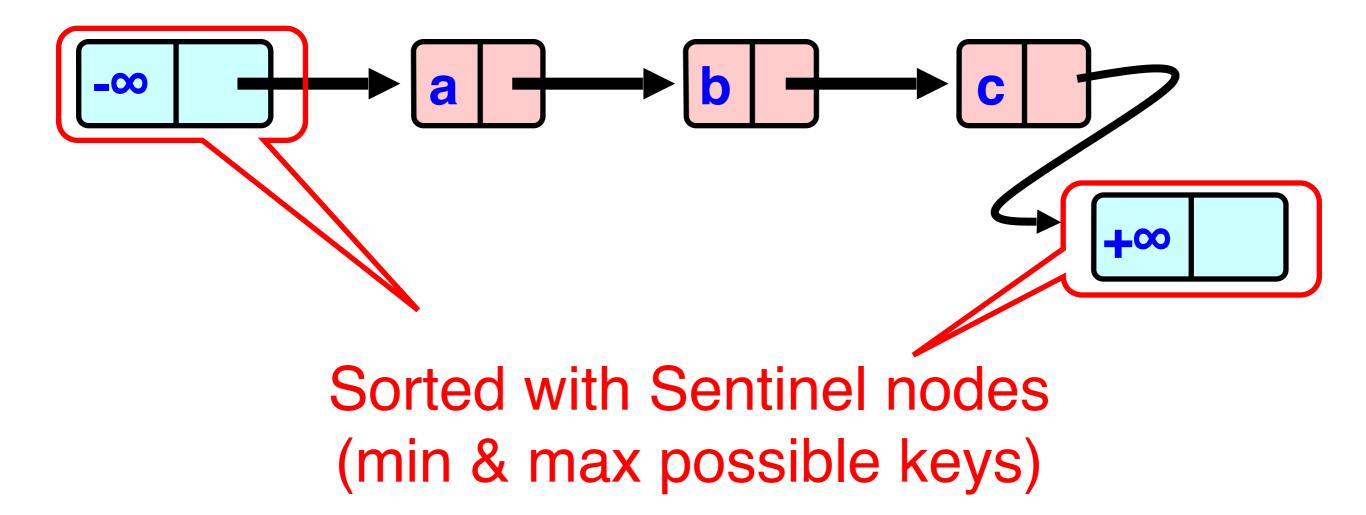
The List-Based Set

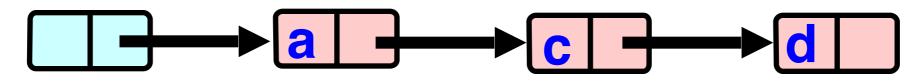


Invariants

- Sentinel nodes
 - tail reachable from head
- Sorted
- No duplicates

Sequential List Based Set

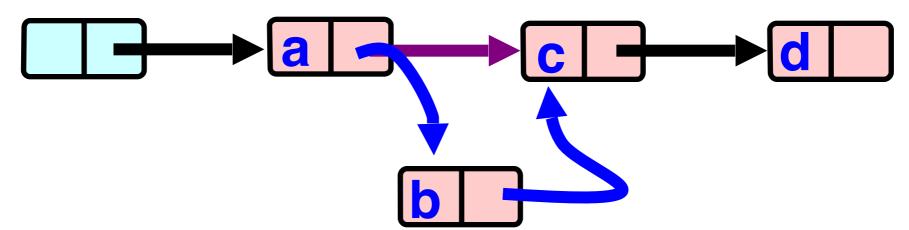
Add()

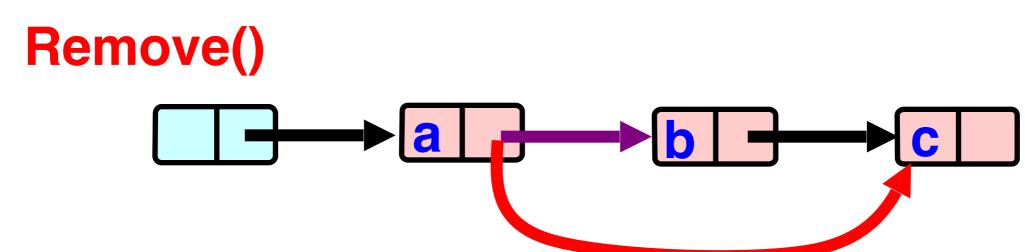


Remove()

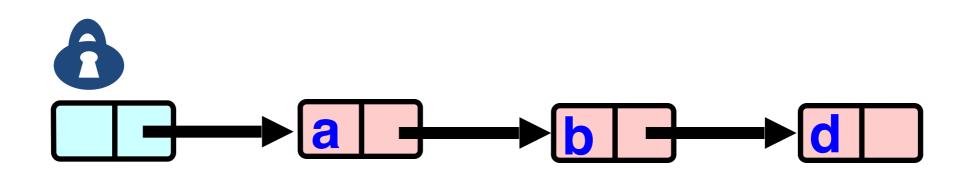
Sequential List Based Set



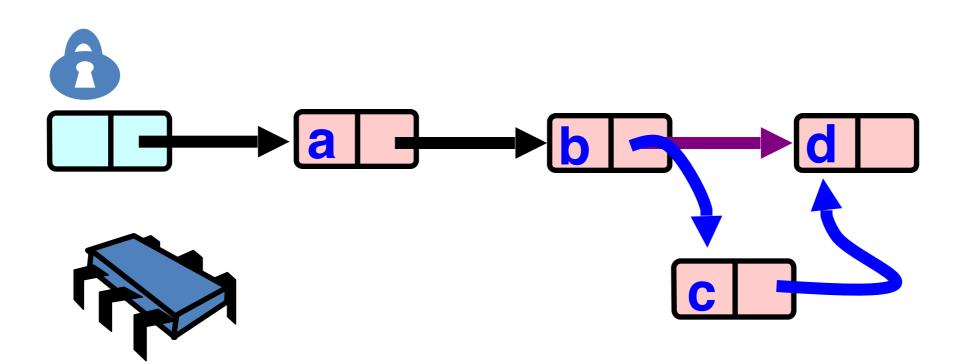




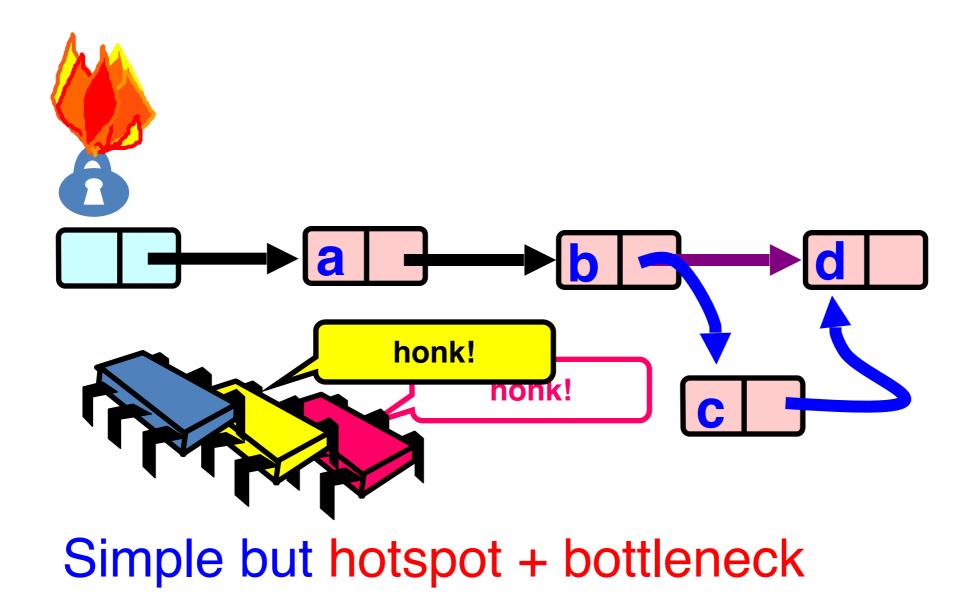
Coarse Grained Locking



Coarse Grained Locking



Coarse Grained Locking



Coarse-Grained Locking

Easy, same as synchronized methods

 "One lock to rule them all ..."

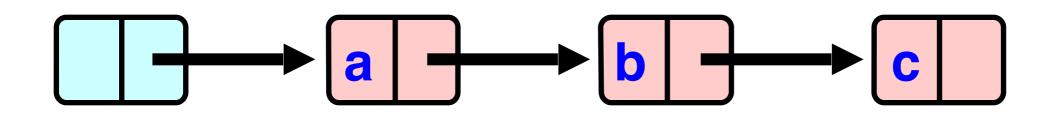
Coarse-Grained Locking

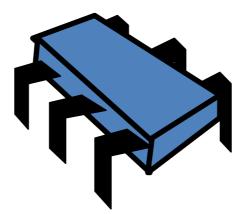
- Easy, same as synchronized methods

 "One lock to rule them all ..."
- Simple, clearly correct
 Deserves respect!
- Works poorly with contention
 - Queue locks help
 - But bottleneck still an issue

