

# CMPT 383

## Quiz #8

### December 1, 2005

1. Imagine a language of expressions for representing integers defined by the syntax rules:

- (i) zero is an expression
- (ii) If  $e$  is an expression, then so are  $\text{succ}(e)$  and  $\text{pred}(e)$ .
- (iii) If  $e_1$  and  $e_2$  are expression, then so is  $\text{add}(e_1, e_2)$ .

An evaluator reduces expressions in this language by applying the following rules repeatedly until no longer possible:

$\text{succ}(\text{pred}(e)) = e$   
 $\text{pred}(\text{succ}(e)) = e$   
 $\text{add}(\text{zero}, e_2) = e_2$   
 $\text{add}(\text{succ}(e_1), e_2) = \text{succ}(\text{add}(e_1, e_2))$   
 $\text{add}(\text{pred}(e_1), e_2) = \text{pred}(\text{add}(e_1, e_2))$

Simplify the expression  $\text{add}(\text{succ}(\text{pred}(\text{zero})), \text{zero})$  using:

a) Innermost reduction sequence.

$\text{add}(\text{succ}(\text{pred}(\text{zero})), \text{zero})$   
 { by  $\text{succ}(\text{pred}(e)) = e$  }  
 $\text{add}(\text{zero}, \text{zero})$   
 { by  $\text{add}(\text{zero}, e) = e$  }  
**zero**

b) Outermost reduction sequence.

$\text{add}(\text{succ}(\text{pred}(\text{zero})), \text{zero})$   
 { by  $\text{add}(\text{succ}(e_1), e_2) = \text{succ}(\text{add}(e_1, e_2))$  }  
 $\text{succ}(\text{add}(\text{pred}(\text{zero}), \text{zero}))$   
 { by  $\text{add}(\text{pred}(e_1), e_2) = \text{pred}(\text{add}(e_1, e_2))$  }  
 $\text{succ}(\text{pred}(\text{add}(\text{zero}, \text{zero})))$   
 { by  $\text{succ}(\text{pred}(e)) = e$  }  
 $\text{add}(\text{zero}, \text{zero})$   
 {  $\text{add}(\text{zero}, e_2) = e_2$  }  
**zero**

c) Can the expression be simplified using another reduction sequence?

$\text{add}(\text{succ}(\text{pred}(\text{zero})), \text{zero})$   
 { by  $\text{add}(\text{succ}(e_1), e_2) = \text{succ}(\text{add}(e_1, e_2))$  }  
 $\text{succ}(\text{add}(\text{pred}(\text{zero}), \text{zero}))$   
 { by  $\text{add}(\text{pred}(e_1), e_2) = \text{pred}(\text{add}(e_1, e_2))$  }  
 $\text{succ}(\text{pred}(\text{add}(\text{zero}, \text{zero})))$   
 { by  $\text{add}(\text{zero}, e_2) = e_2$  }  
 $\text{succ}(\text{pred}(\text{zero}))$   
 { by  $\text{succ}(\text{pred}(e)) = e$  }  
**zero**

2. Which of the following equations are *true* for all  $xs$  and which are *false*?

- |                        |              |                               |              |
|------------------------|--------------|-------------------------------|--------------|
| a) $[]):xs = xs$       | <b>False</b> | g) $[[[]] ++ xs = [xs]$       | <b>False</b> |
| b) $[]):xs = [[], xs]$ | <b>False</b> | h) $[[[]] ++ xs = [[], xs]$   | <b>False</b> |
| c) $xs:[] = xs$        | <b>False</b> | i) $[[[]] ++ [xs] = [[], xs]$ | <b>True</b>  |
| d) $xs:[] = [xs]$      | <b>True</b>  | j) $[xs] ++ [] = [xs]$        | <b>True</b>  |
| e) $xs:xs = [xs, xs]$  | <b>False</b> | k) $[xs] ++ [xs] = [xs, xs]$  | <b>True</b>  |
| f) $[[[]] ++ xs = xs$  | <b>False</b> |                               |              |

3. Define a data type `Direction` whose values describe the four major points of the compass, and define a function `reverse` for reversing direction.

```
data Direction = North | South | East | West
  deriving (Enum)
```

```
reverse :: Direction -> Direction
reverse d = toEnum( mod (fromEnum d + 2) 4 )
```

OR

```
data Direction = North | South | East | West
  deriving (Eq)
```

```
reverse :: Direction -> Direction
reverse d
  | d == North = South
  | d == South = North
  | d == East = West
  | d == West = East
```

4. Evaluate `map (map square) [[1,2],[3,4,5]]`

```
[[1,4],[9,16,25]]
```

5. The Fibonacci numbers  $f_0, f_1, \dots$  are defined by the rule that  $f_0=0, f_1=1$  and  $f_{n+2}=f_n+f_{n+1}$  for all  $n \geq 0$ . Give a definition of the function `fib` that takes an integer  $n$  and returns  $f_n$ .

```
fib :: Integer -> Integer
fib 0 = 0
fib 1 = 1
fib (n+2) = fib n + fib (n+1)
```

6. Using pattern matching, define a function `rev2` that reverses all lists of length 2, but leaves others unchanged. Ensure that the patterns are exhaustive and disjoint.

```
rev2 :: [a] -> [a]
rev2 [] = []
rev2 (x:[]) = [x]
rev2 (x:y:[]) = [y,x]
rev2 (x:y:z:xs) = (x:y:z:xs)
```

7. Define a function `nextlet` that takes a letter of the alphabet and returns the letter coming after it. Assume that letter `'A'` follows `'Z'`.

```
nextlet :: Char -> Char
nextlet c
  | c == 'Z' = 'A'
  | c == 'z' = 'a'
  | otherwise = chr(ord c + 1)
```

8. Explain informally, but clearly and fully, what the function `assert_length`, defined below, does.

```
assert_length :: Integer -> [a] -> [a]
assert_length 0 xs = []
```

```
assert_length (n+1) xs = (head xs) : assert_length n (tail xs)
```

The function `assert_length` takes a nonnegative integer `n` and a list `xs` as arguments. This function selects an initial segment (the first `n` elements) of the given list (`xs`). An error is produced if the list (`xs`) is smaller than the number of elements required (`n`).