LOGICAL AGENTS

CHAPTER 7
Outline

♦ Knowledge-based agents
♦ Wumpus world
♦ Logic in general—models and entailment
♦ Propositional (Boolean) logic
♦ Equivalence, validity, satisfiability
♦ Inference rules and theorem proving
  – forward chaining
  – backward chaining
  – resolution
Knowledge bases

Knowledge base = set of sentences in a formal language

Declarative approach to building an agent (or other system):
   **Tell** it what it needs to know

Then it can **Ask** itself what to do—answers should follow from the KB

Agents can be viewed at the knowledge level
   i.e., **what they know**, regardless of how implemented

Or at the implementation level
   i.e., data structures in KB and algorithms that manipulate them
A simple knowledge-based agent

function KB-Agent(\textit{percept}) returns an \textit{action}

static: \textit{KB}, a knowledge base
\textit{t}, a counter, initially 0, indicating time

\textit{Tell}(\textit{KB}, \text{Make-Percept-Sentence}(\textit{percept}, \textit{t}))
\textit{action} \leftarrow \text{Ask}(\textit{KB}, \text{Make-Action-Query}(\textit{t}))
\textit{Tell}(\textit{KB}, \text{Make-Action-Sentence}(\textit{action}, \textit{t}))
\textit{t} \leftarrow \textit{t} + 1
return \textit{action}

The agent must be able to:
- Represent states, actions, etc.
- Incorporate new percepts
- Update internal representations of the world
- Deduce hidden properties of the world
- Deduce appropriate actions
Wumpus World PEAS description

Performance measure
gold +1000, death -1000
-1 per step, -10 for using the arrow

Environment
Squares adjacent to wumpus are smelly
Squares adjacent to pit are breezy
Glitter iff gold is in the same square
Shooting kills wumpus if you are facing it
Shooting uses up the only arrow
Grabbing picks up gold if in same square
Releasing drops the gold in same square

Actuators Left turn, Right turn,
Forward, Grab, Release, Shoot

Sensors Breeze, Glitter, Smell
Wumpus world characterization

Observable??
Wumpus world characterization

Observable?? No—only local perception

Deterministic??
Wumpus world characterization

**Observable??** No—only **local** perception

**Deterministic??** Yes—outcomes exactly specified

**Episodic??**
Wumpus world characterization

**Observable**  No—only local perception

**Deterministic**  Yes—outcomes exactly specified

**Episodic**  No—sequential at the level of actions

**Static**
Wumpus world characterization

**Observable??** No—only local perception

**Deterministic??** Yes—outcomes exactly specified

**Episodic??** No—sequential at the level of actions

**Static??** Yes—Wumpus and Pits do not move

**Discrete??**
### Wumpus world characterization

<table>
<thead>
<tr>
<th>Characterization</th>
<th>Description</th>
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<tr>
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<td><strong>Single-agent</strong></td>
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Wumpus world characterization

**Observable**? No—only local perception

**Deterministic**? Yes—outcomes exactly specified

**Episodic**? No—sequential at the level of actions

**Static**? Yes—Wumpus and Pits do not move

**Discrete**? Yes

**Single-agent**? Yes—Wumpus is essentially a natural feature
Exploring a wumpus world
Exploring a wumpus world
Exploring a wumpus world

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<th>B</th>
<th>OK</th>
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<tbody>
<tr>
<td>A</td>
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<td>OK</td>
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Chapter 7  15
Exploring a wumpus world
Exploring a wumpus world

Chapter 7 18
Exploring a wumpus world
Exploring a wumpus world
Other tight spots

Breeze in (1,2) and (2,1)  
⇒ no safe actions

Assuming pits uniformly distributed,  
(2,2) has pit w/ prob 0.86, vs. 0.31

Smell in (1,1)  
⇒ cannot move
Other tight spots

Breeze in (1,2) and (2,1)  
⇒  no safe actions

Assuming pits uniformly distributed, (2,2) has pit w/ prob 0.86, vs. 0.31

Smell in (1,1)  
⇒  cannot move

Can use a strategy of coercion:
  shoot straight ahead
  wumpus was there  ⇒  dead  ⇒  safe
  wumpus wasn’t there  ⇒  safe
Logic in general

Logics are formal languages for representing information such that conclusions can be drawn.

Syntax defines the sentences in the language.

Semantics define the “meaning” of sentences; i.e., define truth of a sentence in a world.

E.g., the language of arithmetic

$x + 2 \geq y$ is a sentence; $x^2 + y >$ is not a sentence.

