CMPT 225 - Midterm 2
November 17, 2017
Time: 45 minutes

DO NOT begin until told to do so.
Fill in your name and student ID below.
Turn off your phone and put it in your bag.
Keep your student ID card on your desk at all times.
You can't have anything other than a pen/pencil and your student ID card on you during the test.
Before you begin, please check that your exam consists of 7 consecutively numbered questions and 2 pages of appendix.

Read the questions carefully and more than once.
Good luck!

NAME: ........................................................................................................................................

STUDENT ID: ..............................................................................................................................

SCORE:

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1. Give the in-order traversal of the following tree (5 marks)

```
  10
   / 
  8   11
  /   / 
 2   3 9  14
 / /  / /  
1 6 13 16
  / 
4   5
```

in-order traversal:
1, 2, 4, 6, 5, 8, 3, 9, 11, 13, 14, 16

2. True/False? (5 marks)
   a. the above tree is a min-heap F
   b. the above tree is complete F
   c. the above tree is proper T
   d. the above tree is a binary search tree F
   e. the height of the above tree is 5 F

3. What will the following heap look like after you insert 17 into it? (5 marks)

```
  6
 /   
8     10
 /     /  
24    16 18
/ /  
50 38 20 26 22
```

```
Answer:

4. What is the time complexity of the following in big-O notation in the worst case scenario? You do not need to justify your answers.  

   a. getting the min() from a priority queue that is based on a sorted list  \( O(1) \)

   b. insertion into a priority queue that is based on a sorted list  \( O(n) \)

   c. bottom-up construction of a heap from a list of items  \( O(n) \)

   d. removeMin() in a heap-based priority queue  \( O(\lg n) \)

   e. insertion into a binary search tree  \( O(n) \)

5. What is the time complexity of the following functions in big-O notation? You do not need to justify your answers.

   void mystery(int n){
       if (n>1){
           mystery(n/2);
           mystery(n/2);
       }
       cout << n << endl;
   }

   \( O(n) \)
void mystery2(int n) {
    for (int i=0; i<10; i++)
        for (int j=i+2; j<i+7; j++)
            cout << "*";
    cout << endl;
}

O(1)

6. Give the code for the insert() function in the HeapPriorityQueue. The rest of the class implementation is provided in appendix A for your reference. (10 marks)

```cpp
template <typename E, typename C> // insert element
void HeapPriorityQueue::insert(const E& e)
{
    T.addLast(e); // add e to heap
    Position v = T.last(); // e's position
    while (!T.isRoot(v)) { // up-heap bubbling
        Position u = T.parent(v);
        if (!isLess(*v, *u)) break; // if v in order, we're done
        T.swap(v, u); // . . .else swap with parent
        v = u;
    }
}
```

7. Describe (in pseudo-code or in fewer than five sentences) an algorithm that finds the kth smallest element of a vector of n distinct integers in $O(n + k \log(n))$ time. (10 marks)

build a heap from the items in the vector using bottom-up heap construction $O(n)$
repeatedly remove the min (k-1 times) $O(k \log(n))$
kth smallest element is now at the root. $O(1)$
Appendix A. Implementation of HeapPriorityQueue

template <typename E, typename C>
class HeapPriorityQueue {
public:
    int size() const { return T.size(); } // number of elements
    bool empty() const { return size() == 0; } // is the queue empty?

void insert(const E& e) /* You will fill this.*
    const E& min() { return *(T.root()); } // minimum element

void removeMin(){
    if (size() == 1) // only one node?
        T.removeLast(); // ...remove it
    else {
        Position u = T.root(); // root position
        T.swap(u, T.last()); // swap last with root
        T.removeLast(); // ...and remove last
        while (T.hasLeft(u)) { // down-heap bubbling
            Position v = T.left(u);
            if (T.hasRight(u) && C::isLess(*(T.right(u)), *v))
                v = T.right(u); // v is u's smaller child
            if (C::isLess(*v, *u)) { // is u out of order?
                T.swap(u, v); // ...then swap
                u = v;
            }
        }
        else break; // else we're done
    }
}

private:
    VectorCompleteTree T; // priority queue contents - Implemented in the next page
    C isLess; // less-than comparator
    // shortcut for tree position
    typedef typename VectorCompleteTree::Position Position;
};
Apendix A (continued). Implementation of VectorCompleteTree used in HeapPriorityQueue

template <typename E>
class VectorCompleteTree {
private: // member data
    std::vector V; // tree contents
public: // publicly accessible types
    typedef typename std::vector::iterator Position; // a position in the tree
protected: // protected utility functions
    Position pos(int i) // map an index to a position
    { return V.begin() + i; }
    int idx(const Position& p) const // map a position to an index
    { return p - V.begin(); }
public: VectorCompleteTree() : V(1) {} // constructor
    int size() const { return V.size() - 1; }
    Position left(const Position& p) { return pos(2*idx(p)); }
    Position right(const Position& p) { return pos(2*idx(p) + 1); }
    Position parent(const Position& p) { return pos(idx(p)/2); }
    bool hasLeft(const Position& p) const { return 2*idx(p) <= size(); }
    bool hasRight(const Position& p) const { return 2*idx(p) + 1 <= size(); }
    bool isRoot(const Position& p) const { return idx(p) == 1; }
    Position root() { return pos(1); }
    Position last() { return pos(size()); }
    void addLast(const E& e) { V.push_back(e); }
    void removeLast() { V.pop_back(); }
    void swap(Position& p, const Position& q) { E e = *q; *q = *p; *p = e; }
};