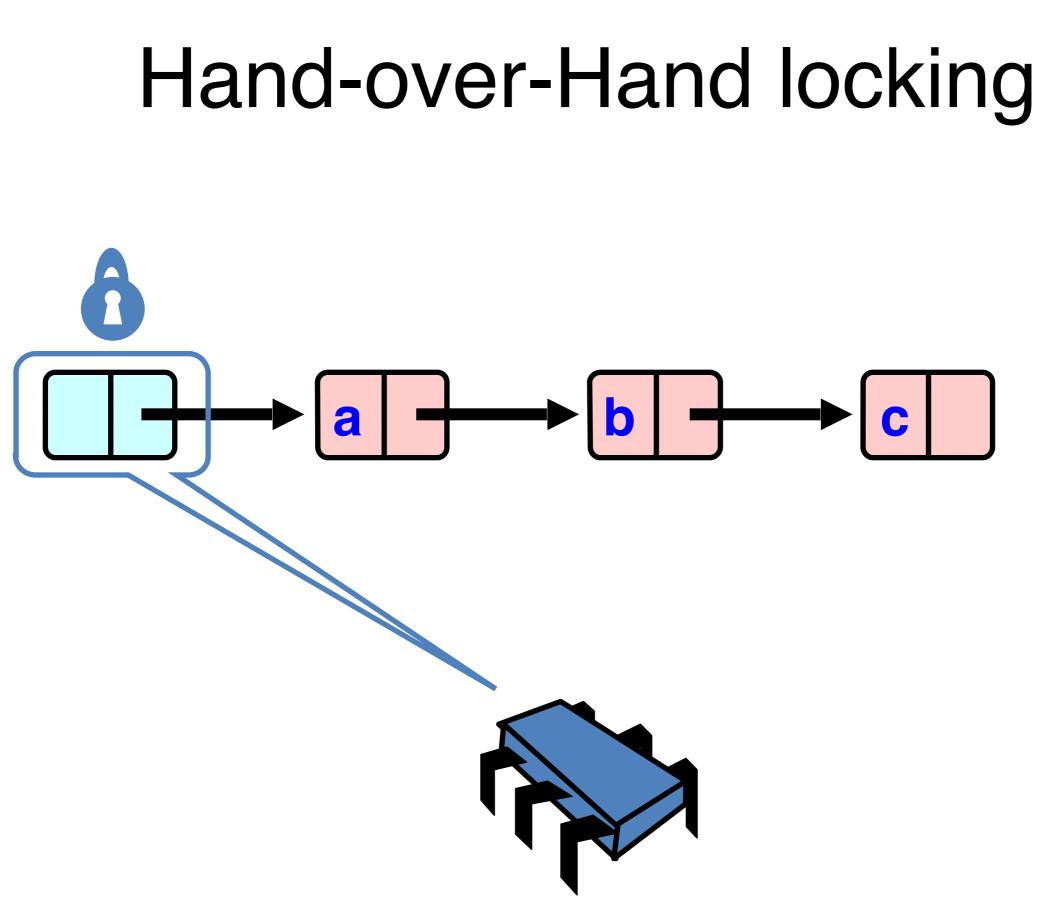
Fine-grained Locking

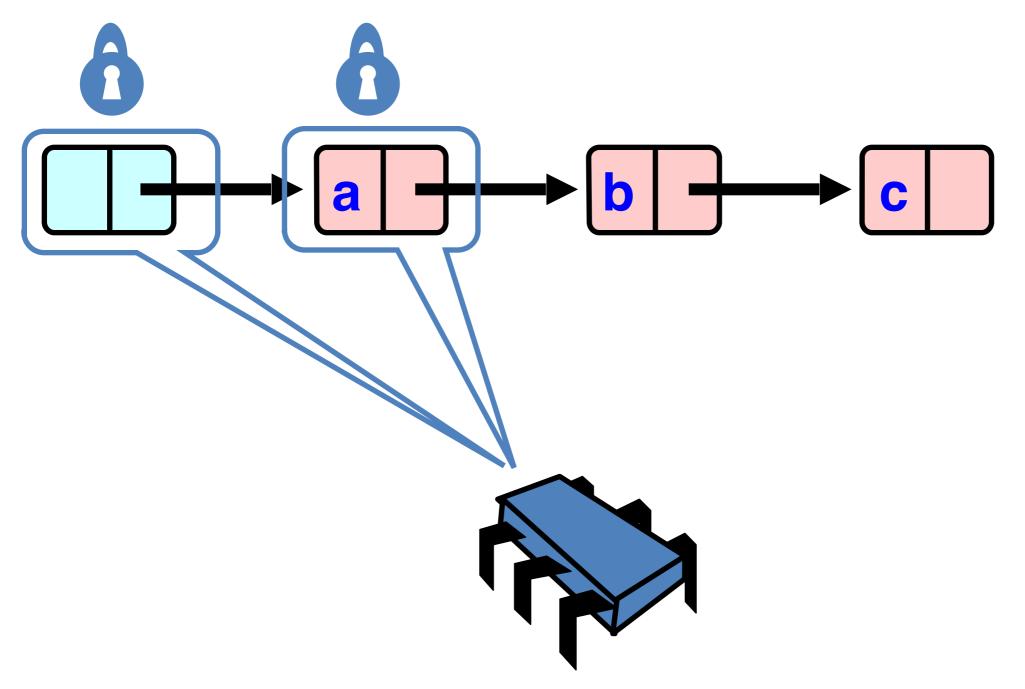
- Requires careful thought
- Split object into pieces
 - Each piece has own lock
 - Methods that work on disjoint pieces need not exclude each other