$x + 2 \geq y$ is true iff the number $x + 2$ is no less than the number $y$.

$x + 2 \geq y$ is true in a world where $x = 7, y = 1$.

$x + 2 \geq y$ is false in a world where $x = 0, y = 6$. 
Entailment means that one thing follows from another:

\[ KB \models \alpha \]

Knowledge base \( KB \) entails sentence \( \alpha \)
if and only if \( \alpha \) is true in all worlds where \( KB \) is true.

E.g., the KB containing “the Giants won” and “the Reds won” entails “Either the Giants won or the Reds won”

E.g., \( x + y = 4 \) entails \( 4 = x + y \)

Entailment is a relationship between sentences (i.e., syntax) that is based on semantics.

Note: brains process syntax (of some sort)
Models

Logicians typically think in terms of models, which are formally structured worlds with respect to which truth can be evaluated.

We say $m$ is a model of a sentence $\alpha$ if $\alpha$ is true in $m$.

$M(\alpha)$ is the set of all models of $\alpha$.

Then $KB \models \alpha$ if and only if $M(KB) \subseteq M(\alpha)$.

E.g. $KB = \text{Giants won and Reds won}$

$\alpha = \text{Giants won}$
Entailment in the Wumpus World

Situation after detecting nothing in [1,1], moving right, breeze in [2,1]

Consider possible models for ?s assuming only pits

3 Boolean choices $\Rightarrow$ 8 possible models
Wumpus models
$KB = \text{wumpus-world rules} + \text{observations}$
\( KB = \text{wumpus-world rules} + \text{observations} \)

\( \alpha_1 = \text{“[1,2] is safe”}, \ KB \models \alpha_1, \text{proved by model checking} \)
\[ KB = \text{wumpus-world rules} + \text{observations} \]
\( KB = \text{wumpus-world rules } + \text{ observations} \)

\( \alpha_2 = \text{“[2,2] is safe”}, \ KB \not\models \alpha_2 \)
Inference

\( KB \vdash_i \alpha = \) sentence \( \alpha \) can be derived from \( KB \) by procedure \( i \)

Consequences of \( KB \) are a haystack; \( \alpha \) is a needle.
Entailment = needle in haystack; inference = finding it

**Soundness:** \( i \) is sound if
whenever \( KB \vdash_i \alpha \), it is also true that \( KB \models \alpha \)

**Completeness:** \( i \) is complete if
whenever \( KB \models \alpha \), it is also true that \( KB \vdash_i \alpha \)

Preview: we will define a logic (first-order logic) which is expressive enough
to say almost anything of interest, and for which there exists a sound and
complete inference procedure.

That is, the procedure will answer any question whose answer follows from
what is known by the \( KB \).
Propositional logic is the simplest logic—illustrates basic ideas

The proposition symbols $P_1$, $P_2$ etc are sentences

If $S$ is a sentence, $\neg S$ is a sentence (negation)

If $S_1$ and $S_2$ are sentences, $S_1 \land S_2$ is a sentence (conjunction)

If $S_1$ and $S_2$ are sentences, $S_1 \lor S_2$ is a sentence (disjunction)

If $S_1$ and $S_2$ are sentences, $S_1 \Rightarrow S_2$ is a sentence (implication)

If $S_1$ and $S_2$ are sentences, $S_1 \Leftrightarrow S_2$ is a sentence (biconditional)
Propositional logic: Semantics

Each model specifies true/false for each proposition symbol

E.g. \( P_{1,2} \), \( P_{2,2} \), \( P_{3,1} \)
\[
\begin{array}{ccc}
\text{true} & \text{true} & \text{false}
\end{array}
\]
(With these symbols, 8 possible models, can be enumerated automatically.)

Rules for evaluating truth with respect to a model \( m \):

\[
\begin{align*}
\neg S & \text{ is true iff } S \text{ is false} \\
S_1 \land S_2 & \text{ is true iff } S_1 \text{ is true and } S_2 \text{ is true} \\
S_1 \lor S_2 & \text{ is true iff } S_1 \text{ is true or } S_2 \text{ is true} \\
S_1 \Rightarrow S_2 & \text{ is true iff } S_1 \text{ is false or } S_2 \text{ is true} \\
\text{i.e., is false iff } & S_1 \text{ is true and } S_2 \text{ is false} \\
S_1 \Leftrightarrow S_2 & \text{ is true iff } S_1 \Rightarrow S_2 \text{ is true and } S_2 \Rightarrow S_1 \text{ is true}
\end{align*}
\]

Simple recursive process evaluates an arbitrary sentence, e.g.,
\[
\neg P_{1,2} \land (P_{2,2} \lor P_{3,1}) = true \land (false \lor true) = true \land true = true
\]
### Truth tables for connectives

<table>
<thead>
<tr>
<th>$P$</th>
<th>$Q$</th>
<th>$\neg P$</th>
<th>$P \land Q$</th>
<th>$P \lor Q$</th>
<th>$P \Rightarrow Q$</th>
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Wumpus world sentences

Let $P_{i,j}$ be true if there is a pit in $[i, j]$. Let $B_{i,j}$ be true if there is a breeze in $[i, j]$.

