

Chapter 16

Logic Programming

Topics

- ◆ Proving Theorems
 - Resolution
 - Instantiation and Unification
- ◆ Prolog
 - Terms
 - Clauses
 - Inference Process
 - Backtracking

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Predicate Calculus and Proving Theorems

- ◆ A use of propositions is to discover new theorems that can be inferred from known axioms and theorems
- ◆ *Resolution*: the process of computing inferred propositions from given propositions
 - Resolution principle is similar to the idea of transitivity in algebra.

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Resolution

- ◆ Making a single inference from a pair of Horn clauses.
 - If h is the head of a Horn clause and it matches with one of the terms of another Horn clause, then that term can be replaced by h .
 - ◆ The Horn clauses:
 $h \leftarrow \text{terms}$
 $t \leftarrow t_1, h, t_2$
 - ◆ The second clause is resolved to
 $t \leftarrow t_1, \text{terms}, t_2$

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Resolution: example

- ◆ Consider the following clauses:

```
speaks(Mary,English).  
talkswith(X,Y) ← speaks(X,L), speaks(Y,L), X≠Y
```

- ◆ Resolution allow us to deduce:

```
talkswith(Mary,Y) ← speaks(Mary,English),  
speaks(Y,English), Mary≠Y
```

- ◆ Variables X and L are assigned the values "Mary" and "English" in the second rule.
- ◆ The assignment of values to values during resolution is called *instantiation*.

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Unification and Instantiation

- ◆ *Unification*: finding values for variables in propositions that allows matching process to succeed
- ◆ *Instantiation*: assigning temporary values to variables to allow unification to succeed
- ◆ After instantiating a variable with a value, if matching fails, may need to backtrack and instantiate with a different value.

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Resolution: Theorem Proving

- ◆ Use proof by contradiction.
- ◆ *Hypotheses*: a set of pertinent propositions
- ◆ *Goal*: negation of theorem stated as a proposition.
- ◆ Theorem is proved by finding an inconsistency.

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The language Prolog

- ◆ The most widely used logic programming language.
- ◆ Prolog in a nutshell
 - Uses Horn clauses
 - ◆ Almost identical notation of Horn clauses, except the implication arrow " \leftarrow " is replaced by a colon followed by a dash ":-".
 - Implements resolution using strict linear "depth first" strategy and a unification algorithm.

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A Prolog Program

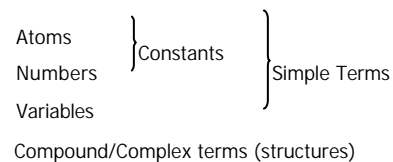
- ◆ Prolog approach.
 - Describe known facts and relationships.
- ◆ A program consists of a sequence of Horn clauses.
 - Clauses are implications of the form
 - ◆ p_1 if $p_2 \dots \&p_k$ written as $p_1 :- p_2, \dots, p_k$ where every p is a term.
 - p_i is called the *head* and p_2, \dots, p_k is the *body*.
 - The body is a list of *goals* separated by commas (conjunctions).

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Terms

- ◆ All Prolog statements are constructed from terms.



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Terms: atoms

- ◆ Atoms are used as names of individuals and predicates.
- ◆ Atoms can be constructed in three ways:
 1. Strings of letters, digits and the underscore character, starting with a lower-case letter.
 - ◆ x
 - ◆ vancouver
 - ◆ ax123aBCD
 - ◆ abc_123_etc

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Terms: atoms

2. Strings of special characters .
 - ◆ Some of these have a predefined meaning (:-).
 - ◆ -->
 - ◆ ==>
3. String of characters enclosed in single quotes.
 - ◆ It can contain any character.
 - ◆ 'Florida'
 - ◆ '12\$12\$'
 - ◆ ' back\\slashes'
 - ◆ '32' is an atom, not equal to the number 32.
 - ◆ A zero length atom is written ''

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Terms: atoms

◆ Internal representation:

- Atoms are not character strings, they are locations in a symbol table.
- Only one copy of each atom is stored.
- All occurrences are replaced by pointers to its location in the symbol table.
 - ◆ `abracadabraabracazam = abracadabraabracazam`
 - ◆ `a = b`
- ◆ Comparison takes the same amount of time.
 - Pointers get compared rather than the strings of characters.