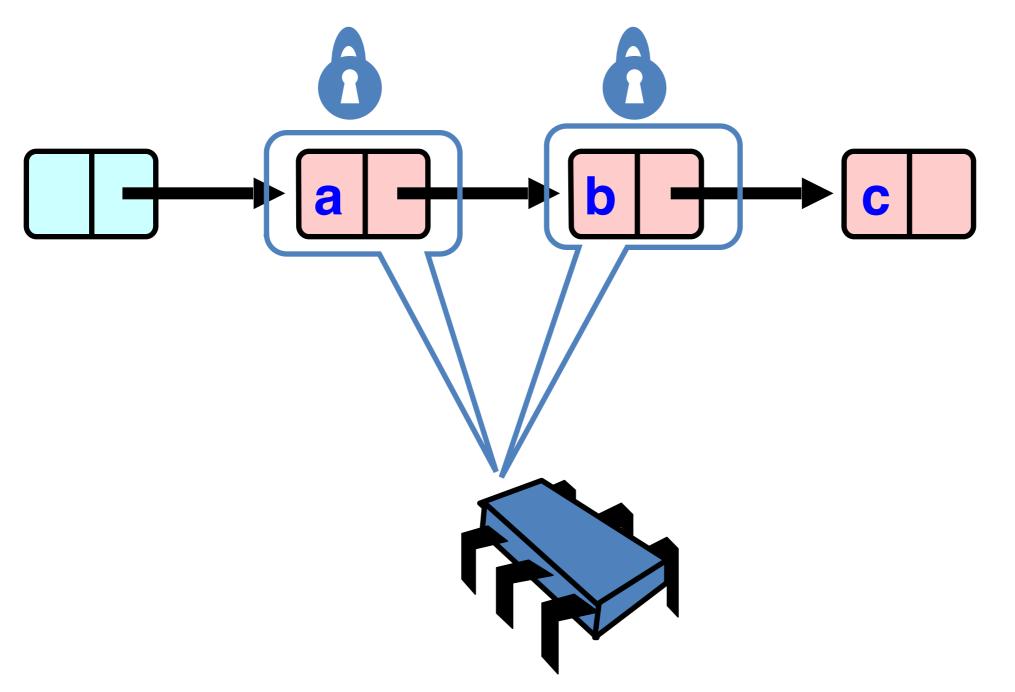


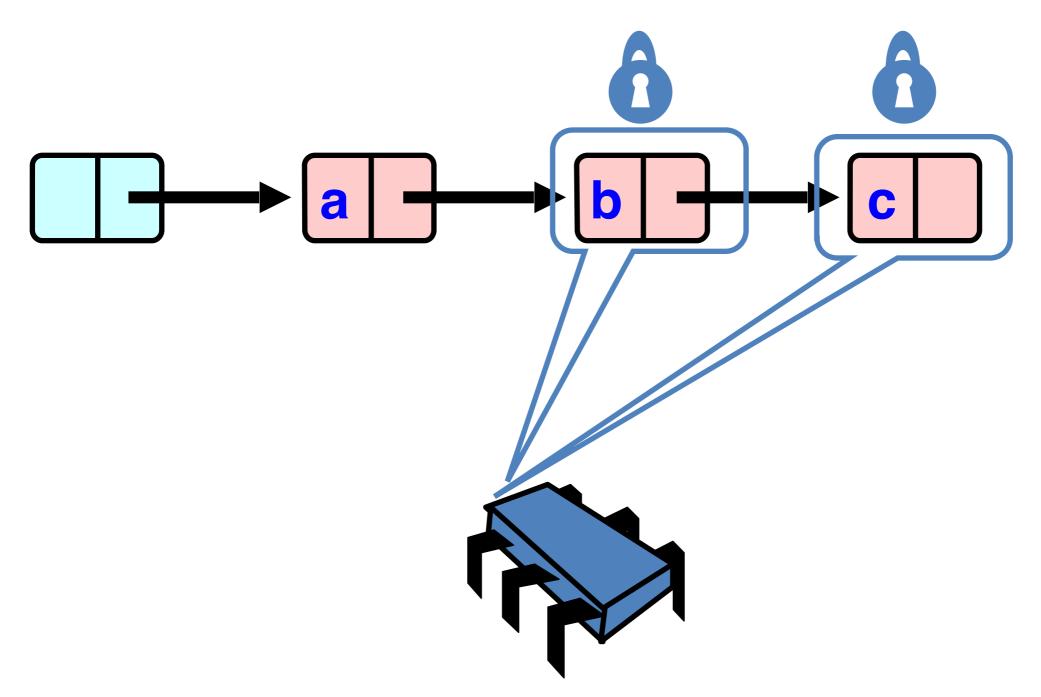


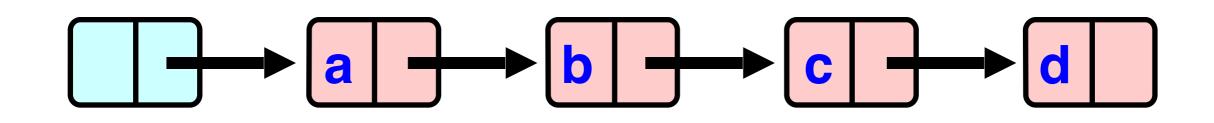
Art of Multiprocessor Programming

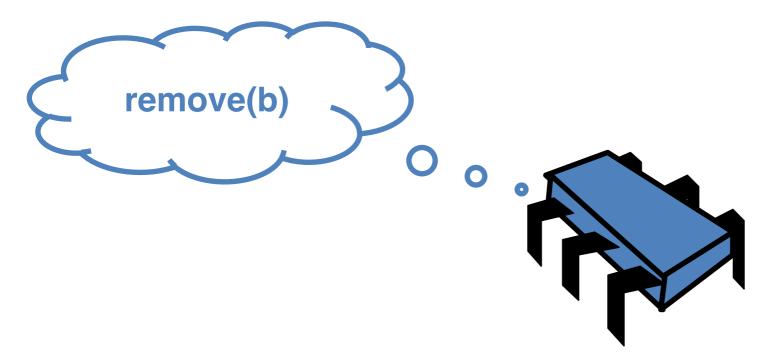


