Insertion Sort

Lecture 10

Today

• Insertion Sort

Insertion Sort Algorithm

Strategy:

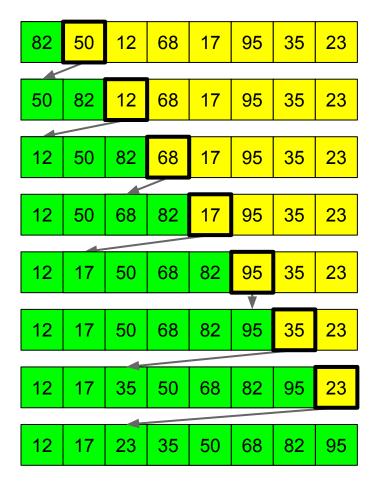
- Insert one element at a time into a sorted list
 - Locate the insertion point
 - Slide array elements to make space while new element < array element



- Array divided into two parts: sorted and unsorted (like Selection Sort)
- Sorted part grows one at a time (like Selection Sort)

Insertion Sort Demo

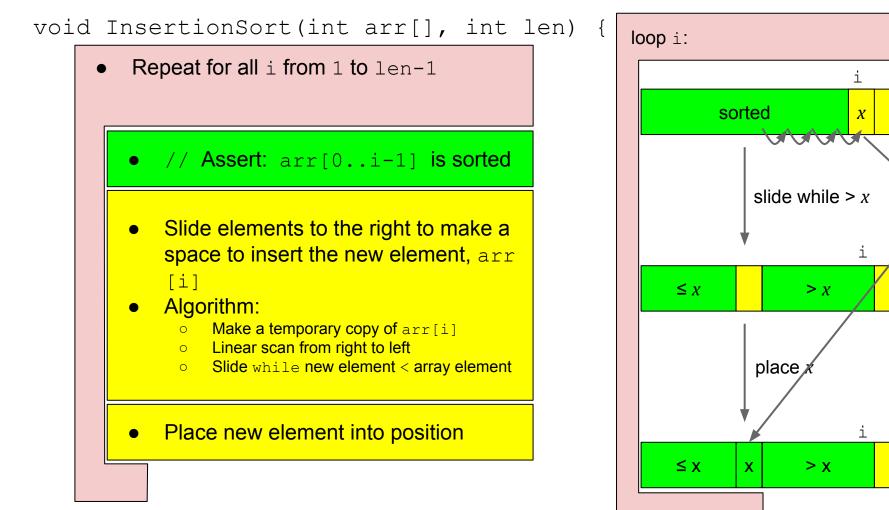
Sort this array using Insertion Sort:



create insertion point in	0 slides	Total number of slides depends on the initial order of
create insertion point in	1 slide	the input.
create insertion point in	2 slides	
create insertion point in	1 slide	What's the worst case for array of length <i>N</i> ?
create insertion point in	3 slides	
create insertion point in	0 slides	What's the best
create insertion point in	4 slides	case?

create insertion point in 5 slides

Insertion Sort in Pseudocode + Assertion Analysis



X

Analysis of Insertion Sort

What's the worst case behaviour on an array of length *N*?

OR . . .

What's the barometer instruction?

Inner loop could be executed i times

• i slides per loop $\Rightarrow O(N^2)$ total slides (in the worst case)

What sort of input leads to the worst case?

• when input array is reverse sorted

Analysis of Insertion Sort

What's the *best* case?

- When the input array is sorted
- Inner loop executed 0 times \Rightarrow 0 slides

Does this mean a running time of O(0)?

- while condition is entry condition (always performed at least once)
- So, O(N) comparisons in the best case
 - to verify the array is indeed sorted

Conclusions

- Insertion Sort algorithm varies greatly with nature of input
 - Worst case $O(N^2)$ vastly differs from best case O(N)
 - Which case carries more meaning?
- Selection Sort vs Insertion Sort
 - are incremental sorts
 - have same asymptotic running times
- Best sorting algorithms run in *O*(*N*log*N*)
 - New paradigm: Divide & Conquer