$\neg P_{1,1}$
$\neg B_{1,1}$
$B_{2,1}$

“Pits cause breezes in adjacent squares”
Wumpus world sentences

Let $P_{i,j}$ be true if there is a pit in $[i, j]$.
Let $B_{i,j}$ be true if there is a breeze in $[i, j]$.

$\neg P_{1,1}$
$\neg B_{1,1}$
$B_{2,1}$

“Pits cause breezes in adjacent squares”

$B_{1,1} \iff (P_{1,2} \lor P_{2,1})$
$B_{2,1} \iff (P_{1,1} \lor P_{2,2} \lor P_{3,1})$

“A square is breezy if and only if there is an adjacent pit”
## Truth tables for inference

<table>
<thead>
<tr>
<th>$B_{1,1}$</th>
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Enumerate rows (different assignments to symbols),
if $KB$ is true in row, check that $\alpha$ is too.
Inference by enumeration

Depth-first enumeration of all models is sound and complete

**Function** \( \text{TT-Entails?}(KB, \alpha) \) \text{ returns } true or false

**Inputs:**
- \( KB \), the knowledge base, a sentence in propositional logic
- \( \alpha \), the query, a sentence in propositional logic

**Symbols:** a list of the proposition symbols in \( KB \) and \( \alpha \)

**Return** \( \text{TT-Check-All}(KB, \alpha, \text{symbols}, []) \)

**Function** \( \text{TT-Check-All}(KB, \alpha, \text{symbols}, \text{model}) \) \text{ returns } true or false

if \( \text{Empty?}(\text{symbols}) \) then

if \( \text{PL-True?}(KB, \text{model}) \) then return \( \text{PL-True?}(\alpha, \text{model}) \)
else return true

else do

\( P \leftarrow \text{First}(\text{symbols}); \text{ rest } \leftarrow \text{Rest}(\text{symbols}) \)

return \( \text{TT-Check-All}(KB, \alpha, \text{rest}, \text{Extend}(P, \text{true}, \text{model})) \) \text{ and }
\( \text{TT-Check-All}(KB, \alpha, \text{rest}, \text{Extend}(P, \text{false}, \text{model})) \)

\( O(2^n) \) for \( n \) symbols; problem is co-NP-complete
Logical equivalence

Two sentences are **logically equivalent** iff true in same models:
\[
\alpha \equiv \beta \text{ if and only if } \alpha \models \beta \text{ and } \beta \models \alpha
\]

\[
\begin{align*}
(\alpha \land \beta) & \equiv (\beta \land \alpha) \quad \text{commutativity of } \land \\
(\alpha \lor \beta) & \equiv (\beta \lor \alpha) \quad \text{commutativity of } \lor \\
((\alpha \land \beta) \land \gamma) & \equiv (\alpha \land (\beta \land \gamma)) \quad \text{associativity of } \land \\
((\alpha \lor \beta) \lor \gamma) & \equiv (\alpha \lor (\beta \lor \gamma)) \quad \text{associativity of } \lor \\
\neg(\neg\alpha) & \equiv \alpha \quad \text{double-negation elimination} \\
(\alpha \Rightarrow \beta) & \equiv (\neg \beta \Rightarrow \neg \alpha) \quad \text{contraposition} \\
(\alpha \Rightarrow \beta) & \equiv (\neg \alpha \lor \beta) \quad \text{implication elimination} \\
(\alpha \iff \beta) & \equiv ((\alpha \Rightarrow \beta) \land (\beta \Rightarrow \alpha)) \quad \text{biconditional elimination} \\
\neg(\alpha \land \beta) & \equiv (\neg \alpha \lor \neg \beta) \quad \text{De Morgan} \\
\neg(\alpha \lor \beta) & \equiv (\neg \alpha \land \neg \beta) \quad \text{De Morgan} \\
(\alpha \land (\beta \lor \gamma)) & \equiv ((\alpha \land \beta) \lor (\alpha \land \gamma)) \quad \text{distributivity of } \land \text{ over } \lor \\
(\alpha \lor (\beta \land \gamma)) & \equiv ((\alpha \lor \beta) \land (\alpha \lor \gamma)) \quad \text{distributivity of } \lor \text{ over } \land
\end{align*}
\]
Validity and satisfiability