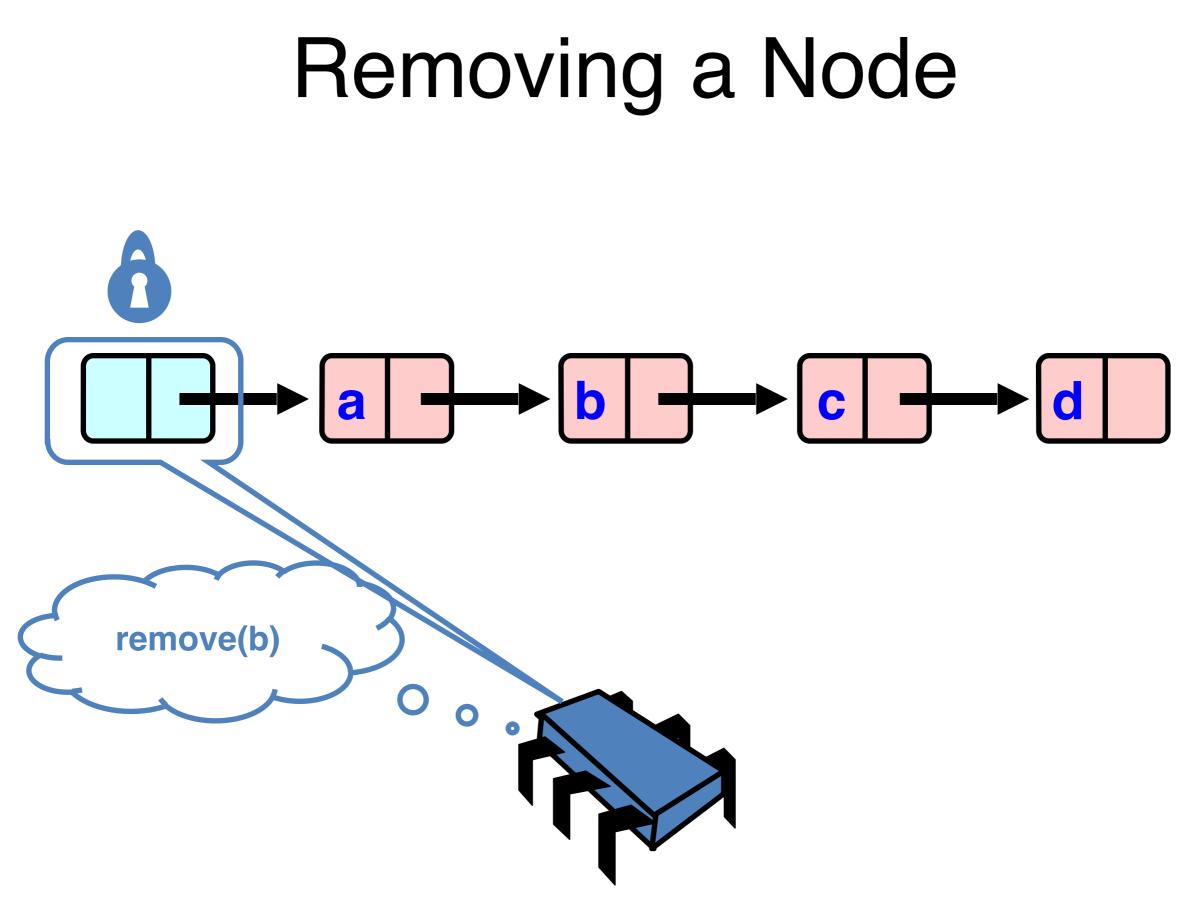


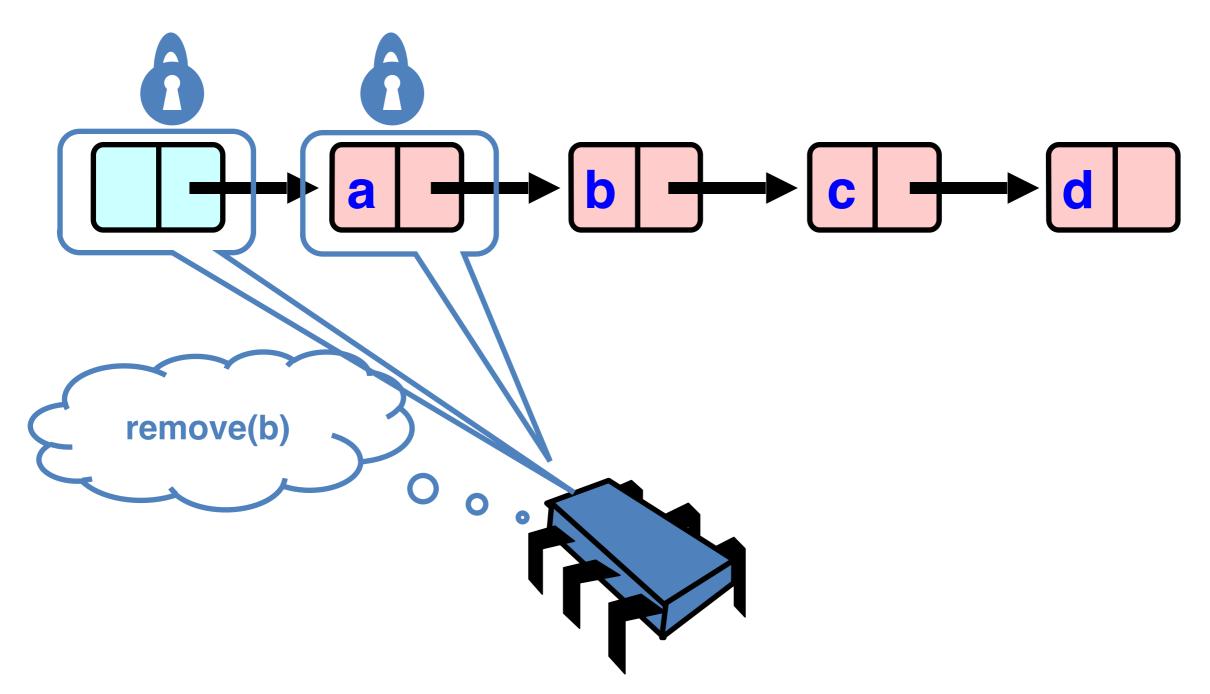


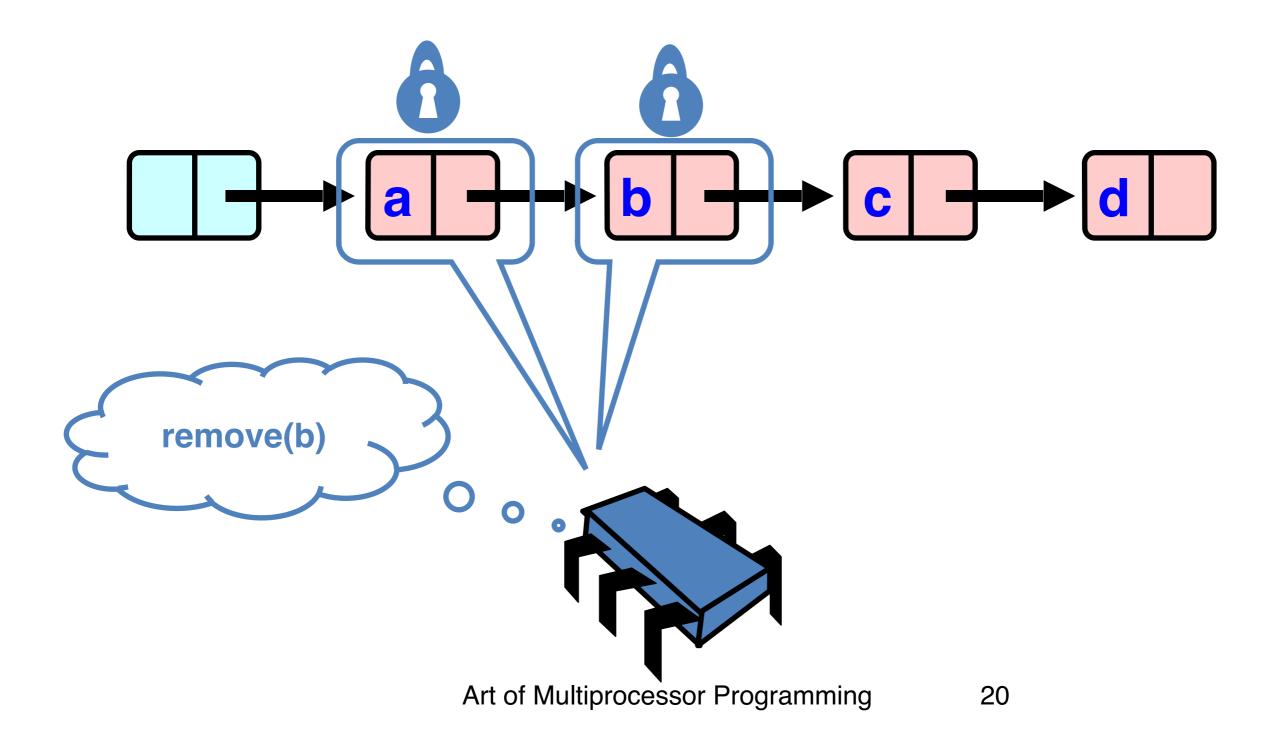


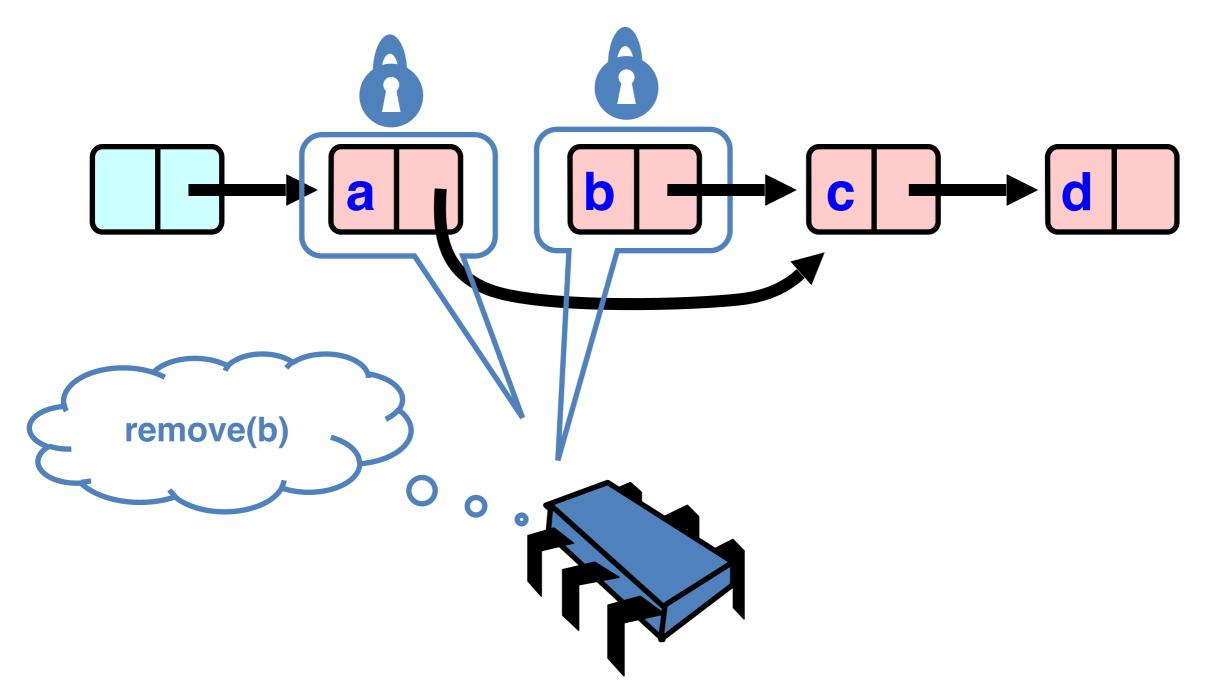


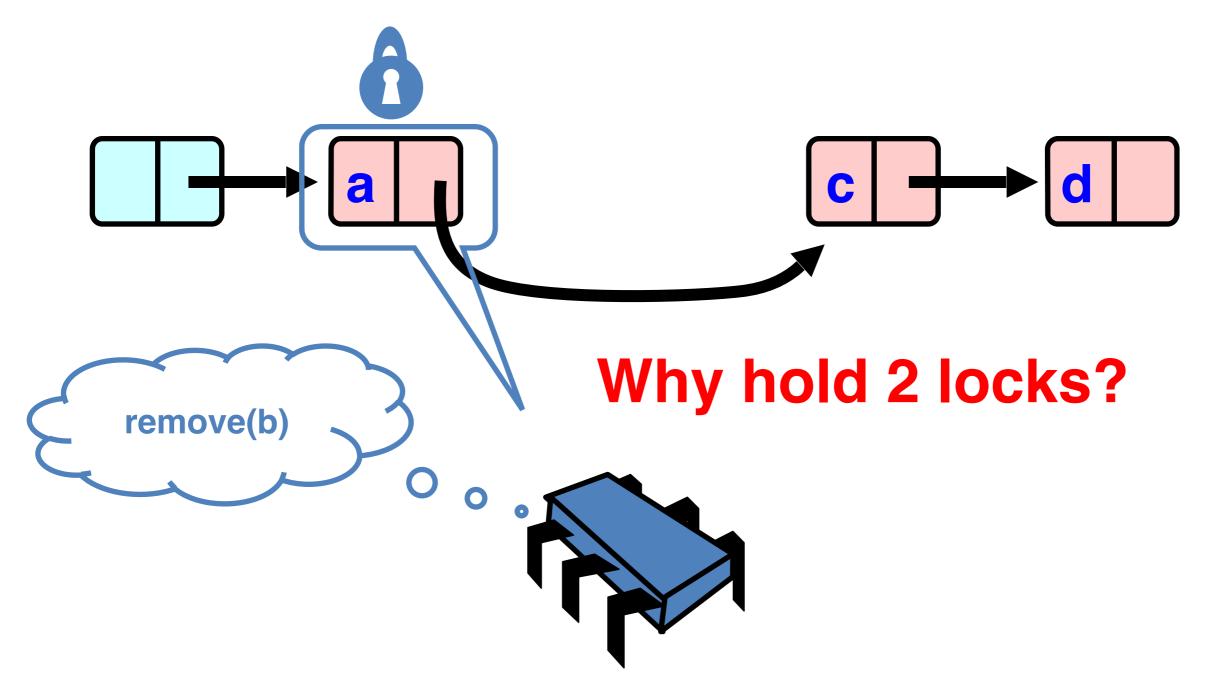