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Terms: numbers

◆ Integer

- Useful for task such as counting the elements of a list.

◆ Real

- Not used very much in typical Prolog programming.

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Terms: variables

- ◆ Stand for objects that we cannot name.
- ◆ Begins with a capital letter or `'_'`.
 - `X`
 - `_value`
 - `Mother`
- ◆ A variable can be
 - *Instantiated* when there is an object that it stands for.
 - *Uninstantiated* when what the variable it stands for is not yet known.

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Terms: variables

◆ Question containing a variable:

- Search through all facts to find an object that the variable could stand for.

◆ Anonymous variable (`_`)

- A special variable that matches anything, but never takes on a value.
- Successive anonymous variables in the same clause do not take on the same value.

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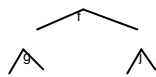
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Terms: structures

◆ Consists of a *functor* followed by a sequence of *arguments*.

- The functor must be an atom.
- Arguments can be any kind of terms (including other structures).
- The outermost functor (`f/2`, in this case) is called the *principal functor* of the structure.

◆ `f(g(h,i),j(k,l))`



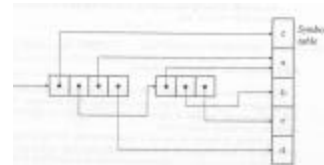
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Terms: structures

◆ Internal representation

- A structure is a linked tree made of pointers to its substructures and to entries in the symbol table.



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Prolog Clauses

- ◆ Prolog clauses are of three types:
 - *Facts*: declare things that are always, unconditionally true.
 - *Rules*: declare things that are true, depending on a given condition.
 - By means of *questions* users can ask the program what things are true.

Prolog Clauses: characteristics

- ◆ A Prolog clause:
 - Arbitrary number of arguments (parameters).
 - A predicate that takes N arguments is called *N-placed predicate*.
 - A one-place predicate describes a *property* of one individual; a two-place predicate describes a *relation* between two individuals.
 - The number of arguments that a predicate takes is called its *arity*.

Prolog Clauses: characteristics

- Two distinct predicates can have the same name if they have different arities.
 - ◆ `mother(pam)` meaning Pam is a mother.
 - ◆ `mother(pam,bob)` meaning Pam is the mother of Bob.
- A predicate is identified by giving its name, a slash, and its arity.
 - ◆ `mother/1`.
 - ◆ `mother/2`.
- Every Prolog statement is terminated by a period.

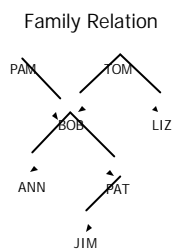
Facts

- ◆ Define relationships between objects.
- ◆ Sometimes called *ground clauses* because they are the basis from which other information is inferred.
- ◆ Facts are clauses that have an empty body.
- ◆ Facts are written as:


```
name_of_relationship(object,...,object).
```
- ◆ A collection of facts and rules is called a *database* or *knowledge base*.

Facts: examples

```
parent(tom,bob).
parent(pam,bob).
parent(tom,liz).
parent(bob,ann).
parent(bob,pat).
parent(pat,jim).
```



Rules

- ◆ Specify things that are true if some condition is satisfied.
- ◆ Rules are used to say that a fact depends on a group of other facts.
- ◆ A rule consists of a head and a body.
 - Connected by the symbol `:-` (if).
 - The head describes what fact the rule is intended to define.
 - The body describes the conjunctions of goals that must be satisfied for the head to be true.