A sentence is **valid** if it is true in **all** models,
e.g., $True$, $A \lor \neg A$, $A \Rightarrow A$, $(A \land (A \Rightarrow B)) \Rightarrow B$

Validity is connected to inference via the **Deduction Theorem**:
$KB \models \alpha$ if and only if $(KB \Rightarrow \alpha)$ is valid

A sentence is **satisfiable** if it is true in **some** model
e.g., $A \lor B$, $C$

A sentence is **unsatisfiable** if it is true in **no** models
e.g., $A \land \neg A$

Satisfiability is connected to inference via the following:
$KB \models \alpha$ if and only if $(KB \land \neg \alpha)$ is unsatisfiable
i.e., prove $\alpha$ by **reductio ad absurdum**
Proof methods

Proof methods divide into (roughly) two kinds:

Model checking
- truth table enumeration (always exponential in $n$)
- improved backtracking, e.g., Davis–Putnam–Logemann–Loveland
- heuristic search in model space, search for model where $(KB \land \neg \alpha)$ satisfied (unsound but complete)
  - e.g., min-conflicts-like hill-climbing algorithms

Application of inference rules
- Legitimate (sound) generation of new sentences from old
- Proof = a sequence of inference rule applications
  - Can use inference rules as operators in a standard search alg.
- Typically require translation of sentences into a normal form
Horn Clauses

A Horn Clause is a disjunction of literals of which at most one is positive

E.g., $A, \neg B \lor C, \neg D \lor \neg E$

Useful, as they are equivalent to

◊ proposition symbol; or
◊ (conjunction of symbols) $\Rightarrow$ symbol; or
◊ (conjunction of symbols) $\Rightarrow$ false

A natural way to describe facts
Forward and backward chaining

A KB is in **Horn Form** if

KB = conjunction of (restricted) **Horn clauses**

◊ Restricted to **exactly** one positive literal

**Modus Ponens** (for Horn Form): complete for Horn KBs

\[
\alpha_1, \ldots, \alpha_n, \quad \alpha_1 \land \cdots \land \alpha_n \Rightarrow \beta
\]

Can be used with **forward chaining** or **backward chaining**.
These algorithms are very natural and run in **linear** time
Forward chaining

Idea: fire any rule whose premises are satisfied in the $KB$, add its conclusion to the $KB$, until query is found

- $P \Rightarrow Q$
- $L \land M \Rightarrow P$
- $B \land L \Rightarrow M$
- $A \land P \Rightarrow L$
- $A \land B \Rightarrow L$
- $A$
- $B$
Forward chaining algorithm

function PL-FC-Entails? (KB, q) returns true or false

inputs: KB, the knowledge base, a set of propositional Horn clauses
        q, the query, a proposition symbol

local variables: count, a table, indexed by clause, initially the number of premises
                 inferred, a table, indexed by symbol, each entry initially false
                 agenda, a list of symbols, initially the symbols known in KB

while agenda is not empty do
    p ← Pop(agenda)
    unless inferred[p] do
        inferred[p] ← true
        for each Horn clause c in whose premise p appears do
            decrement count[c]
            if count[c] = 0 then do
                if HEAD[c] = q then return true
                Push(HEAD[c], agenda)
    return false
Forward chaining example
Forward chaining example
Forward chaining example
Forward chaining example
Forward chaining example
Forward chaining example
Forward chaining example
Proof of completeness

FC derives every atomic sentence that is entailed by $KB$

1. FC reaches a **fixed point** where no new atomic sentences are derived
2. Consider the final state as a model $m$, assigning true/false to symbols
3. Every clause in the original $KB$ is true in $m$
   **Proof:** Suppose a clause $a_1 \land \ldots \land a_k \Rightarrow b$ is false in $m$
   Then $a_1 \land \ldots \land a_k$ is true in $m$ and $b$ is false in $m$
   Therefore the algorithm has not reached a fixed point!
4. Hence $m$ is a model of $KB$
5. If $KB \models q$, $q$ is true in **every** model of $KB$, including $m$
Backward chaining

Idea: work backwards from the query $q$:
   to prove $q$ by BC,
   check if $q$ is known already, or
   prove by BC all premises of some rule concluding $q$

Avoid loops: check if new subgoal is already on the goal stack

Avoid repeated work: check if new subgoal
   1) has already been proved true, or
   2) has already failed
Backward chaining example
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Backward chaining example

Diagram of backward chaining example with nodes A, B, L, M, P, and Q.
Backward chaining example
Forward vs. backward chaining

FC is data-driven, cf. automatic, unconscious processing,
  e.g., object recognition, routine decisions

May do lots of work that is irrelevant to the goal

BC is goal-driven, appropriate for problem-solving,
  e.g., Where are my keys? How do I get into a PhD program?