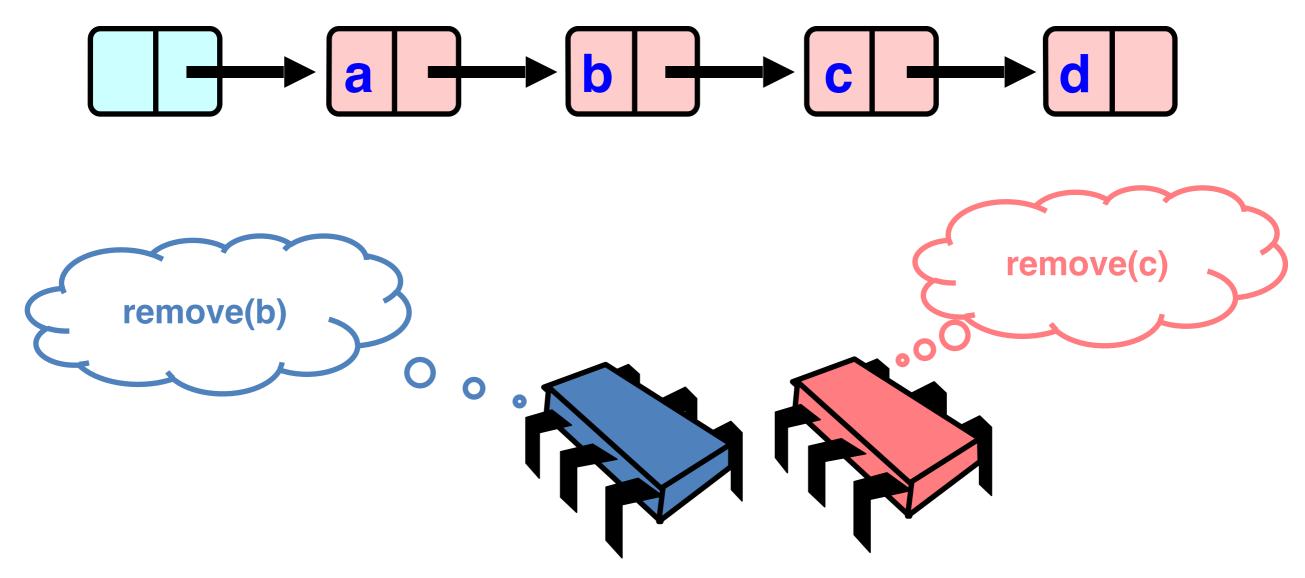


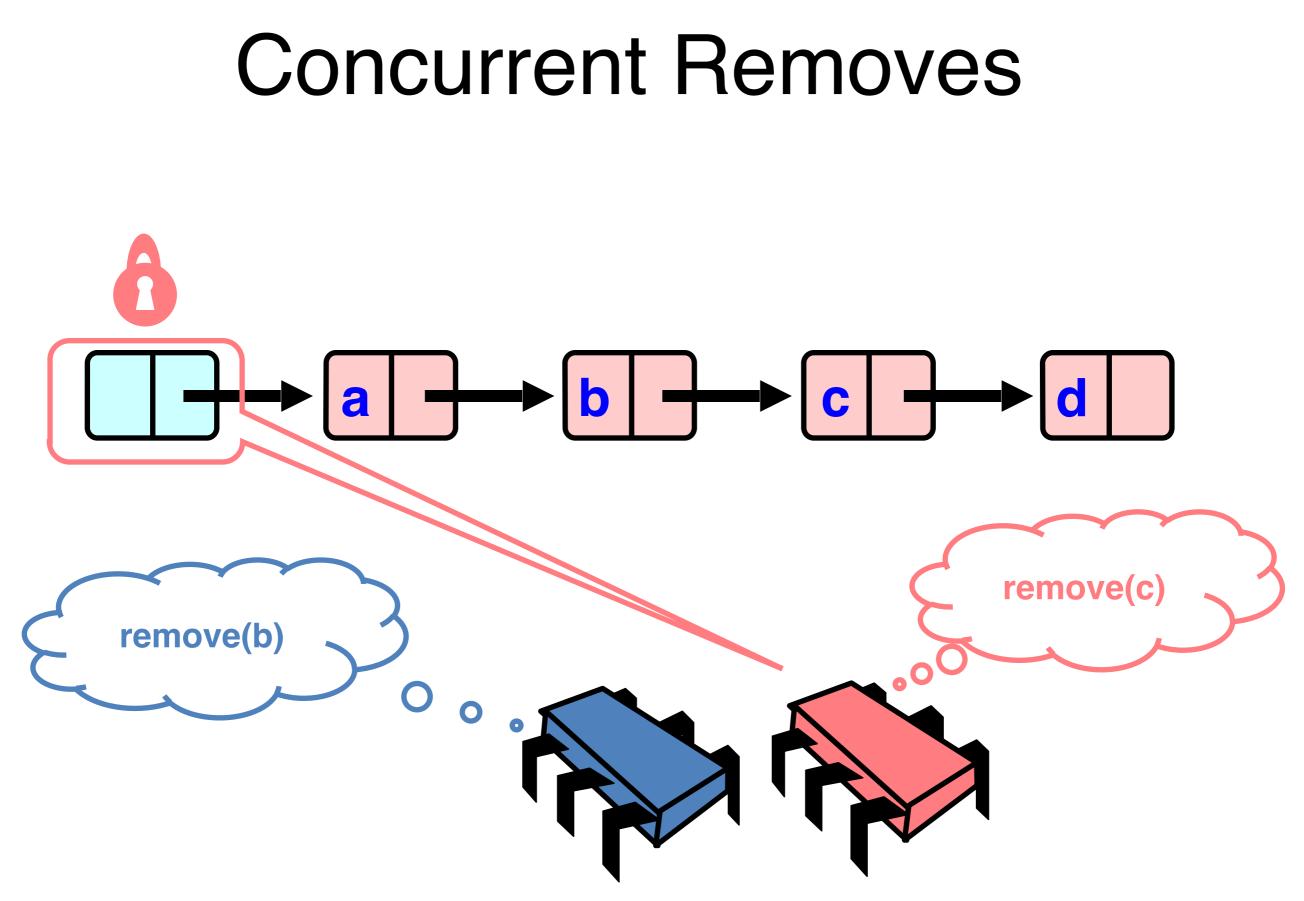


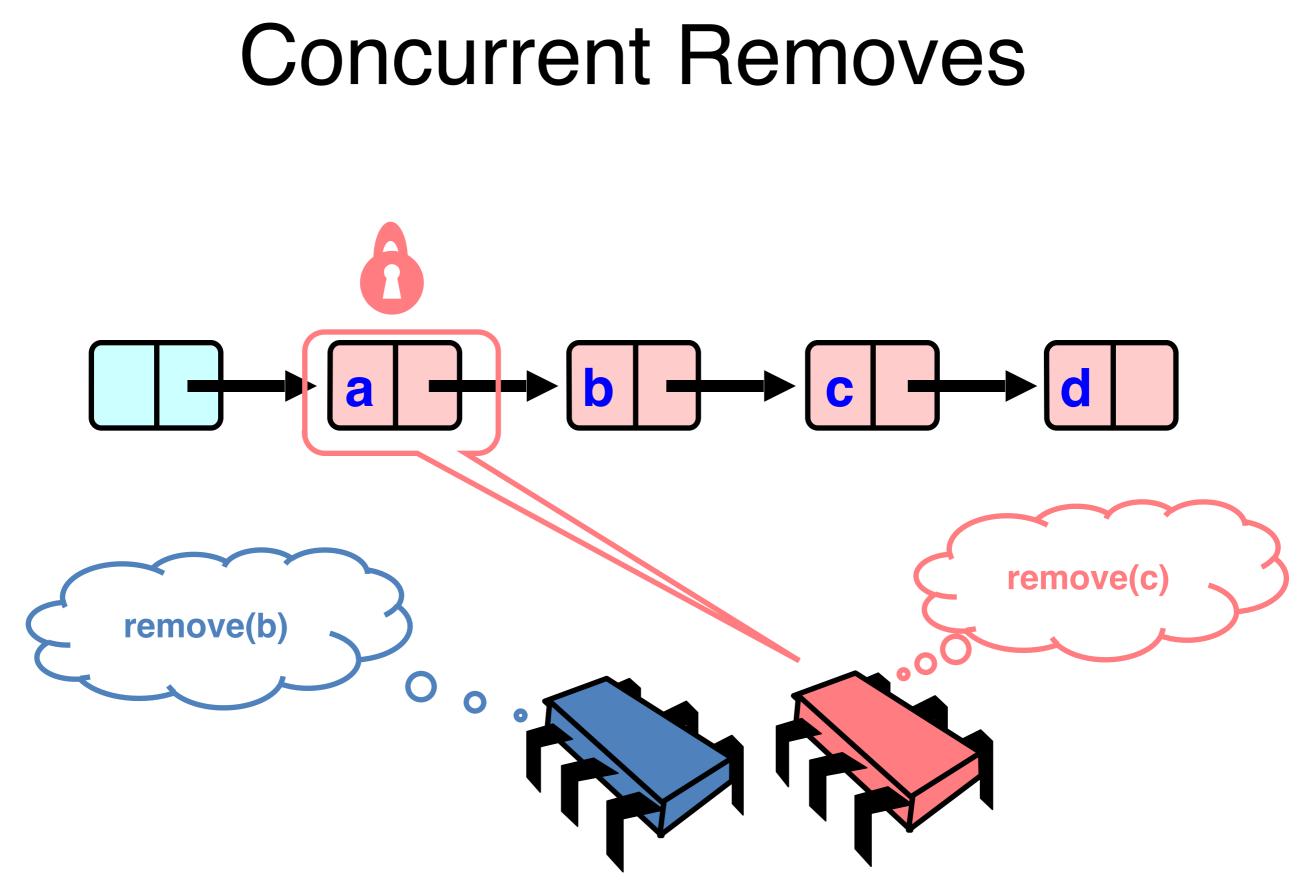


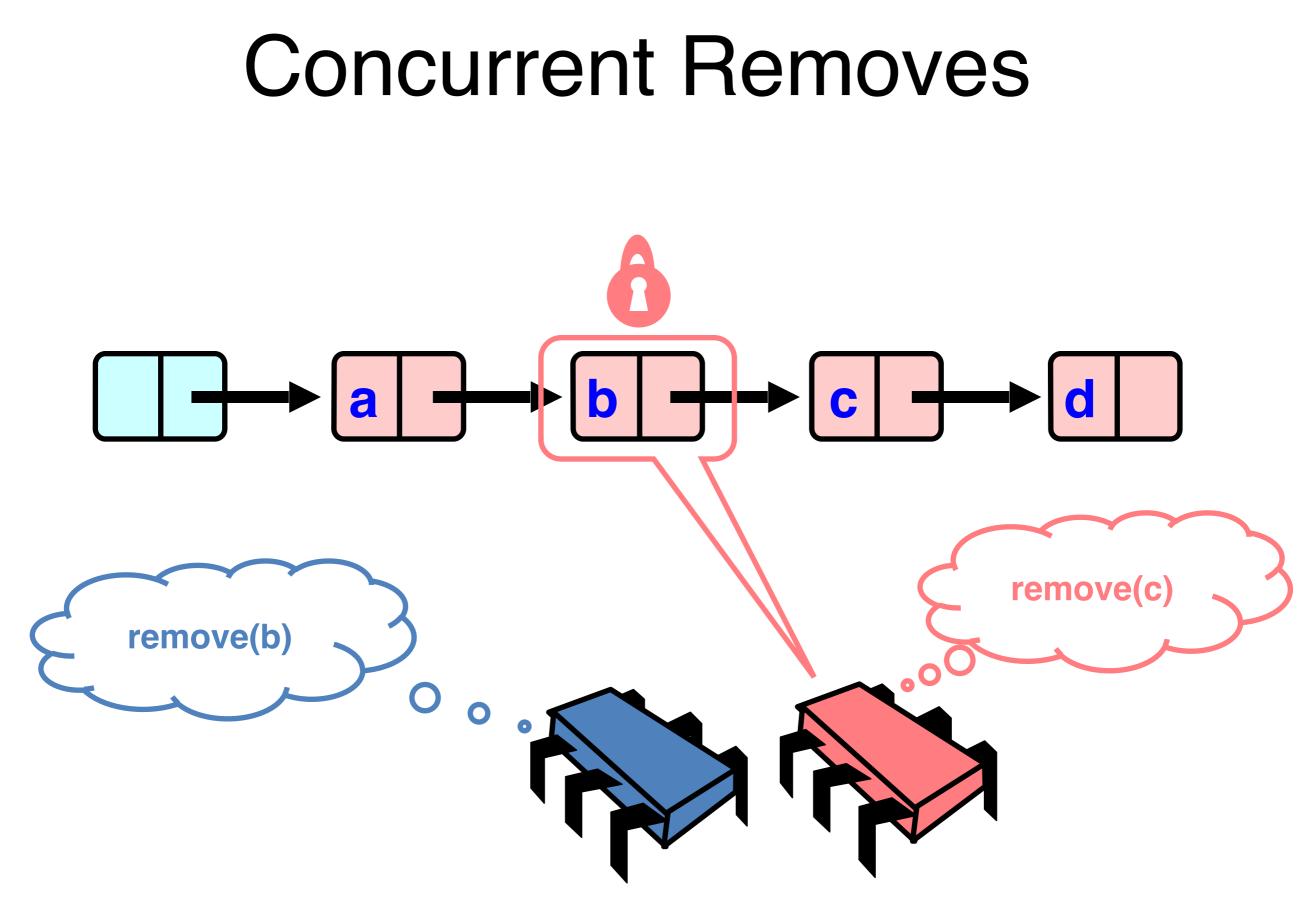


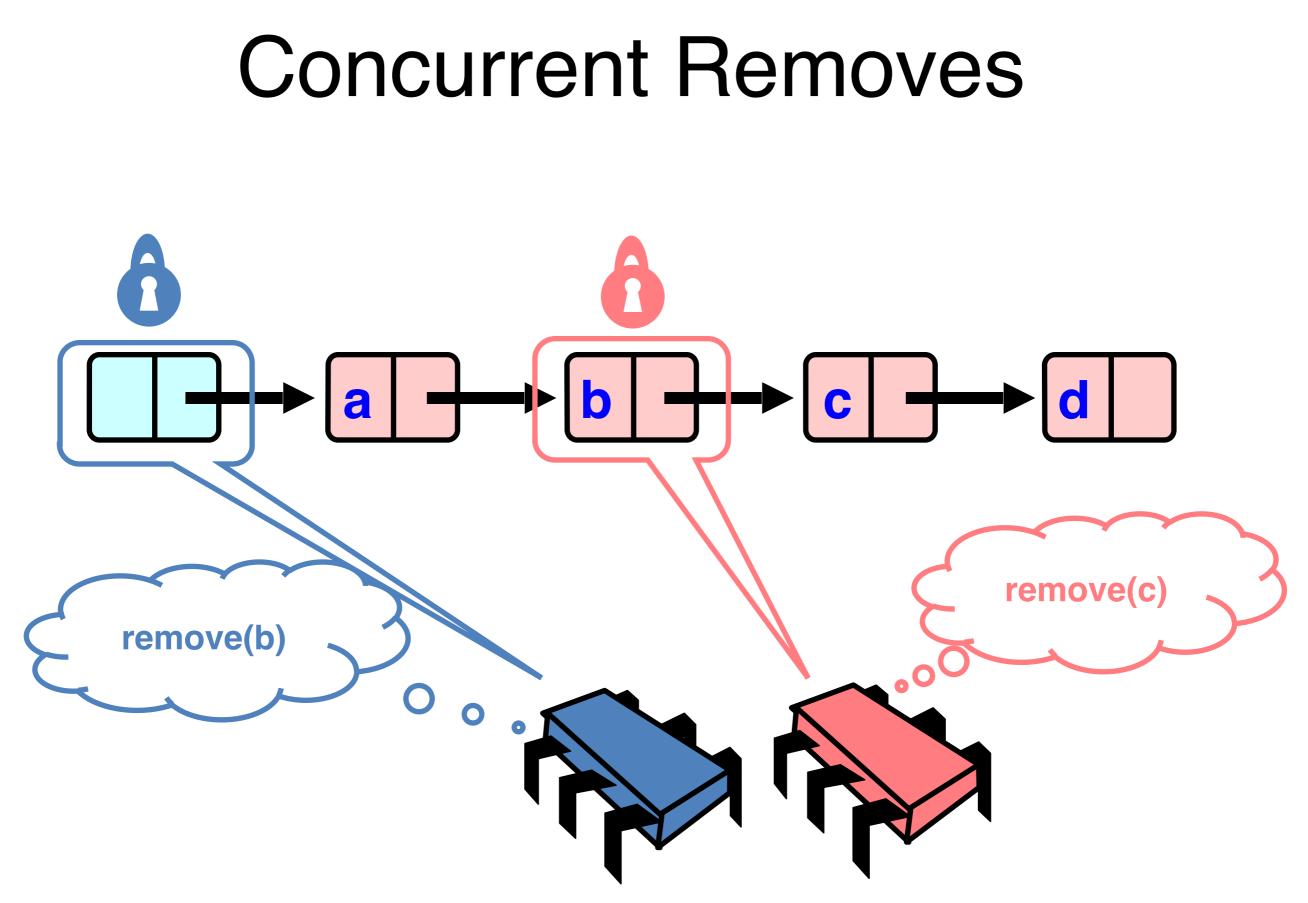
Concurrent Removes

