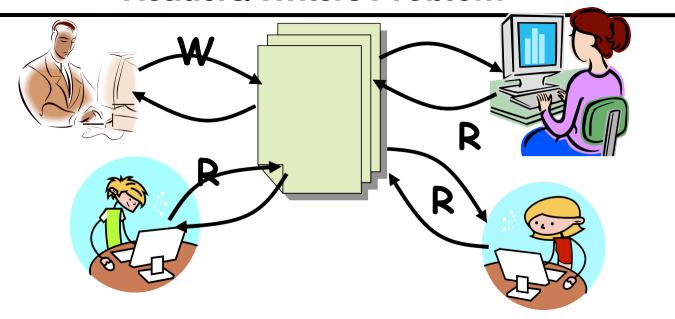
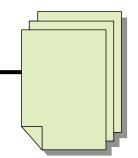
# Reader/Writer Synchronization

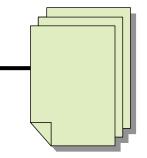
#### **Readers/Writers Problem**



- · Motivation: Consider a shared database
  - Two classes of users:
    - » Readers never modify database
    - » Writers read and modify database
  - Is using a single lock on the whole database sufficient?
    - » Like to have many readers at the same time
    - » Only one writer at a time



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  - Readers can access database when no writers
  - Writers can access database when no readers
  - Only one thread manipulates state variables at a time



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  - Reader()

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Wait until no writers
Access data base
Check out - wake up a waiting writer
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Access data base Check out - wake up a waiting writer

- Writer()

Wait until no active readers or writers Access database Check out - wake up waiting readers or writer

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  - Readers can access database when no writers
  - Writers can access database when no readers
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- Basic structure of a solution:
  - Reader()

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- Writer()

Wait until no active readers or writers Access database Check out - wake up waiting readers or writer

- State variables (Protected by a lock called "lock):
  - » int AR: Number of active readers; initially = 0
  - » int WR: Number of waiting readers; initially = 0
  - » int AW: Number of active writers; initially = 0
  - » int WW: Number of waiting writers; initially = 0
  - » Condition okToRead = NIL

```
Reader() {
   // First check self into system
  lock.Acquire();
```

```
Reader() {
  // First check self into system
  lock.Acquire();
 while ((AW + WW) > 0) { // Is it safe to read?
                          // No. Writers exist
    WR++;
    okToRead.wait(&lock); // Sleep on cond var
                          // No longer waiting
 WR--;
                          // Now we are
 AR++;
  lock.release();
                              active!
 // Perform actual read-only access
 AccessDatabase(ReadOnly);
```

```
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  // First check self into system
  lock.Acquire();
 while ((AW + WW) > 0) { // Is it safe to read?
                          // No. Writers exist
    WR++;
    okToRead.wait(&lock); // Sleep on cond var
                          // No longer waiting
 WR--;
 AR++;
                          // Now we are
 lock.release();
                              active!
 // Perform actual read-only access
 AccessDatabase(ReadOnly);
 // Now, check out of system
  lock.Acquire();
                          // No longer active
 AR--;
  if (AR == 0 && WW > 0) // No other active readers
    okToWrite.signal(); // Wake up one writer
 lock.Release();
```

```
Reader() {
  // First check self into system
  lock.Acquire();
 while ((AW + WW) > 0) { // Is it safe to read?
                           // No. Writers exist
    WR++;
    OkToRead.wait() // Sleep on cond var
    WR--;
                Why is lock realeased here?
 AR++;
  lock.release();
  // Perform actual read-only access
 AccessDatabase(ReadOnly);
  // Now, check out of system
  lock.Acquire();
                           // No longer active
 AR--;
  if (AR == 0 && WW > 0) // No other active readers
    okToWrite.signal(); // Wake up one writer
  lock.Release();
```

```
Writer() {
   // First check self into system
   lock.Acquire();
```

```
Writer() {
 // First check self into system
 lock.Acquire();
 while ((AW + AR) > 0) { // Is it safe to write?
   WW++;
                         // No. Active users exist
    okToWrite.wait(&lock); // Sleep on cond var
WW--; // No longer waiting
   WW--;
  }
                         // Now we are active!
 AW++;
 lock.release();
 // Perform actual read/write access
 AccessDatabase (ReadWrite);
 // Now, check out of system
 lock.Acquire();
                         // No longer active
 AW--:
 if (WW > 0) {
                       // Give priority to writers
 okToRead.broadcast(); // Wake all readers
 lock.Release();
```

- · Consider the following sequence of operators:
  - R1, R2, W1, R3

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First, R1 comes along:

```
AR = 1, WR = 0, AW = 0, WW = 0
```

- Consider the following sequence of operators:
  - R1, R2, W1, R3
- First, R1 comes along:

$$AR = 1$$
,  $WR = 0$ ,  $AW = 0$ ,  $WW = 0$ 

Next, R2 comes along:

$$AR = 2$$
,  $WR = 0$ ,  $AW = 0$ ,  $WW = 0$ 

Can't start because of readers, so go to sleep:
 AR = 2, WR = 0, AW = 0, WW = 1

Can't start because of readers, so go to sleep:

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AR = 2, WR = 0, AW = 0, WW = 1
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Finally, R3 comes along:
 AR = 2, WR = 1, AW = 0, WW = 1

Can't start because of readers, so go to sleep:

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```

- Finally, R3 comes along:
   AR = 2, WR = 1, AW = 0, WW = 1
- Now, say that R2 finishes before R1:
   AR = 1, WR = 1, AW = 0, WW = 1

Can't start because of readers, so go to sleep:

```
AR = 2, WR = 0, AW = 0, WW = 1
```

- Finally, R3 comes along:
   AR = 2, WR = 1, AW = 0, WW = 1
- Now, say that R2 finishes before R1:
   AR = 1, WR = 1, AW = 0, WW = 1
- Finally, last of first two readers (R1) finishes and wakes up writer:

```
• Next, W1 comes along:
    while ((AW + AR) > 0) { // Is it safe to write?
        WW++;
        okToWrite.wait(&lock); // Sleep on cond var
        WW--;
    }
    AW++;
```

Can't start because of readers, so go to sleep:

```
AR = 2, WR = 0, AW = 0, WW = 1
```

- Finally, R3 comes along:
   AR = 2, WR = 1, AW = 0, WW = 1
- Now, say that R2 finishes before R1:
   AR = 1, WR = 1, AW = 0, WW = 1
- Finally, last of first two readers (R1) finishes and wakes up writer:

```
if (AR == 0 && WW > 0)  // No other active readers
  okToWrite.signal();  // Wake up one writer
```

· When writer wakes up, get:

$$AR = 0$$
,  $WR = 1$ ,  $AW = 1$ ,  $WW = 0$ 

When writer wakes up, get:

$$AR = 0$$
,  $WR = 1$ ,  $AW = 1$ ,  $WW = 0$ 

· Then, when writer finishes:

When writer wakes up, get:

```
AR = 0, WR = 1, AW = 1, WW = 0
```

· Then, when writer finishes:

When writer wakes up, get:

```
AR = 0, WR = 1, AW = 1, WW = 0
```

· Then, when writer finishes:

- Writer wakes up reader, so get:

```
AR = 1, WR = 0, AW = 0, WW = 0
```

When writer wakes up, get:

$$AR = 0$$
,  $WR = 1$ ,  $AW = 1$ ,  $WW = 0$ 

· Then, when writer finishes:

- Writer wakes up reader, so get:

$$AR = 1$$
,  $WR = 0$ ,  $AW = 0$ ,  $WW = 0$ 

· When writer completes, we are finished