

Supplementary Exercises- Mathematical Induction

Summations

- Prove by induction on positive n that $\sum_{i=1}^n \frac{1}{i(i+1)} = \frac{n}{n+1}$.
- Prove by induction that for all $n \geq 1$, $\sum_{i=1}^n i^3 = \frac{n^2(n+1)^2}{4}$.
- Using a proof by induction, show that $\sum_{i=1}^n (i)(i!) = (n+1)! - 1$, for all $n \geq 1$.
- Prove that $\sum_{i=0}^n 2^i = 2^{n+1} - 1$ for all $n \geq 0$ by induction.

Divisibility

- Show that $n^3 + 2n$ is divisible by 3 for all $n \geq 1$ by induction.
- Show that $n^4 - 4n^2$ is divisible by 3 for all $n \geq 2$ by induction.
- Show that $7^n - 1$ is divisible by 6 for all $n \geq 1$ by induction.
- Let a, b, m be positive integers such that $m \mid (a - b)$. Prove by induction on $k \geq 0$ that $m \mid (a^k - b^k)$.

Inequalities

- Prove that $\sum_{i=1}^n \frac{1}{\sqrt{i}} \geq \sqrt{n}$ for all $n \geq 1$ by induction.
- Prove by induction that for $n \in \mathbb{Z}^+$, $\sum_{i=1}^n \frac{1}{i^2} \leq 2 - \frac{1}{n}$.
- Using induction, prove that $3^n < n!$ for all $n \geq 7$.

Sequences and Strong Induction

- Let $F_0 = F_1 = 1$ and $F_n = F_{n-1} + F_{n-2}$ for all $n \geq 2$. Prove for all $n \geq 0$:

$$- \sum_{i=0}^n F_i = F_{n+2} - 1.$$

$$- \sum_{i=0}^n (F_i)^2 = F_n \cdot F_{n+1}.$$

$$- \sum_{i=0}^{\lfloor n/2 \rfloor} \binom{n-i}{i} = F_n.$$

$$- \gcd(F_{n+1}, F_n) = 1.$$

- Show that postage of 24 cents or more can be achieved by using only 5-cent and 7-cent stamps.
- One way of generating the n^{th} Fibonacci number is by using a recursive algorithm like the following.

```
int fibonacci(int n)    // generates Fn
  if ((n = 0) OR (n = 1)) then
    return 1;
  else
    return fibonacci(n-1) + fibonacci(n-2);
  end if
end fibonacci;
```

Effectively, the algorithm mimics the definition of the Fibonacci sequence shown above.

In computing science, the efficiency of an algorithm is measured by the amount of computing time expended as a function of the input. In this example, the input is the integer n , so we will let T_n represent the amount of time expended to run `fibonacci(n)`.

Note that this algorithm is *recursive* which means that it calls an instance of itself to help solve the problem. Function calls are very costly time-wise, so we can measure T_n by counting the total number of calls to `fibonacci` during execution.

For example, an initial call to `fibonacci(3)` will call `fibonacci(2)` and `fibonacci(1)`. `fibonacci(2)` will (in turn) call `fibonacci(1)` and `fibonacci(0)`. There are a total of 5 calls to `fibonacci`, which means $T_3 = 5$. (Note also that $T_2 = 3$.)

- Determine the values of T_0 and T_1 .
- Write a recursive definition for T_n , for $n \geq 2$, $n \in \mathbb{N}$.
- Prove by induction that for all $n \in \mathbb{N}$, $T_n = 2F_n - 1$.