

## ***Integrity Constraints – Practice Questions Solution***

1. Consider the schema:

*employee(employee-name, street, city)*  
*works(employee-name, company-name, salary)*  
*company(company-name, city)*  
*manages(employee-name, manager-name)*

Give an SQL DDL definition for the tables of this database. Identify referential-integrity constraints that should hold and include them in the DDL definition.

```
create table employee
(person-name char(20) not null,
street char(30),
city char(30),
primary key(person-name))
```

```
create table works
(person-name char(20) not null,
company-name char(15) not null,
salary integer,
foreign key (person-name) references employee,
foreign key (company-name) references company)
```

```
create table company
(company-name char(15) not null,
city char(30))
```

```
create table manages
(person-name char(20) not null,
manager-name char(20) not null,
foreign key (person-name) references employee,
foreign key (manager-name) references employee)
```

Note that your domain selections can be different, as long as they are compatible where one column is a foreign key of another.

2. Given the following *relation*, list all of the nontrivial functional dependencies satisfied in it:

A	B	C
a <sub>1</sub>	b <sub>1</sub>	c <sub>1</sub>
a <sub>1</sub>	b <sub>1</sub>	c <sub>2</sub>
a <sub>2</sub>	b <sub>1</sub>	c <sub>1</sub>
a <sub>2</sub>	b <sub>1</sub>	c <sub>3</sub>

The nontrivial functional dependencies are:

A → B  
 C → B  
 AC → B

Note that the last functional dependency is actually logically implied by the first two (which axiom(s)?). There are also 19 trivial functional dependencies of the form  $\alpha \rightarrow \beta$ , where  $\beta \subseteq \alpha$ .

3. Given the relation schema  $R = (A, B, C, D, E)$  and the set of functional dependencies:

F = {  
     E → AB  
     BC → D  
     D → E  
     AB → BC  
     BC → E  
 } }

Compute the canonical cover  $F_c$ . Show your steps clearly!

Using the algorithm for computing a canonical cover in text 6.5.4:

1. Use the union rule to replace

BC → D  
 BC → E

With

BC → DE

The left side of each functional dependency in F is now unique, so there are no more functional dependencies to replace using the union

rule.

2. The attribute B in BC of  $AB \rightarrow BC$  is extraneous because from the algorithm from page 209 of the text,  $AB \rightarrow C$  logically implies  $AB \rightarrow BC$ , so replace  $AB \rightarrow BC$  with  $AB \rightarrow C$ .
3. The attribute E in  $BC \rightarrow DE$  is extraneous because  $E \in DE$  and  $(F - \{BC \rightarrow DE\}) \cup \{BC \rightarrow (DE - E)\}$  logically implies F.

This is true because  $BC \rightarrow D$  is one of the given functional dependencies, so replace  $BC \rightarrow DE$  with  $BC \rightarrow D$ .

4. There are no more extraneous attributes, since none of the attributes on the left side or right side of any remaining functional dependency is extraneous. Therefore, the canonical cover is:

$$F_c = \left\{ \begin{array}{l} E \rightarrow AB \\ BC \rightarrow D \\ D \rightarrow E \\ AB \rightarrow C \end{array} \right\}$$