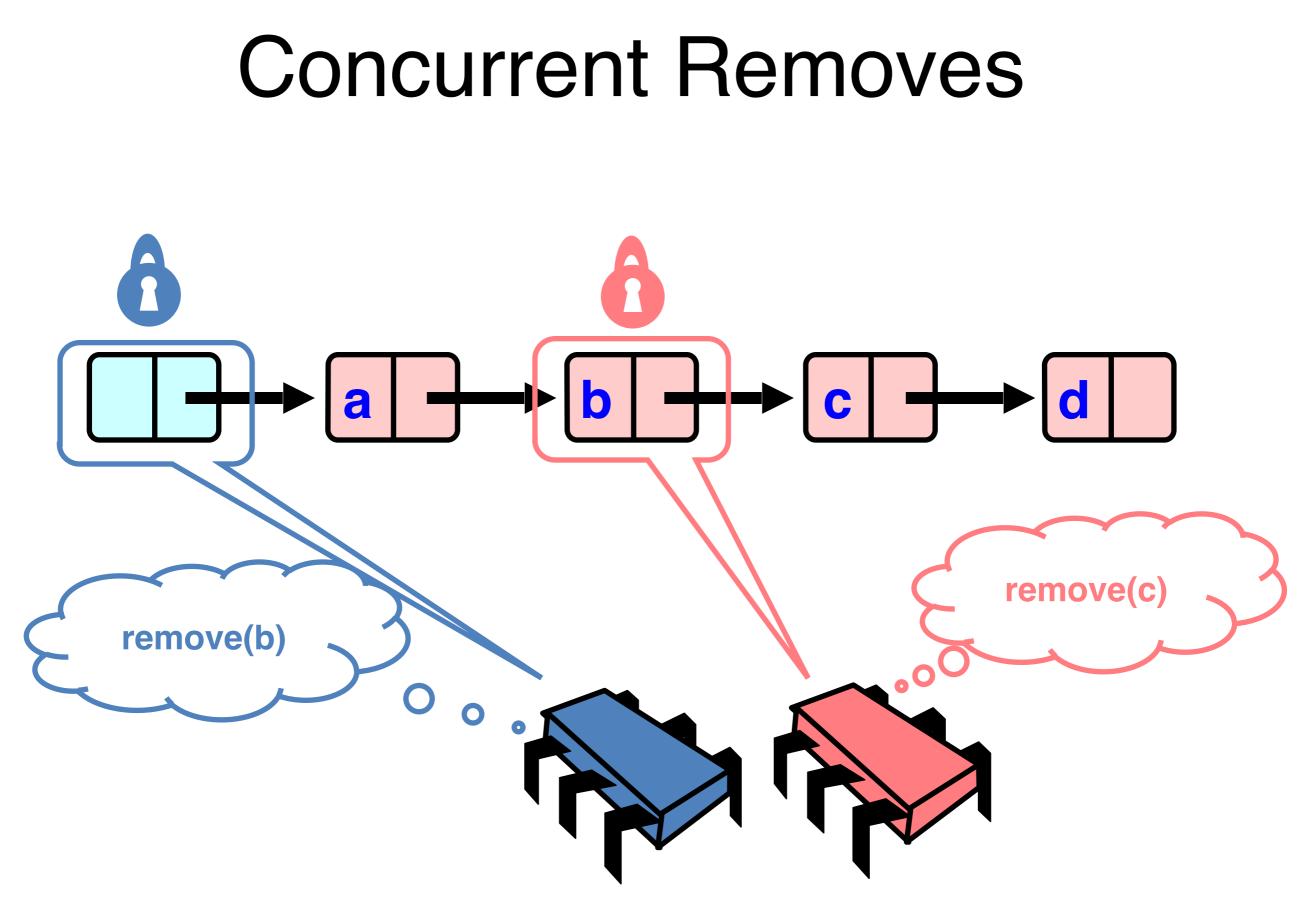


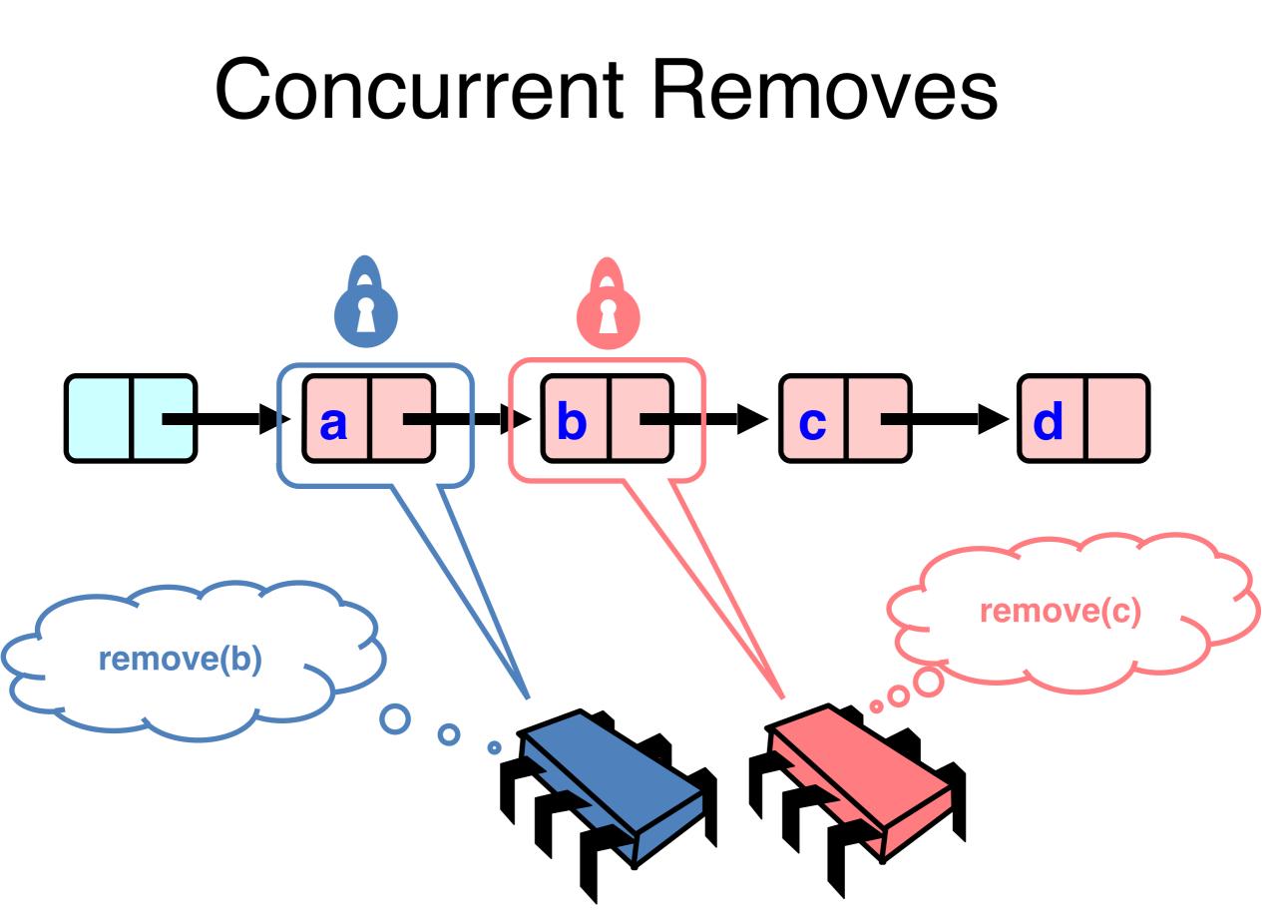


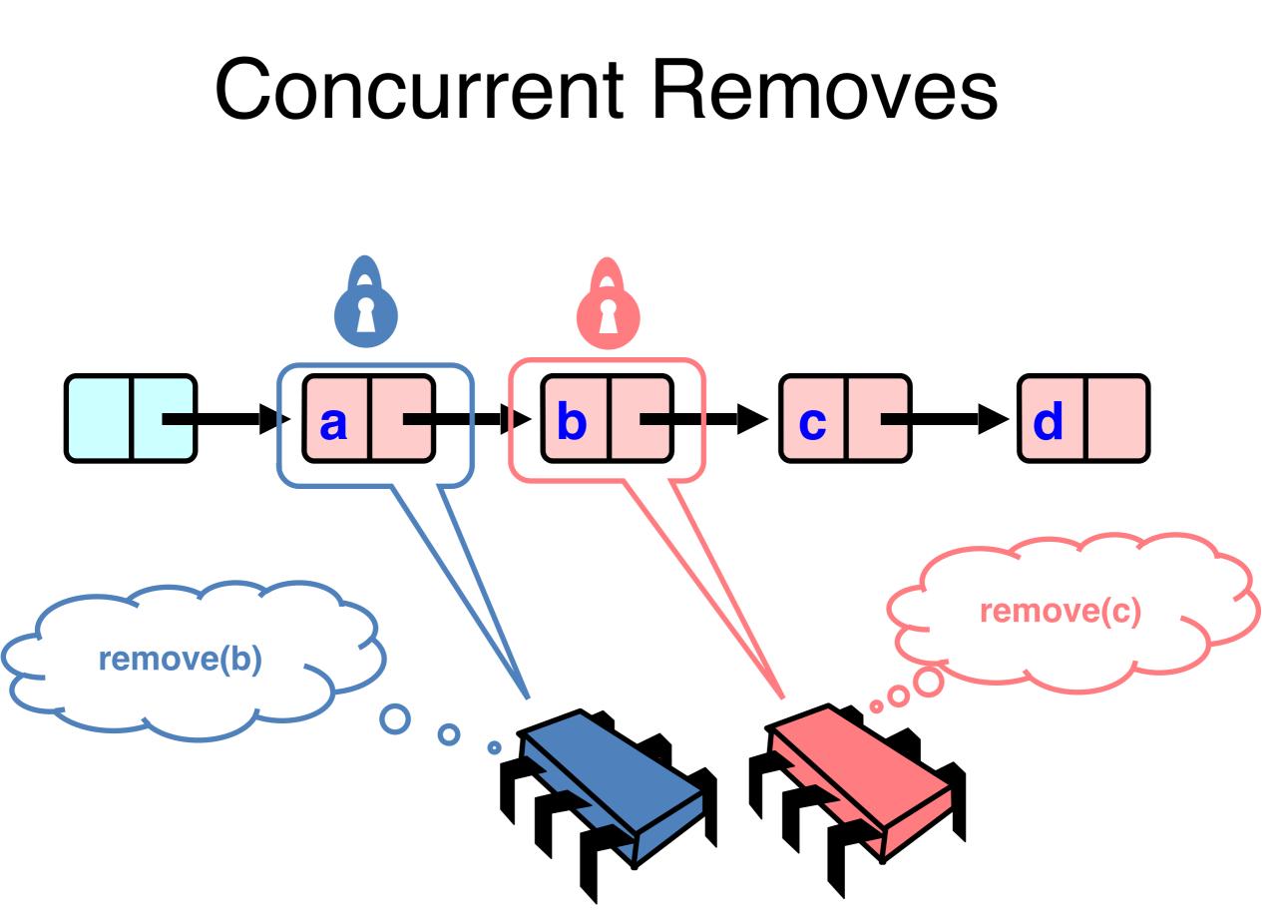


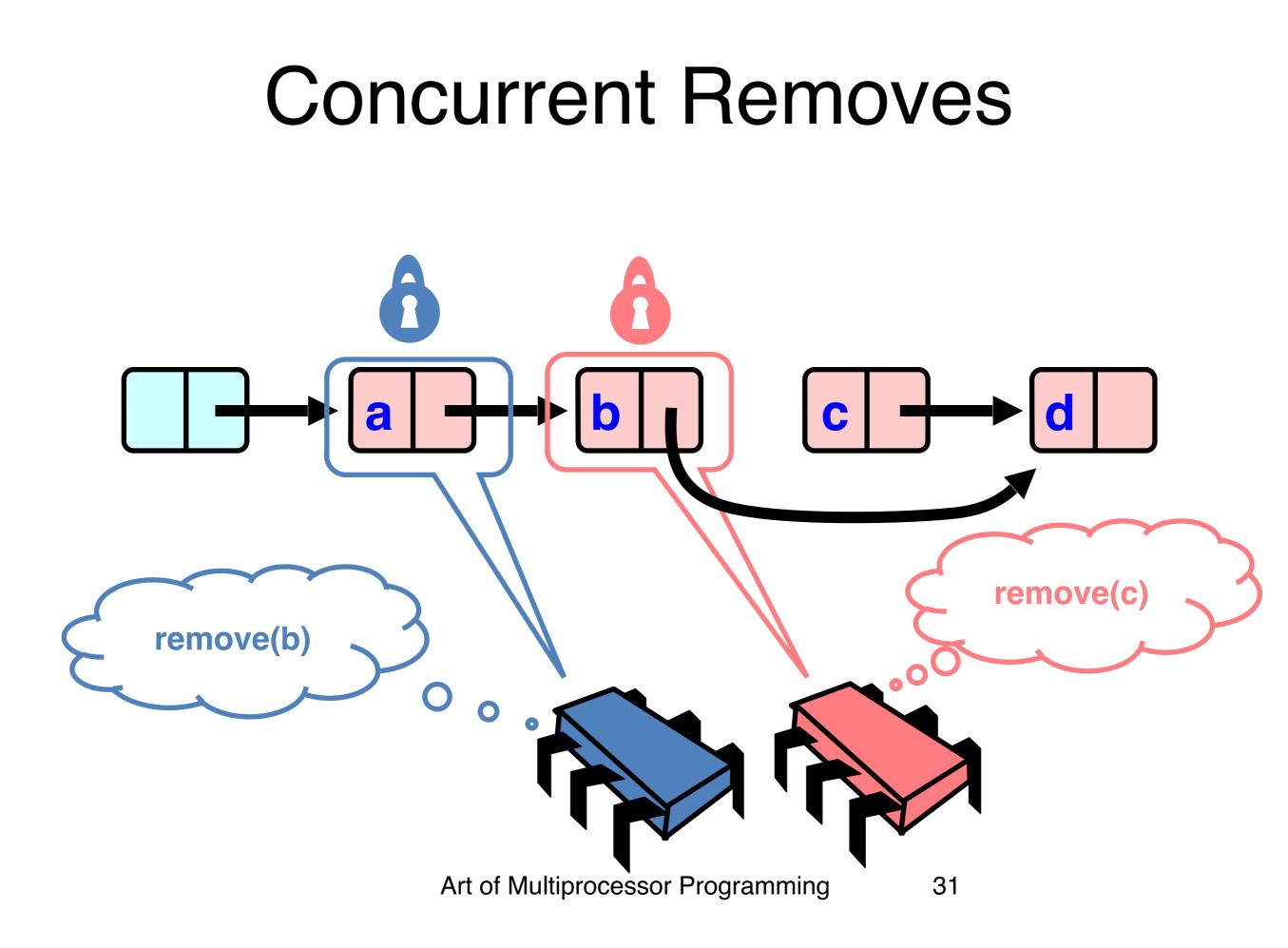


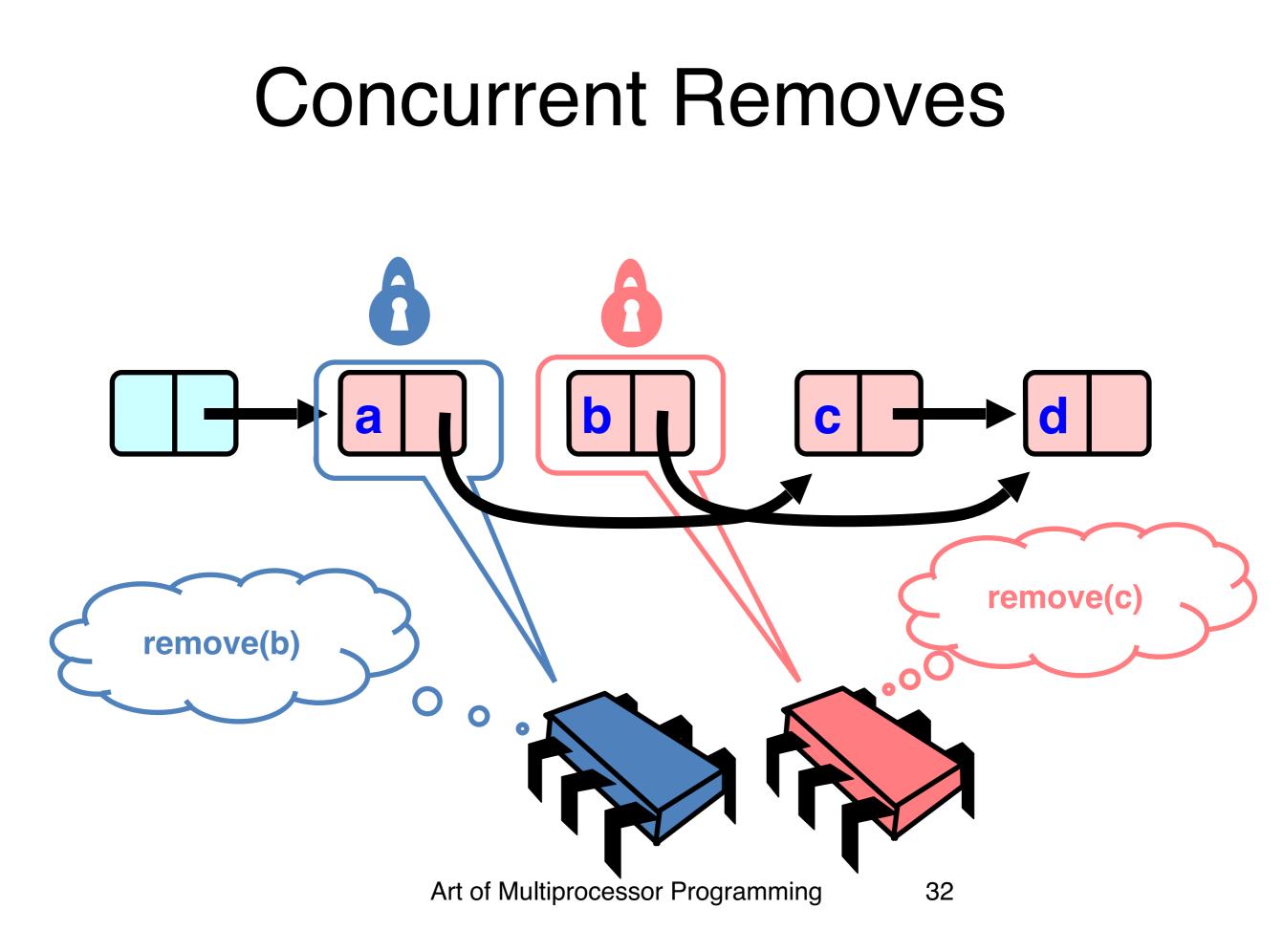