Rules: examples

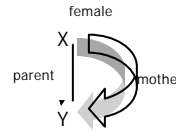
For all X and Y,
Y is an offspring of X if
X is a parent of Y.

$$?- \underbrace{\text{offspring}(Y, X)}_{\text{head/conclusion}} : - \underbrace{\text{parent}(X, Y)}_{\text{body/condition}} .$$

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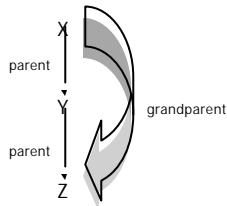
Rules: examples

$$?- \text{mother}(X, Y) :- \text{parent}(X, Y), \text{female}(X)$$


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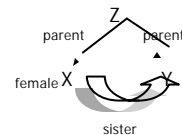
Rules: examples

$$?- \text{grandparent}(X, Z) :- \text{parent}(X, Y), \text{parent}(Y, Z) .$$


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Rules: examples

$$?- \text{sister}(X, Y) :- \text{parent}(Z, X), \text{parent}(Z, Y), \text{female}(X) .$$


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Rules: exercise

◆ Define the relation `aunt(X, Y)`

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Questions

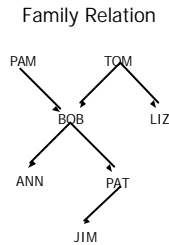
- ◆ How do we use Prolog programs?
 - By posing queries (a query is a request for the computer to do something.)
- ◆ Questions are clauses that only have the body.
 - Search through the database.
 - Look for facts that *match* the fact in question.
 - ◆ Two facts *match* if their predicates and corresponding arguments are the same.
 - YES: Prolog finds a match
 - No: Prolog does not find a match (nothing matches the question vs. false)

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Questions: example

User Queries	Prolog's Answers
?-parent(bob,pat).	yes
?-parent(liz,pat).	no
?-parent(tom,ben).	no
?-parent(X,liz).	X=tom
?-parent(bob,X).	X=ann X=pat
?-parent(X,Y).	X=pam Y=bob; X=tom Y=bob;



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Conjunctions

- ◆ More complicated relationships.
- ◆ Satisfy two or more separate goals.
- ◆ Commas are understood as conjunctions.
- ◆ Question containing conjunctions:
 - Try to satisfy each goal in turn (searching the database).
 - All goals have to be satisfied.

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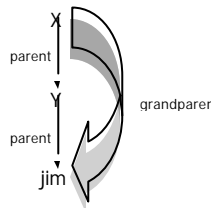
Complex Queries: example

Who is a grandfather of Jim?

1. Who is a parent of Jim?
Assume that this is some Y.
2. Who is a parent of Y?
Assume that this is some X.

?-parent(X,Y),parent(Y,jim).

X=bob
Y=pat



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Complex Queries: example

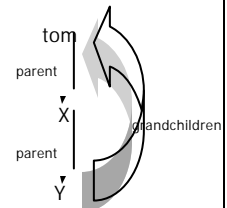
Who are Tom's grandchildren?

1. Who is a child of Tom?
Assume that this is some X.
2. Who is a child of X?
Assume that this is some Y.

?-parent(tom,X),parent(X,Y).

X=bob
Y=ann

X=bob
Y=pat



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Complex Queries: example

Do Ann and Pat have a common parent?

1. Who is a parent, X, of Ann?
2. Is (this same) X a parent of Pat?

?-parent(X,ann),parent(X,pat).

X=bob

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Complex Queries: exercises

What will be Prolog's answers to the following questions?

1. ?-parent(jim,X).
2. ?-parent(X,jim).
3. ?-parent(pam,X),parent(X,pat).
4. ?-parent(pam,X),parent(X,Y),parent(Y,jim).

Formulate in Prolog the following questions:

1. Who is Pat's parent?
2. Does liz have a child?
3. Who is Pat's grandparent?

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Matching / Unification

- Two terms ($term_1$ and $term_2$) can be unified (matched) if they are alike or can be made alike by instantiation.
 - Instantiation: make one variable the same as another.

