

# Merge Sort

## Basics

- Merge sort is usually described recursively.
  - The recursive calls work on smaller and smaller part of the list.
- Idea:
  1. Split the list into two halves.
  2. Recursively sort each half.
  3. “Merge” the two halves into a single sorted list.

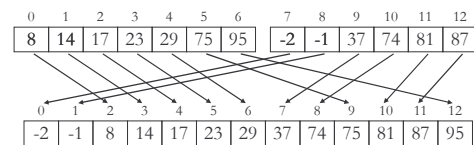
## Example

- |    |   |    |    |    |    |    |    |    |    |    |    |    |
|----|---|----|----|----|----|----|----|----|----|----|----|----|
| 0  | 1 | 2  | 3  | 4  | 5  | 6  | 7  | 8  | 9  | 10 | 11 | 12 |
| 17 | 8 | 75 | 23 | 14 | 95 | 29 | 87 | 74 | -1 | -2 | 37 | 81 |
1. Split:

0	1	2	3	4	5	6	7	8	9	10	11	12
17	8	75	23	14	95	29	87	74	-1	-2	37	81
  2. Recursively sort:

0	1	2	3	4	5	6	7	8	9	10	11	12
8	14	17	23	29	75	95	-2	-1	37	74	81	87
  3. Merge?

## Example Merge



- Must put the next-smallest element into the merged list at each step.
- each “next-smallest” could come from either half

## Merge Sort Algorithm

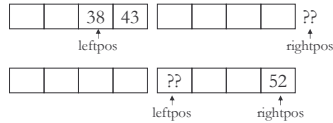
```
mergeSort(array, first, last):  
    // sort array[first] to array[last-1]  
    if last - first ≤ 1:  
        return // arrays of length 0, 1 are already sorted  
    mid = (first + last)/2  
    mergeSort(array, first, mid) // recursive call 1  
    mergeSort(array, mid, last) // recursive call 2  
    merge(array, first, mid, last)
```

## Merge Algorithm (incorrect)

```
merge(array, first, mid, last):  
    // merge array[first to mid-1] and array[mid to last-1]  
    leftpos = first  
    rightpos = mid  
    for newpos from 0 to last-first:  
        if array[leftpos] ≤ array[rightpos]:  
            newarray[newpos] = array[leftpos]  
            leftpos++  
        else:  
            newarray[newpos] = array[rightpos]  
            rightpos++  
    copy newarray to array[first to last-1]
```

## Problem?

- This algorithm starts correctly, but has an error as it finishes.
  - Eventually, one of the halves will empty.
- Then, the “if” will compare against ???
  - the element past the end of one of the halves
- one of:



## Solution

- Must prevent this: we can only look at the correct parts of the array.
- So, compare only until we reach the end of one half.
  - Then, just copy the rest over.

## Corrected Merge Algorithm (1)

```
merge(array, first, mid, last):
    leftpos = first
    rightpos = mid
    newpos = 0
    while leftpos < mid and rightpos <= last-1:
        if array[leftpos] <= array[rightpos]:
            newarray[newpos] = array[leftpos]
            leftpos++; newpos++
        else:
            newarray[newpos] = array[rightpos]
            rightpos++; newpos++
```

(continues)

## Corrected Merge Algorithm (2)

```
merge(array, first, mid, last):
    ... // code from last slide

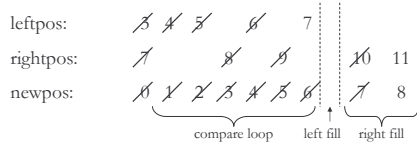
    while leftpos < mid: // copy rest of left half (if any)
        newarray[newpos] = array[leftpos]
        leftpos++; newpos++
    while rightpos <= last-1: // copy rest of right half (if any)
        newarray[newpos] = array[rightpos]
        rightpos++; newpos++

    copy newarray to array[first to (last-1)]
```

## Example 1

merge(array, 3, 7, 11):

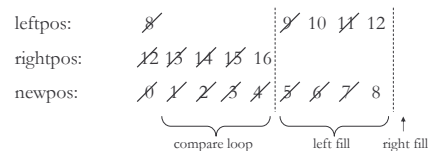
array:	3	4	5	6	7	8	9	10
	10	20	40	50	30	40	60	70
newarray:	0	1	2	3	4	5	6	7
	10	20	30	40	40	50	60	70



## Example 2

merge(array, 8, 12, 15):

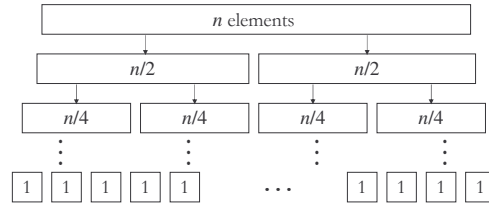
array:	8	9	10	11	12	13	14	15
	50	60	70	80	10	20	30	40
newarray:	0	1	2	3	4	5	6	7
	10	20	30	40	50	60	70	80



## Running Time

- What is merge sort's running time?
- recursive calls  $\times$  work per call?
  - yes, but overly-simplified.
  - Work per call changes.
- We know: the merge algorithm takes  $O(m)$  work to merge a total of  $m$  elements.

## Merge Sort Recursive Calls



- Each level has a total of  $n$  elements.

## Running Time

- $O(n)$  total time to merge each level
- $O(\log n)$  levels
- Total time for merge sort:  $O(n \cdot \log n)$ 
  - Much faster than insertion sort, which takes  $O(n^2)$ .
- In general, no sorting algorithm can do better than  $O(n \cdot \log n)$ .
  - There are some algorithms that are faster for limited cases.

## In-place Sorting

- Merging requires extra storage.
  - an extra array with  $n$  elements (can be reused by all merges)
  - Insertion sort requires no extra storage (except a few numeric variables).
- An algorithm that uses at most  $O(1)$  extra storage is called “in-place”.
  - Insertion sort is in-place; merge sort isn't.

## Stable Sorting

- In merge sort, equal elements are kept in-order.
  - Think of sorting a spreadsheet with other columns.
  - Rows with equal values in the sort column stay in order.
- A “stable” sorting algorithm has this property.
  - Our implementations of insertion sort & merge sort are both stable.