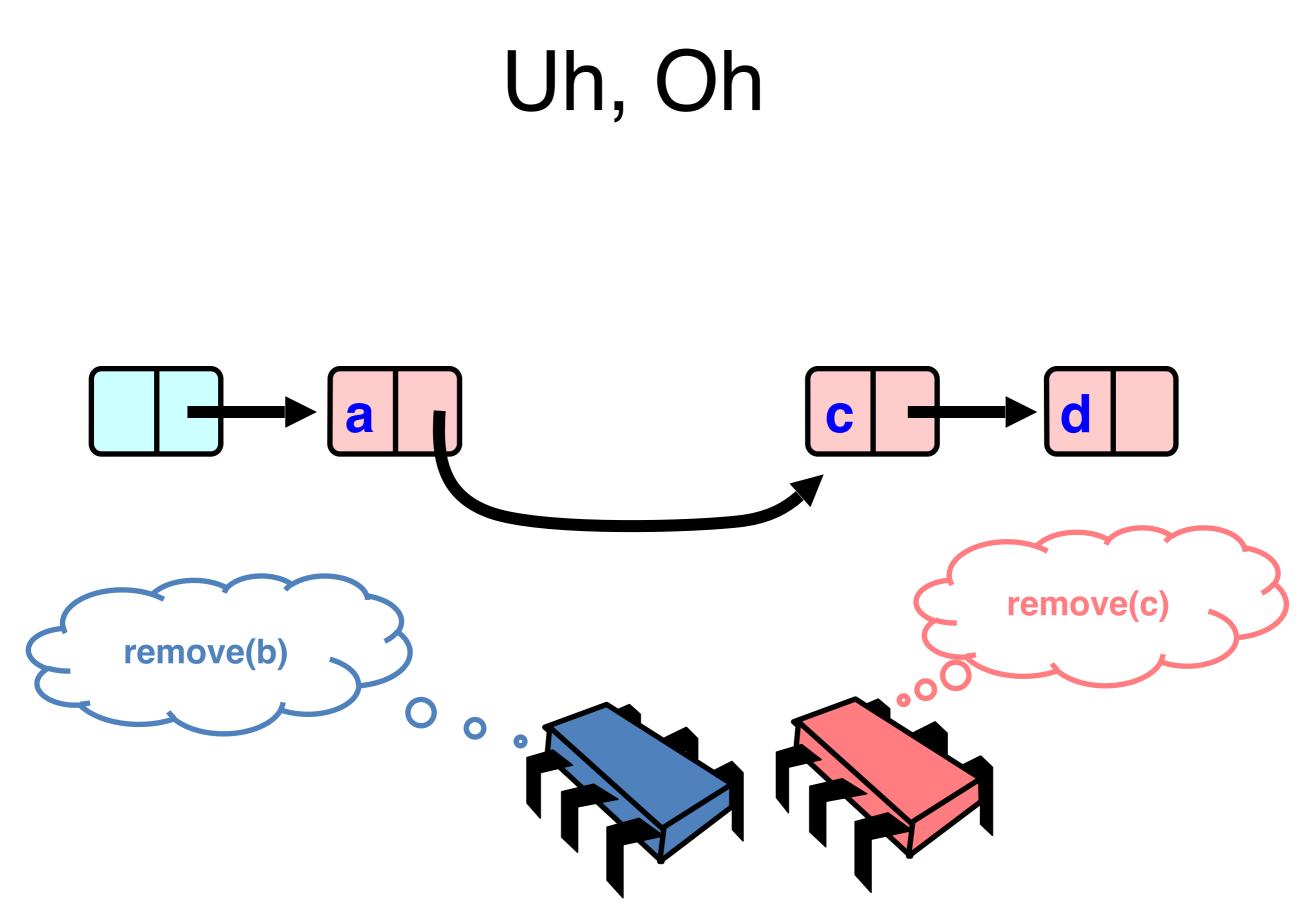




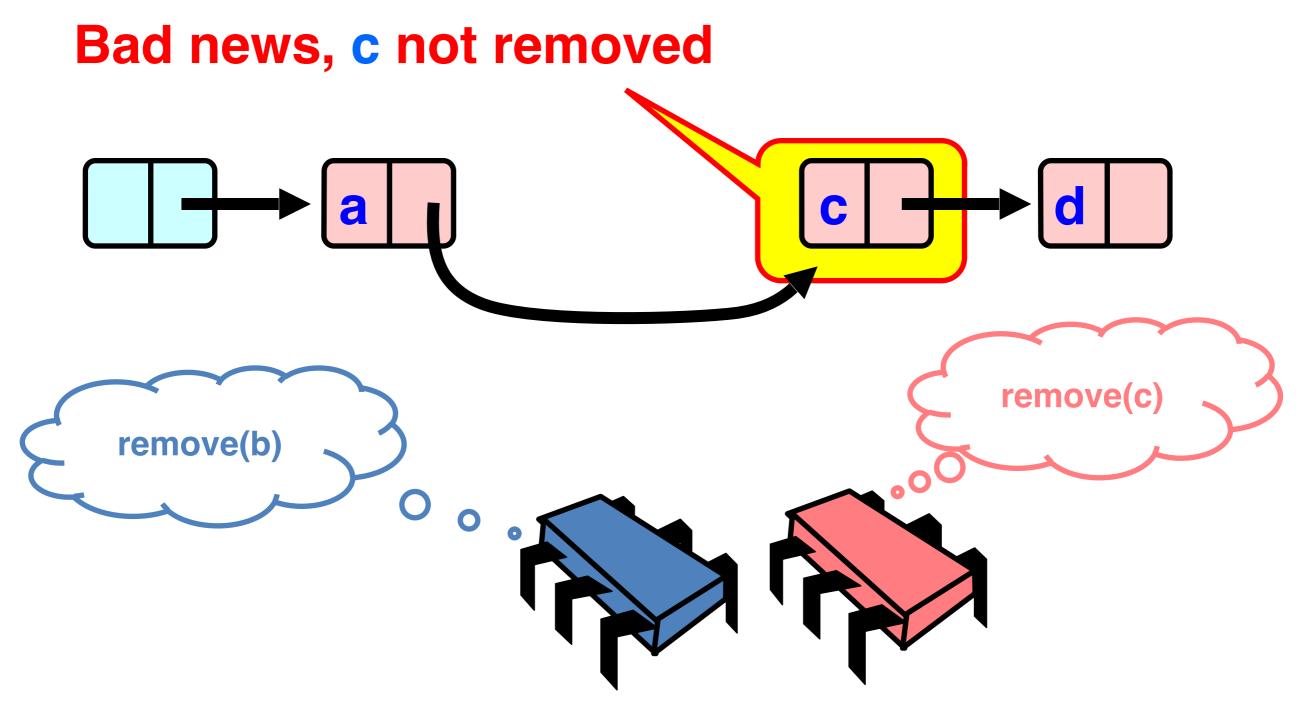






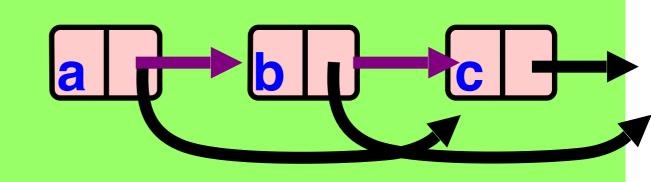


Uh, Oh



Problem

- To delete node c
 - Swing node b's next field to d
- Problem is,
 - Someone deleting b concurrently could direct a pointer



Insight

- If a node is locked
 - No one can delete node's successor
- If a thread locks
 - Node to be deleted
 - And its predecessor
 - Then it works