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Matching / Unification

- If $term_1$ and $term_2$ are constants
 - Same atom or same number.
- If $term_1$ is a variable and $term_2$ is any type of term
 - $term_1$ is instantiated to $term_2$.
- If $term_1$ and $term_2$ are variables
 - They are instantiated to each other (share values).
- If $term_1$ and $term_2$ are complex terms
 - They have the same functor and arity.
 - All their corresponding arguments match.
 - The variable instantiations are compatible.
- Two terms match if and only if it follows from the previous four clauses that they match.

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Unification: examples

User Queries	Prolog's Answers
?-2=2.	yes
?-bob=jim.	no
?-'bob'='bob'.	yes
?-'2'='2'.	no
?-bob=X.	X=bob
?-X=bob.	yes
?-X=Y.	yes
?-X=bob, X=jim.	no
?-kill(shoot(gun),Y)=kill(X,stab(knife)).	X=shoot(gun) Y=stab(knife)
?-kill(shoot(gun),stab(knife))=kill(X,stab(Y)).	yes
?-kill(shoot(gun),stab(knife))=kill(X,stab(Y)).	X=shoot(gun) Y=knife
	yes

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Inference Process of Prolog

- Queries are called *goals*
- If a goal is a compound proposition, each of the facts is a *subgoal*
- To prove a goal is true, must find a chain of inference rules and/or facts. For goal Q:
 - B :- A
 - C :- B
 - ...
 - Q :- P

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Inference Process of Prolog

- Process of proving a subgoal is called *matching*, *satisfying*, or *resolution*.
- Bottom-up resolution, forward chaining
 - Begin with facts and rules of database and attempt to find sequence that leads to goal
 - Works well with a large set of possibly correct answers
- Top-down resolution, backward chaining
 - Begin with goal and attempt to find sequence that leads to set of facts in database.

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Inference Process of Prolog

- works well with a small set of possibly correct answers
- Prolog implementations use backward chaining
- When goal has more than one subgoal, can use either
 - Depth-first search: find a complete proof for the first subgoal before working on others
 - Breadth-first search: work on all subgoals in parallel.
- Prolog uses depth-first search

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A Simple Prolog Knowledge Base

- ◆ A Prolog knowledge base that describes the location of certain North American cities.

```

/* 1 */ located_in(atlanta,georgia).
/* 2 */ located_in(houston,texas).
/* 3 */ located_in(austin,texas).
/* 4 */ located_in(toronto,ontario).
/* 5 */ located_in(X,usa) :- located_in(X,georgia).
/* 6 */ located_in(X,usa) :- located_in(X,texas).
/* 7 */ located_in(X,canada) :- located_in(X,ontario).
/* 8 */ located_in(X,north_america) :-
    located_in(X,usa).
/* 9 */ located_in(X,north_america) :-
    located_in(X,canada).

```

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Inference Process: example

- ◆ Query:

```
?- located_in(austin,north_america).
```

Unifies with the head of Clause 8 by instantiating *x* as *austin*.

The right-hand side of Clause 8 becomes the new goal.

Goal: `?- located_in(austin,north_america).`

Clause 8: `located_in(X,north_america) :-`

`located_in(X,usa).`

Instantiation: *X* = *austin*

New goal: `?-located_in(austin,usa).`

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Inference Process: example

Unify the new query with Clause 6:

Goal: `?- located_in(austin,usa).`

Clause 6: `located_in(X,usa) :-`
`located_in(X,texas).`

Instantiation: *X* = *austin*

New goal: `?-located_in(austin,texas).`

This query matches Clause 3. Since Clause 3 does not contain an "if", no new query is generated and the process terminates successfully.

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Backtracking

- ◆ If several rules can unify with a query, how does Prolog know which one to use?
- ◆ Prolog does not know in advance which clause will succeed but it does know how to black out of blind alleys.
- ◆ Prolog tries the rules in order in which they are given in the knowledge base.
 - If a rule does not lead to success, it backs up and tries another

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Backtracking

- ◆ The query `?-located_in(austin,usa)` will try to unify with Clause 5 and then, when that fails, the computer will back up and try Clause 6.
- ◆ A good way to conceive of backtracking is to arrange all possible paths of computation into a tree.
- ◆ Consider the query:
`?- located_in(toronto,north_america).`

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Backtracking: example

- ◆ The following tree shows all the paths that the computation might follow.