Complexity of BC can be much less than linear in size of KB
Resolution

Conjunctive Normal Form (CNF—universal)

conjunction of disjunctions of literals

clauses

E.g., \((A \lor \neg B) \land (B \lor \neg C \lor \neg D)\)

Resolution inference rule (for CNF): complete for propositional logic

\[
\ell_1 \lor \cdots \lor \ell_k, \quad m_1 \lor \cdots \lor m_n
\]

\[
\ell_1 \lor \cdots \lor \ell_{i-1} \lor \ell_{i+1} \lor \cdots \lor \ell_k \lor m_1 \lor \cdots \lor m_{j-1} \lor m_{j+1} \lor \cdots \lor m_n
\]

where \(\ell_i\) and \(m_j\) are complementary literals. E.g.,

\[
P_{1,3} \lor P_{2,2}, \quad \neg P_{2,2}
\]

\[
P_{1,3}
\]

Resolution is sound and complete for propositional logic
Conversion to CNF

\[ B_{1,1} \iff (P_{1,2} \lor P_{2,1}) \]

1. Eliminate \( \iff \), replacing \( \alpha \iff \beta \) with \( (\alpha \Rightarrow \beta) \land (\beta \Rightarrow \alpha) \).

\[ (B_{1,1} \Rightarrow (P_{1,2} \lor P_{2,1})) \land ((P_{1,2} \lor P_{2,1}) \Rightarrow B_{1,1}) \]

2. Eliminate \( \Rightarrow \), replacing \( \alpha \Rightarrow \beta \) with \( \neg \alpha \lor \beta \).

\[ (\neg B_{1,1} \lor P_{1,2} \lor P_{2,1}) \land (\neg (P_{1,2} \lor P_{2,1}) \lor B_{1,1}) \]

3. Move \( \neg \) inwards using de Morgan’s rules and double-negation:

\[ (\neg B_{1,1} \lor P_{1,2} \lor P_{2,1}) \land ((\neg P_{1,2} \land \neg P_{2,1}) \lor B_{1,1}) \]

4. Apply distributivity law (\( \lor \) over \( \land \)) and flatten:

\[ (\neg B_{1,1} \lor P_{1,2} \lor P_{2,1}) \land (\neg P_{1,2} \lor B_{1,1}) \land (\neg P_{2,1} \lor B_{1,1}) \]
Resolution algorithm

Proof by contradiction, i.e., show $KB \land \neg \alpha$ unsatisfiable

function PL-Resolution($KB, \alpha$) returns true or false

inputs: $KB$, the knowledge base, a sentence in propositional logic
        $\alpha$, the query, a sentence in propositional logic

$clauses \leftarrow$ the set of clauses in the CNF representation of $KB \land \neg \alpha$
$new \leftarrow \{\}$

loop do
    for each $C_i, C_j$ in $clauses$ do
        $resolvents \leftarrow$ PL-Resolve($C_i, C_j$)
        if $resolvents$ contains the empty clause then return true
        $new \leftarrow new \cup resolvents$
        if $new \subseteq clauses$ then return false
    $clauses \leftarrow clauses \cup new$
Resolution example

\[ KB = (B_{1,1} \iff (P_{1,2} \lor P_{2,1})) \land \neg B_{1,1} \quad \alpha = \neg P_{1,2} \]
Logical agents apply inference to a knowledge base to derive new information and make decisions.

Basic concepts of logic:
- **syntax**: formal structure of sentences
- **semantics**: truth of sentences wrt models
- **entailment**: necessary truth of one sentence given another
- **inference**: deriving sentences from other sentences
- **soundness**: derivations produce only entailed sentences
- **completeness**: derivations can produce all entailed sentences

Wumpus world requires the ability to represent partial and negated information, reason by cases, etc.

Forward, backward chaining are linear-time, complete for Horn clauses. Resolution is complete for propositional logic.

Propositional logic lacks expressive power.