example