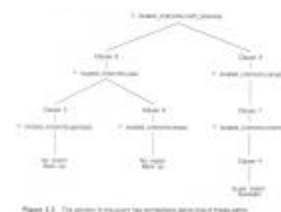


Figure 1.1. The search tree for the query `?-located_in(toronto,north_america).`

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Backtracking: example

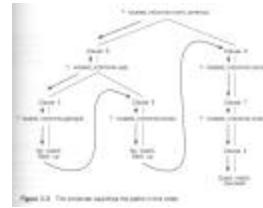
- ◆ We can prove that Toronto is in North America if we can prove that it is in either the U.S.A. or Canada.
- ◆ If we try the U.S.A., we have to try several states and then Canada.
- ◆ Almost all paths are blind alleys.
- ◆ Only the rightmost one leads to a successful solution.

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Backtracking: example

- ◆ Same diagram with arrows added to show the order in which the possibilities are tried.



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Backtracking: example

- ◆ Whenever the computer finds it has gone down a blind alley, it backs up to the most recent query for which there are still untried alternatives, and tries another path.
- ◆ When a successful answer is found, the process stops.

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Prolog Syntax: comments

- ◆ Two ways to delimit comments
 - Anything bracketed by `/*` and `*/`
 - `/* This is a comment */`
- ◆ Anything between `%` and the end of the line
 - `% This is also a comment`

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Sicstus

- ◆ <http://www.cs.sfu.ca/CC/SW/Prolog/>
- ◆ Linux and SunOS machines
 - CSIL SunOS machines have an additional Prolog implementation: BinProlog.
- ◆ Running Sicstus Prolog
 - `orion% sicstus`

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Running Sicstus

- ◆ Interactive definition mode:
 - The special goal `[user]` is used to enter interactive definition mode.
 - In interactive mode, Prolog expects goals to establish.
 - `^D` (i.e., the ctrl-D key) exits definition mode.
- ◆ Consulting a file
 - The special query
 - `?-consult('family.pl').`asks prolog to read the definitions from the named file in quotes.
 - Goals must be terminated with a period or Prolog just waits until you enter the period.

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Running Sicstus

- To load the same program use `reconsult` instead of `consult`.
 - Otherwise, there will be two copies of it in memory at the same time.
- ◆ To exit from Prolog just type the special query `?- halt.`
- ◆ If a single query has multiple solutions, Prolog finds one solution and then asks whether to look for another (until all alternatives are found or you stop asking for them).

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Running Sicstus

```
?- located_in(X,texas).  
X = houston  
More (y/n) ? y  
X = austin  
More (y/n) ? Y  
no
```

- The “no” at the end means there are no more solutions.
- ◆ Any of the arguments of a predicate can be queried.
 - `?- located_in(austin,X).` % Names of regions that contain Austin
 - `?- located_in(X,texas).` % Names of cities that are in Texas
 - `?- located_in(X,Y).` % “What is in what?”
 - `?- located_in(X,X).` % “What is in itself?”

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Sicstus: examples

◆ Sample Prolog session

```
1: orion% sicstus  
SICStus 3.8.4 (sparc-solaris-5.6): Mon Jun 12 18:49:23 MET DST 2000  
Licensed to cs.sfu.ca  
| ?- [user].  
| member(X, [X, _]).  
| member(X, [_|More]) :- member(X, More).  
| ^D  
| (consulted user in module user, 0 msec 336 bytes)  
| ?- [data.pl].  
| (consulting /cs/gard1/dma/family.pl...)  
| (consulted /cs/gard1/dma/family.pl in module user, 10 msec 160 bytes)  
yes  
| ?- parent(X,liz).  
X = tom;  
no  
| ?- halt.  
2: orion%
```

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